

# HPCC STREAM and RA in Chapel Performance and Potential

---

Steve Deitz  
Cray Inc.

## What is Chapel?

- A new parallel language
  - Under development at Cray Inc.
  - Supported through the DARPA HPCS program
- Goals
  - **Improve programmer productivity**
  - Improve the programmability of parallel computers
  - Match or improve performance of MPI/UPC/CAF
  - Provide better portability than MPI/UPC/CAF
  - Improve robustness of parallel codes
  - Support multi-core and multi-node systems

## Outline

- What is Chapel?
- Chapel's Parallel Programming Model
- HPCC STREAM Triad in Chapel
- HPCC RA in Chapel
- Summary and Future Work

## Fragmented vs. Global-View: Definitions

- Programming model

*The mental model of a programmer*

- Fragmented models

*Programmers take point-of-view of a single processor/thread*

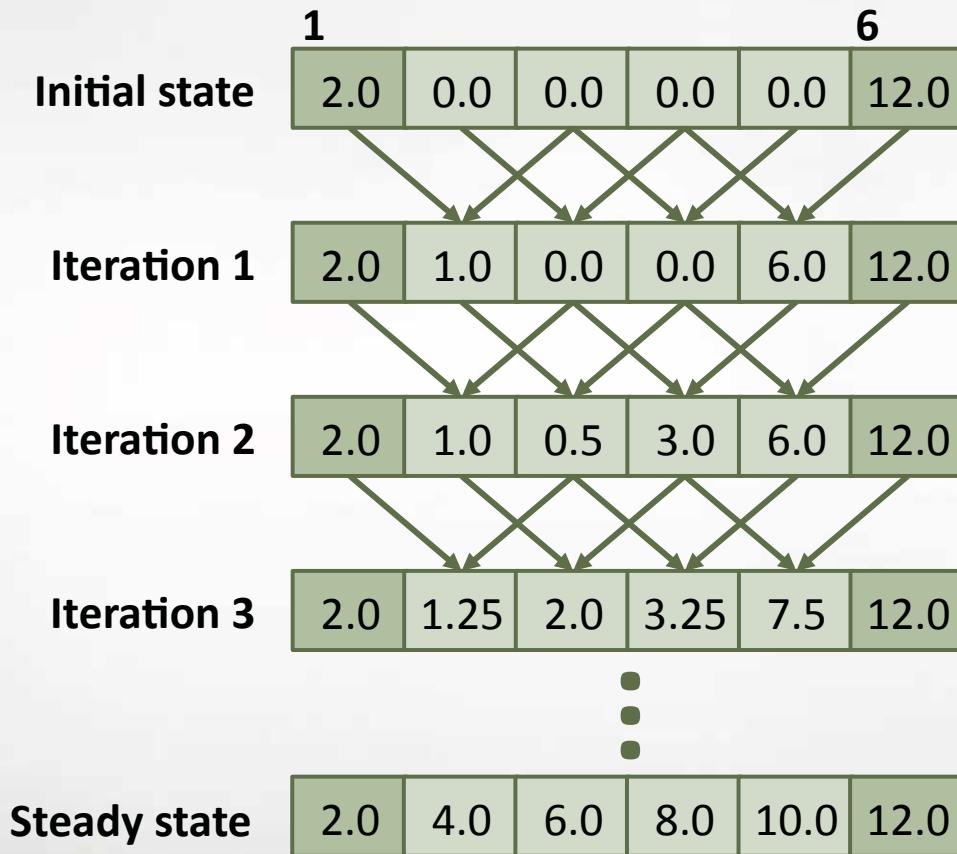
- SPMD models (Single Program, Multiple Data)

