



Strategies & Obstacles in Converting a Large Production Application to FORTRAN 90

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May 19, 1999

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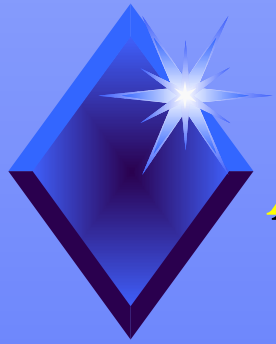
Topics of Discussion

- ◆ Intro. - Rationale for converting to FORTRAN 90
- ◆ Advantages
- ◆ Migration strategies
- ◆ Debugging techniques
- ◆ Obstacles / examples / clean-up
- ◆ Tools used
- ◆ System verification
- ◆ Resources & conclusions



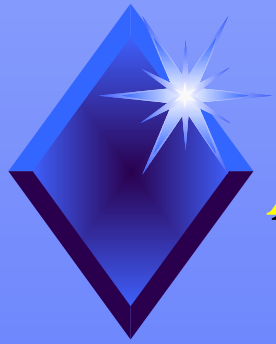
Rationale for Converting to f90

- ◆ Poor performance of PV+ CPUs
- ◆ CMOS technology - GigaFlop performance
- ◆ Phaseout of T90s and support for them
- ◆ Availability of spares - reliability
- ◆ Vendor limited support of FORTRAN 77



Advanatage of FORTRAN 90

- ◆ FORTRAN 90 is vendor standard
- ◆ Upward compatible libraries & object files
- ◆ Dynamic memory allocation
- ◆ Array operations (syntax)
- ◆ More intrinsic procedures
- ◆ Derived data types



Advantages (continued)

- ◆ Performance improvements
 - ◆ Improved vectorization
 - ◆ Unrolling of loops
 - ◆ More in-lining of code
- ◆ Software can be simplified
- ◆ Reduced maintenance cost
- ◆ Improved portability



Migration Strategies

- ◆ Talk with other sites already using f90
- ◆ Convert small-modern applications first
- ◆ Verify external f90 libraries
- ◆ Use FORTRAN 90 compiler to locate noncompliances
 - ◆ Triton
 - ◆ Workstations (IBM RS6000, HPs)