*Fragmented models with multiple copies of one program*

- Global-view models

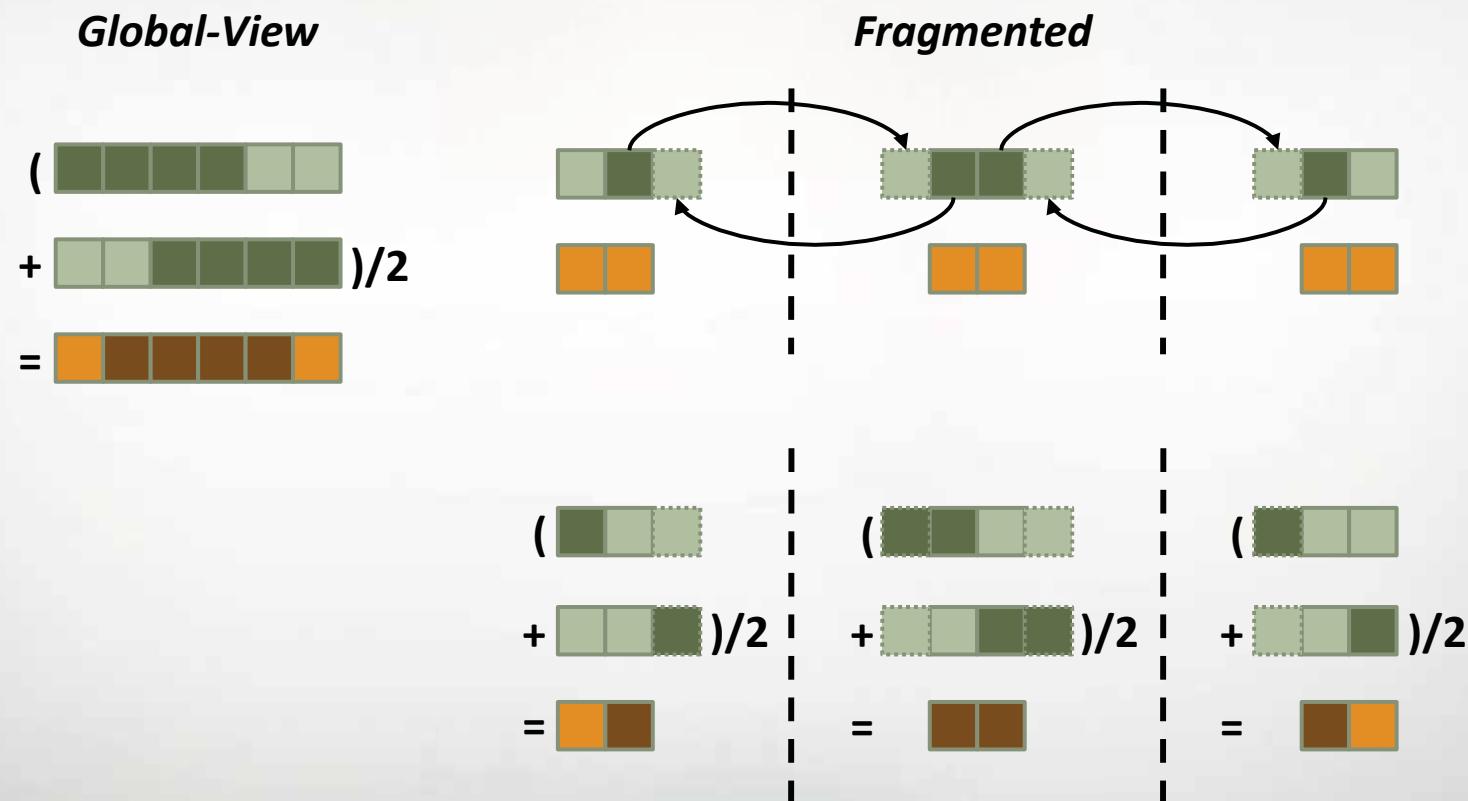
*Programmers write code to describe computation as a whole*

# 3-Point Stencil Example (n=6)



## 3-Point Stencil Example

### Global-View vs. Fragmented Computation



## 3-Point Stencil Example: Code

### Global-View vs. Fragmented Code

#### Global-View

```
def main() {  
    var n = 1000;  
    var A, B: [1..n] real;  
  
    forall i in 2..n-1 do  
        B(i) = (A(i-1)+A(i+1))/2;  
}
```

#### Fragmented

```
def main() {  
    var n = 1000;  
    var me = commID(), p = commProcs(),  
          myN = n/p, myLo = 1, myHi = myN;  
    var A, B: [0..myN+1] real;  
  
    if me < p {  
        send(me+1, A(myN));  
        recv(me+1, A(myN+1));  
    } else myHi = myN-1;  
    if me > 1 {  
        send(me-1, A(1));  
        recv(me-1, A(0));  
    } else myLo = 2;  
    for i in myLo..myHi do  
        B(i) = (A(i-1)+A(i+1))/2;  
}
```

Assumes p divides n

# NAS MG Stencil in Fortran + MPI, in Chapel



```

def rprj3(S, R) {
    const Stencil = [-1..1, -1..1, -1..1],
        W: [0..3] real = (0.5, 0.25, 0.125, 0.0625),
        W3D = [(i,j,k) in Stencil] W((i!=0)+(j!=0)+(k!=0));
    forall inds in S.domain do
        S(inds) =
            + reduce [offset in Stencil] (W3D(offset) *
                R(inds + offset*R.stride)));
}

```

## Outline

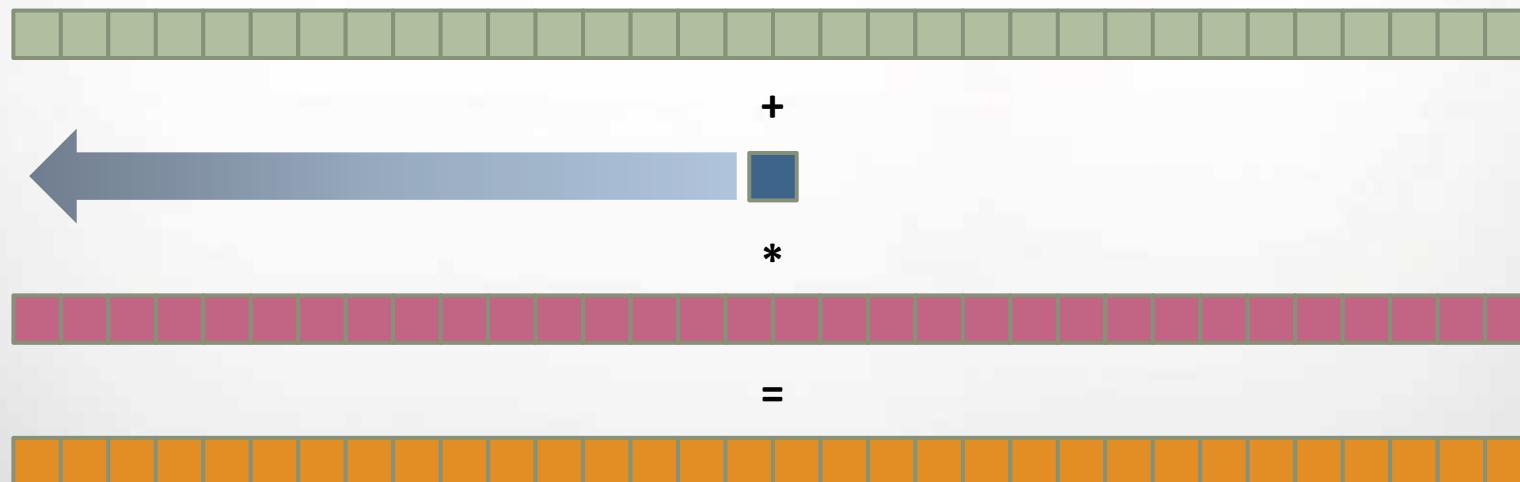
- What is Chapel?
- Chapel's Parallel Programming Model
- HPCC STREAM Triad in Chapel
- HPCC RA in Chapel
- Summary and Future Work

# Introduction to STREAM Triad

Given:  $m$ -element vectors  $A, B, C$

Compute: `forall i in 1..m do`

$$\mathbf{A}(i) = \mathbf{B}(i) + \alpha * \mathbf{C}(i);$$

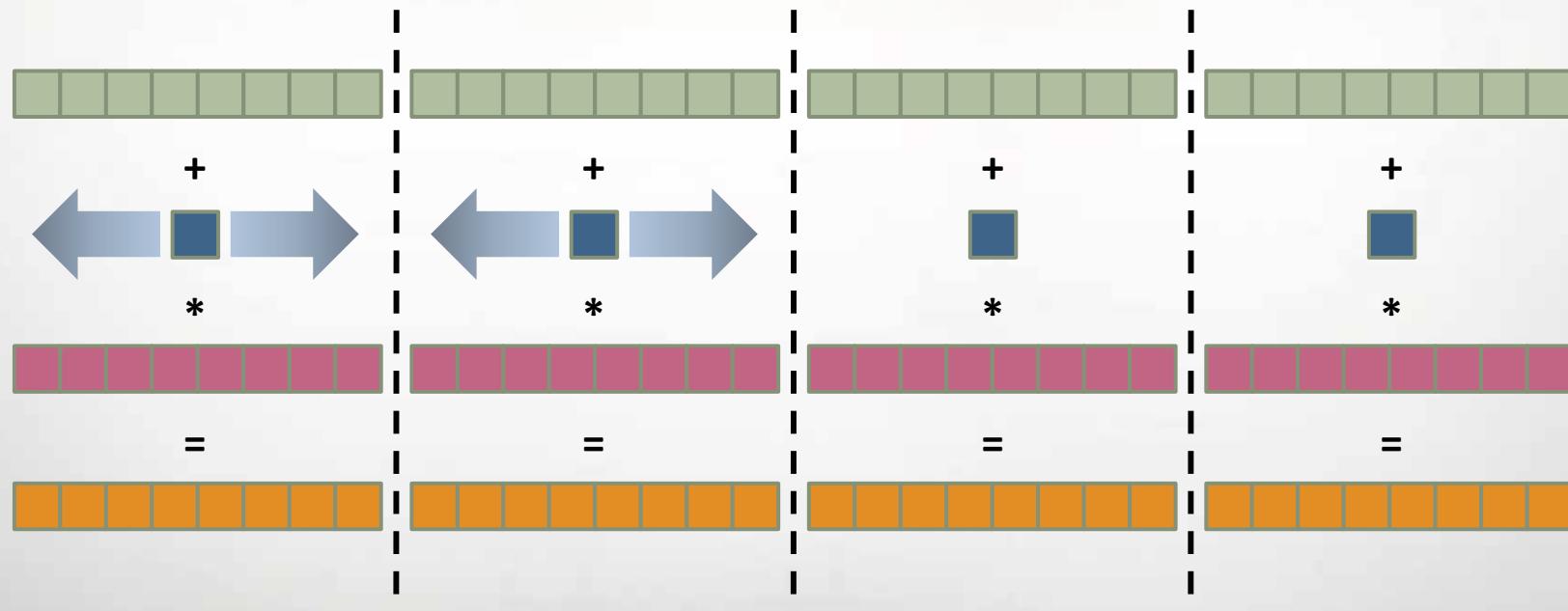


# Distributed Parallelization of STREAM Triad

Given:  $m$ -element vectors  $A, B, C$

Compute: `forall i in 1..m do`

$$\mathbf{A}(i) = \mathbf{B}(i) + \alpha * \mathbf{C}(i);$$

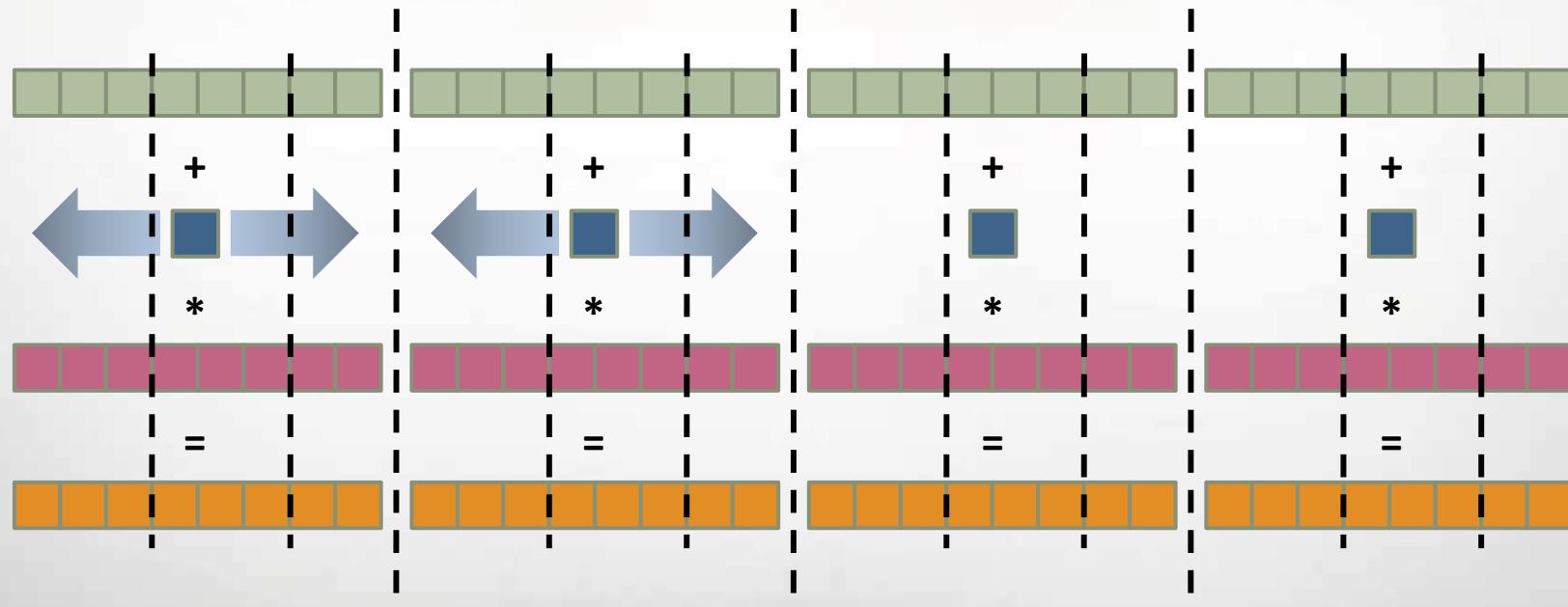


# Further Parallelization of STREAM Triad

Given:  $m$ -element vectors  $A, B, C$

Compute: `forall i in 1..m do`

$$\textcolor{orange}{A}(i) = \textcolor{green}{B}(i) + \alpha * \textcolor{red}{C}(i);$$



# STREAM Triad in Chapel: Single Locale

Given:  $m$ -element vectors  $A, B, C$

Compute: **forall**  $i$  **in**  $1..m$  **do**  
     $\text{A}(i) = \text{B}(i) + \alpha * \text{C}(i);$

```
config const m: int(64) = ...;
const alpha: real = 3.0;
const ProblemSpace: domain(1,int(64)) = [1..m];
var A, B, C: [ProblemSpace] real;

forall i in ProblemSpace do
    A(i) = B(i) + alpha * C(i);
```

# STREAM Triad in Chapel: Single Locale

Given:  $m$ -element vectors  $A, B, C$

Compute: `forall i in 1..m do`  
`A(i) = B(i) + alpha * C(i);`

```
config const m: int(64) = ...;
const alpha: real = 3.0;
const ProblemSpace: domain(1,int(64)) = [1..m];
var A, B, C: [ProblemSpace] real;
```

`A = B + alpha * C;`

More concise variation  
using whole array operations

# STREAM Triad in Chapel: Single Locale

Given:  $m$ -element vectors  $A, B, C$

Compute: **forall**  $i$  **in**  $1..m$  **do**  
 $\text{A}(i) = \text{B}(i) + \alpha * \text{C}(i);$

```
config const m: int(64) = ...;
const alpha: real = 3.0;
const ProblemSpace: domain(1,int(64)) = [1..m];
var A, B, C: [ProblemSpace] real;
```

```
forall (a,b,c) in (A,B,C) do
    a = b + alpha * c;
```