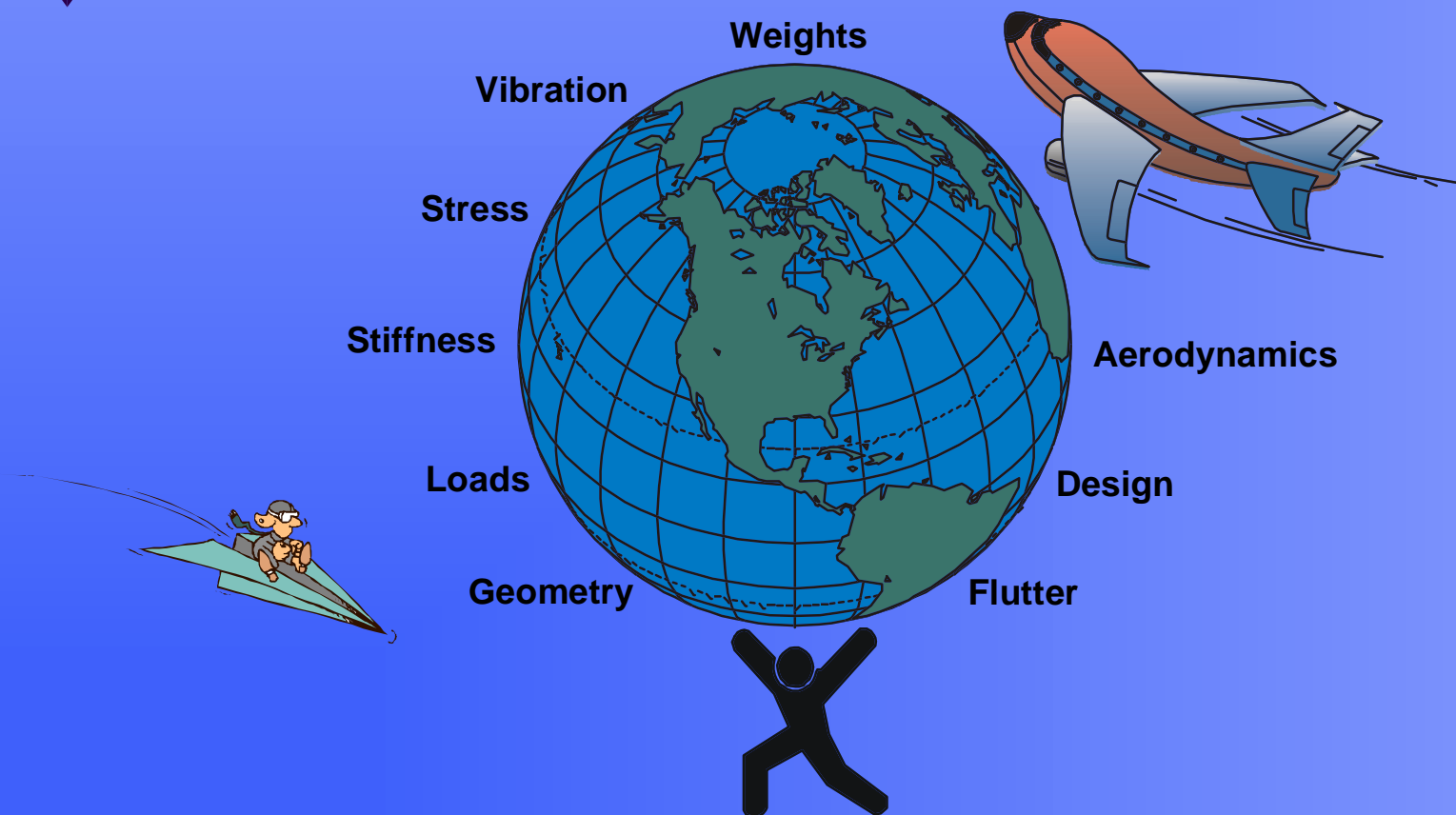
Migration Strategies (continued)

- ◆ Subdivide large applications
 - ◆ Support libraries
 - ◆ Selective loading/testing
 - ◆ Precompilers
 - ◆ Preprocessors, processors, postprocessors
 - ◆ Utilities (e.g. third party interfaces)
- ◆ Set number of CPUs 1



ATLAS

An Integrated Structural Analysis and Design System
(1.4 Million Lines of Code)





Debugging Techniques

- ◆ Address one problem at a time
- ◆ Interactive debugger (Totalview) for aborts
- ◆ Try to duplicate problem on cft77 system
 - ◆ Successful
 - ◆ Our code changed or
 - ◆ System libraries changed
- ◆ Isolate the f90 routine
 - ◆ Mix of f90 and cft77 objects



Debugging Techniques (cont.)

- ◆ Restrict or eliminate optimization
- ◆ Check incoming & outgoing arguments
- ◆ Split routine in question into several
 - ◆ Mix of f90 and cft77 objects
- ◆ Use Totalview
 - ◆ Step through f90 version
 - ◆ Compare with f77 version



Obstacles Encountered (Examples)

```
dimension a(10,3), b(9), c(10,9)
```

```
      .  
call vecadd (a ,10 ,b ,3)
```

```
      .  
      .  
call vecadd (c, 10, b, 6)
```

```
-----  
subroutine vecadd (x, nrow, b, num)
```

```
dimension x (nrow,3), b (num)
```

```
do i=1, num
```

```
  x (nrow, i) = x (nrow, i) + b (i)
```

```
enddo
```

```
return
```