Variation that iterates  
directly over the arrays

# STREAM Triad in Chapel: Multi-Locale

Given:  $m$ -element vectors  $A, B, C$

Compute: `forall i in 1..m do`

`A(i) = B(i) +  $\alpha$  * C(i);`

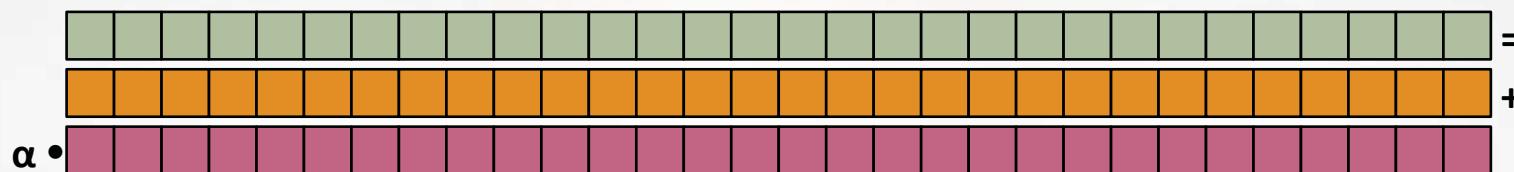
```
config const m: int(64) = ... , tpl = ...;
const alpha: real = 3.0;
const BlockDist = new Block(1,int(64),[1..m],tpl);
const ProblemSpace: domain(1, int(64))
                      distributed BlockDist = [1..m];
var A, B, C: [ProblemSpace] real;

forall (a,b,c) in (A,B,C) do
  a = b + alpha * c;
```

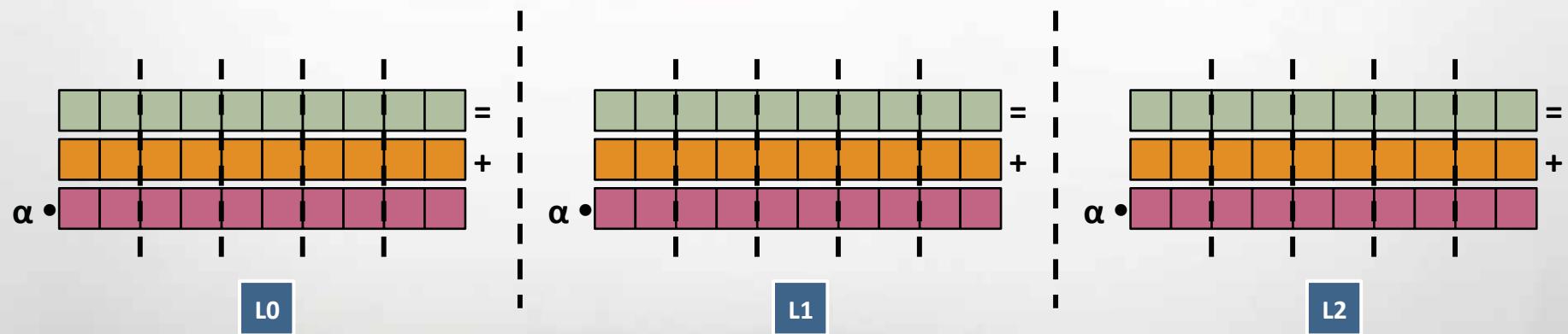
# What is a Distribution?

A “recipe” for distributed arrays that...

Instructs the compiler how to Map the global view...



...to a fragmented, per-processor implementation



# STREAM Triad in Chapel: Multi-Locale

Given:  $m$ -element vectors  $A, B, C$

Compute: `forall i in 1..m do`

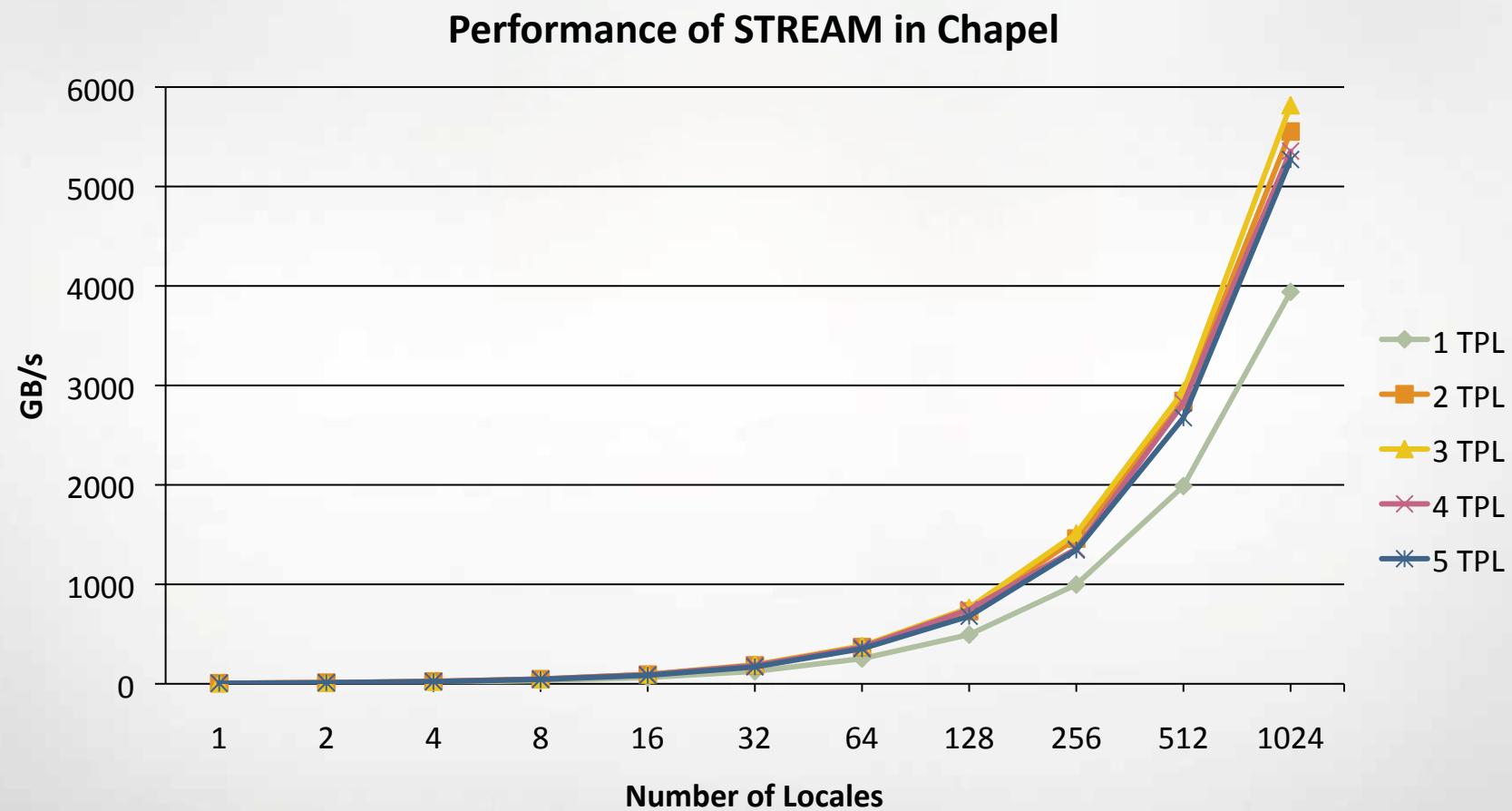
`A(i) = B(i) +  $\alpha$  * C(i);`

```
config const m: int(64) = ... , tpl = ...;