Examples (optimization problems)

dimension $a(1) <> a(*) <> a(n)$

dimension $b(n,1) <> b(n,*) <> b(n,m)$

dimension $c(n,3) <> c(n,*) <> c(n,m)$

dimension $d(1) <> d(n*m) <> d(n,m)$

dimension $e(n,m,1) <> e(n,m,*) <> e(n,m,k)$



Examples (optimization continued)

```
dimension ifile(1)
equivalence( ifile, arnf)
common / kqrndm / arnf, brnf, ... , zrnf
```

```
do i = 1, n      <> call dropfil (arnf, n)
  close (ifile (i) )  where:
enddo            subroutine dropfil(ifiles,n)
                  dimension ifiles (n)
                  do i=1, n
                    close ( ifiles (i) )
                  enddo
```



Examples (optimization continued)

dimension a (n,m), b (n*m)

```
do i= 1, n*m
```

```
  a ( i ,1) = b(i)
```

```
enddo
```

| unpredictable results

```
do j =1 , m
```

```
  do i = 1, n
```

```
    a ( i, j ) = b ( i + (j-1)*n )
```

```
  enddo
```

```
enddo
```

```
do j=1,m
```

```
<>  a(:,j) = b((j-1)*n+1:)
```

```
enddo
```



More Obstacles Encountered

- ◆ Loop error for variables with L - format
(e.g. 3Labc ... 32 bit loop register)
- ◆ `keybig = -mask(1) <> JMHCON(3)`
- ◆ Missing routine arguments
- ◆ `call writms (ntp8, nsizeb, 240, 3)`
 - ◆ `call writms (ntp8, nsizeb, 240, 3, irr)`
 - ◆ `call writms (ntp8, nsizeb, 240, 3, -1, 0)`



Obstacles Encountered (compiler)

- ◆ Formats

- ◆ $3x5e16.8$ \leftrightarrow $3x, 5e16.8$

- ◆ Dimension na (nxt, 5)

- ◆ $\text{int} = \text{na}(\text{nxt})$ \leftrightarrow $\text{int} = \text{na}(\text{nxt}, 1)$

- ◆ Round-off differences

- ◆ Different variable memory locations

- ◆ Common block ordering



Examples (Mixed Arrays)

cft77 equivalence (cntrl(1), icntrl(1))
 rval = icntrl (5) .or. 0
 rval = rmove (icntrl(5))
 rval = cntrl (5)
 ival = rval

f90 rval = transfer (icntrl(5), rval)
 ival = rval

f90 ival = transfer (icntrl(5), rval)



Examples (clean - up)

- ◆ General format clean-up
 - ◆ O22 <> I9
 - ◆ nL <> nH
- ◆ Change pointers to allocatable arrays
- ◆ Replace loops or routine calls with f90 syntax where practical
- ◆ Automatic array allocations



Tools Used / Developed

- ◆ Internal program to process SCCS files
 - ◆ (1), (x, 1)
 - ◆ dimension, real, integer, complex
- ◆ grep (0L, 1L, 2L, ... 9L)
- ◆ f90 Compilers (Triton and RS6000)
- ◆ SCCS
- ◆ Cflist & Totalview



Regression Testing

- ◆ Component testing
 - ◆ Libraries (system, data center, internal)
 - ◆ Preprocessors, postprocessing, processors
- ◆ 256 validation cases out of 443
- ◆ Continuous developer and user testing (over 11 months)
- ◆ Block point release validation (190 cases)



Resources

- ◆ Flow time: March 1998 to February 1999
- ◆ Labor - Hours:
 - ◆ 370 Analyst
 - ◆ 95 Engineering
- ◆ 2302 routines of 6423 modified
- ◆ Triton T916 with 512 MW
- ◆ Minimal machine resource impact



Conclusions

- ◆ Not tested . . . It won't work !
- ◆ Successful conversion
 - ◆ Code executes more efficiently
 - ◆ Discovered many underlying array size errors
 - ◆ Applications are now more portable
 - ◆ Cost reductions
 - ◆ Maintenance (do more with f90 and easier)
 - ◆ Development