const alpha: real = 3.0;
const BlockDist = new Block(1,int(64),[1..m],tpl);
const ProblemSpace: domain(1, int(64))
                      distributed BlockDist = [1..m];
var A, B, C: [ProblemSpace] real;

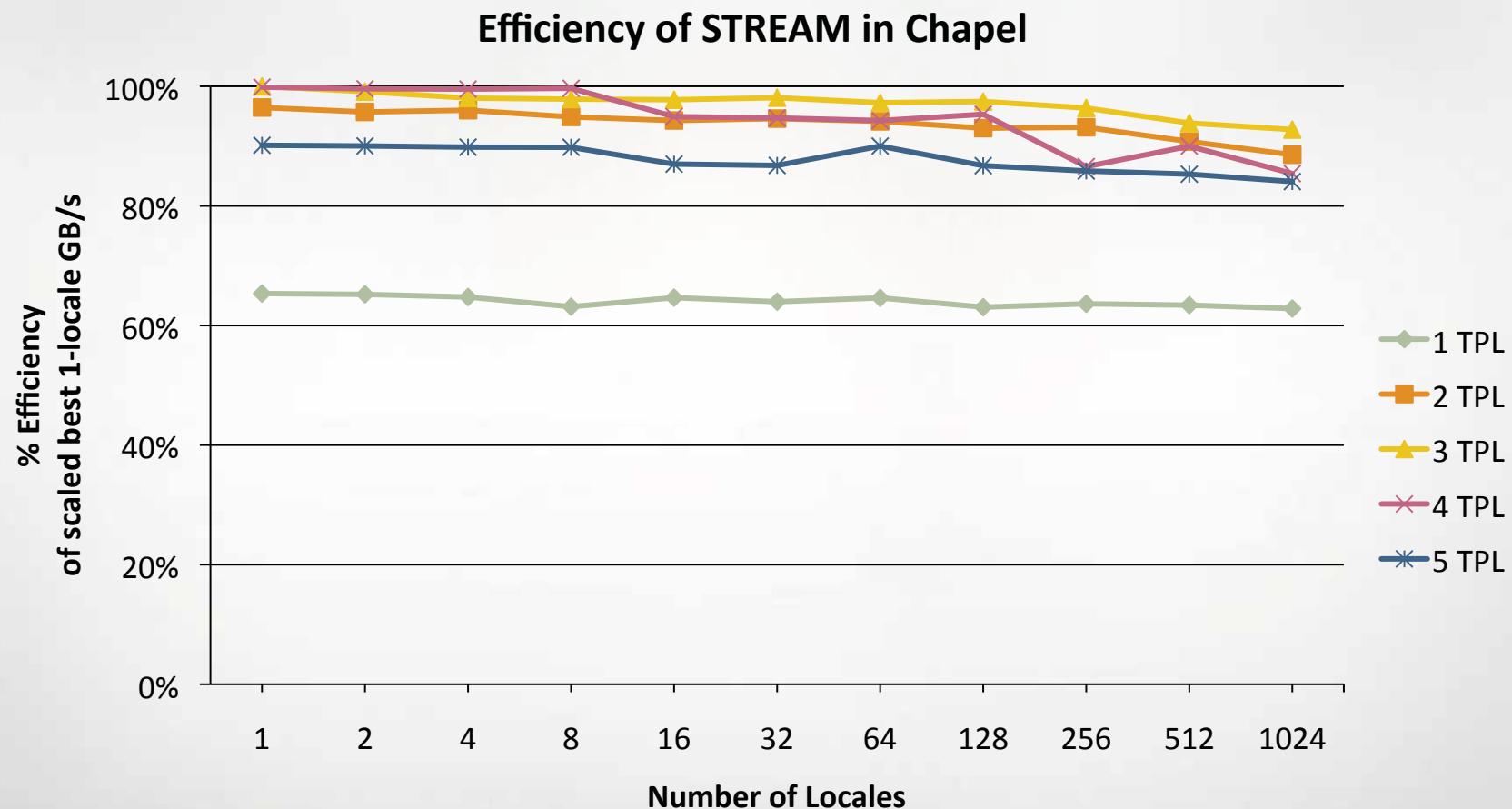
forall (a,b,c) in (A,B,C) do
  a = b + alpha * c;
```

# HPCC STREAM Performance

**CRAY**  
THE SUPERCOMPUTER COMPANY



# HPCC STREAM Efficiency



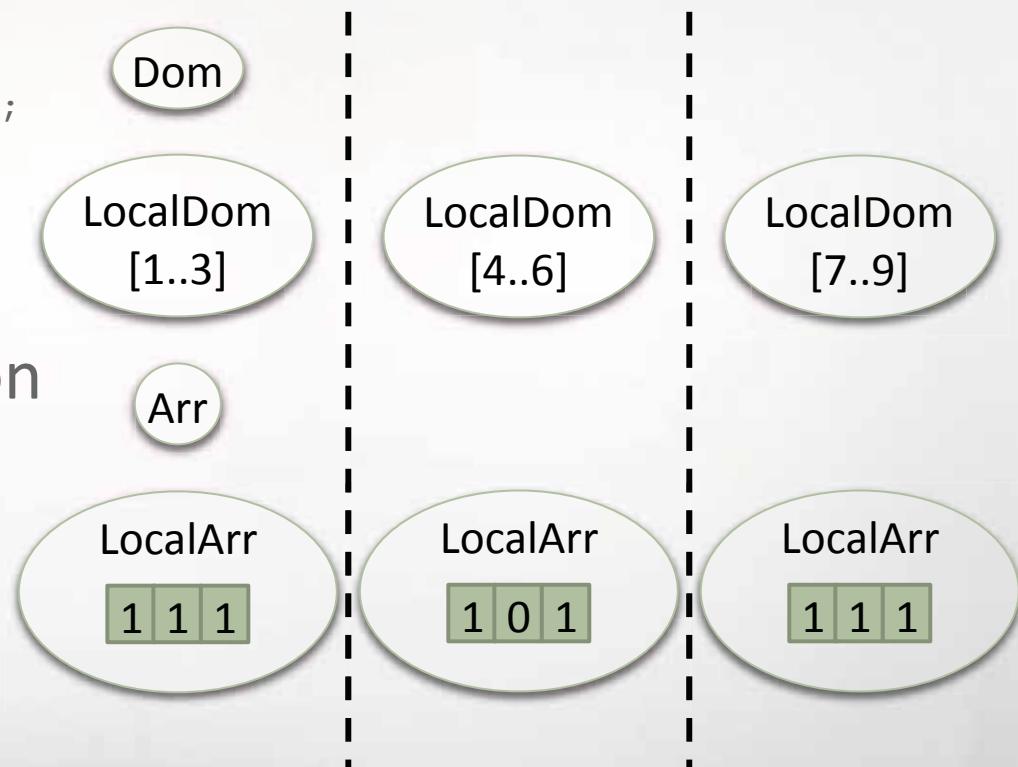
# Optimization: Privatization

## Simple example

```
var Dist: Block(1,int(64));
var Dom: domain(1,int(64))
    distributed Dist;
var Arr: [Dom] int;
```

Reference to local data  
requires communication

```
on Locales(1) {
    Arr(5) = 0;
}
```



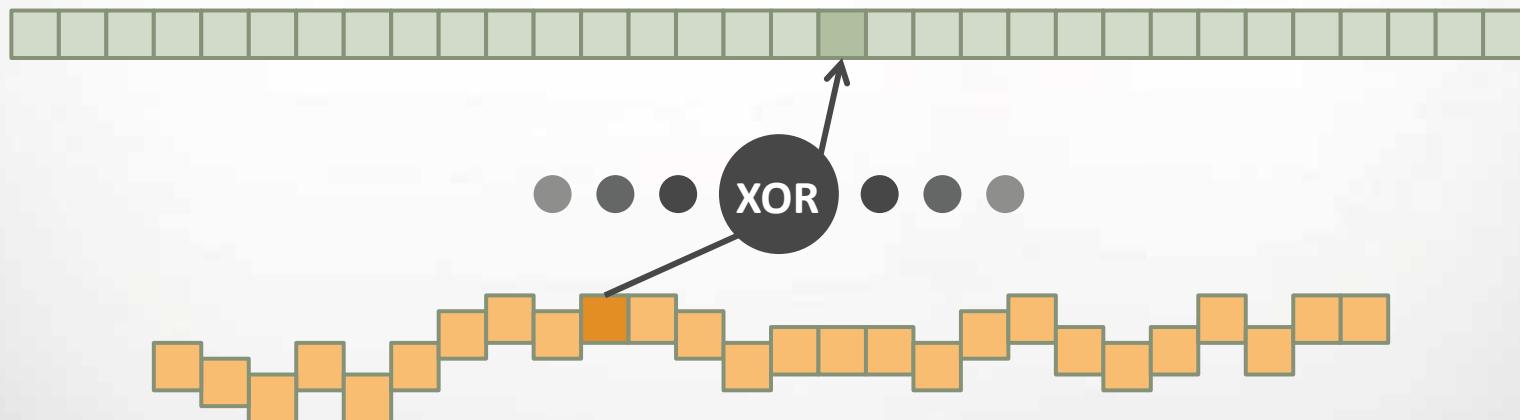
## Outline

- What is Chapel?
- Chapel's Parallel Programming Model
- HPCC STREAM Triad in Chapel
- HPCC RA in Chapel
- Summary and Future Work

# Introduction to RA

Given:  $m$ -element table  $T$  (where  $m = 2^n$ )

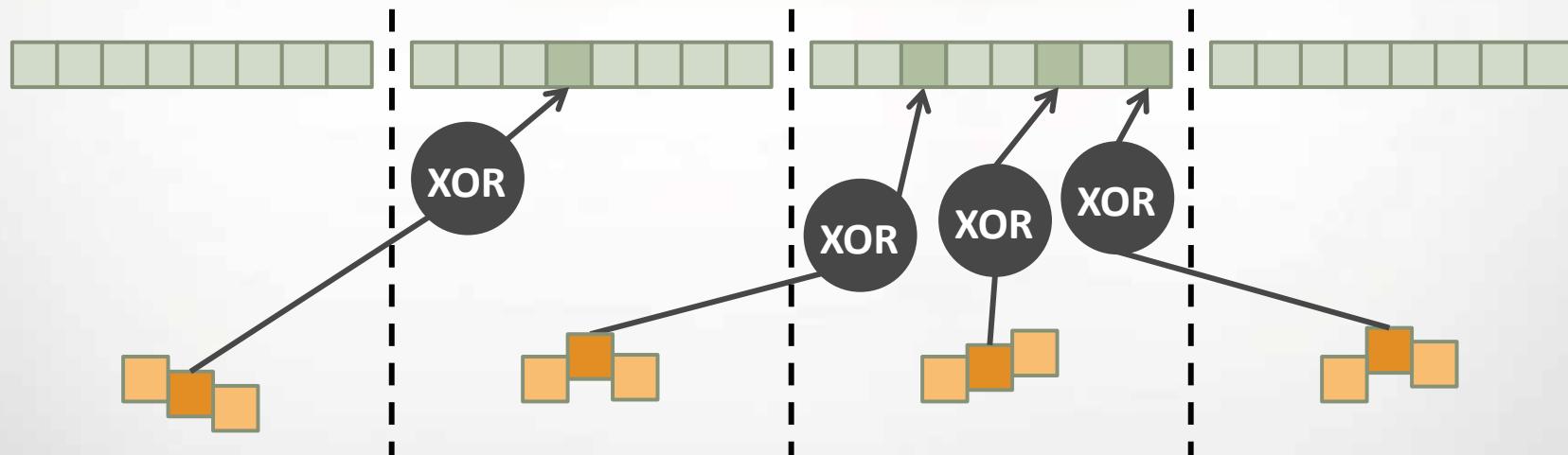
Compute: `forall r in RandomUpdates do`  
`$T(r \& (m-1)) \leftarrow r;$`



# Distributed Parallelization of RA

Given:  $m$ -element table  $T$  (where  $m = 2^n$ )

Compute: `forall r in RandomUpdates do`  
`$T(r \& (m-1)) \leftarrow r;$`



# RA in Chapel: Single Locale

Given:  $m$ -element table  $T$  (where  $m = 2^n$ )

Compute: **forall**  $r$  in RandomUpdates do  
     $T(r \& (m-1)) \wedge= r;$

```
config const m = ..., N_U = ...;
const TableSpace: domain(1,uint(64)) = [0..m-1],
    Updates: domain(1,uint(64)) = [0..N_U-1];
var T: [TableSpace] uint(64);

forall (i,r) in (Updates,RAStream( )) do
    T(r & (m-1)) \wedge= r;
```

## RA in Chapel: Multi-Locale

Given:  $m$ -element table  $T$  (where  $m = 2^n$ )

Compute: **forall**  $r$  in RandomUpdates **do**

**T**( $r \& (m-1)$ )  $\hat{=}$   $r$ ;

```
config const m = ..., N_U = ..., tpl = ...;
const TableDist = new Block(1,uint(64),[0..m-1],tpl),
      UpdateDist = new Block(1,uint(64),[0..N_U-1],tpl),
      TableSpace: domain(1,uint(64))
                  distributed TableDist = [0..m-1],
      Updates: domain(1,uint(64))
                  distributed UpdateDist = [0..N_U-1];
var T: [TableSpace] uint(64);

forall (i,r) in (Updates,RAStream()) do
  on T( $r \& (m-1)$ ) do
    T( $r \& (m-1)$ )  $\hat{=}$   $r$ ;
```

## RA in Chapel: Multi-Locale

Given:  $m$ -element table  $T$  (where  $m = 2^n$ )

Compute: **forall**  $r$  in RandomUpdates **do**

**T**( $r \& (m-1)$ )  $\hat{=}$   $r$ ;

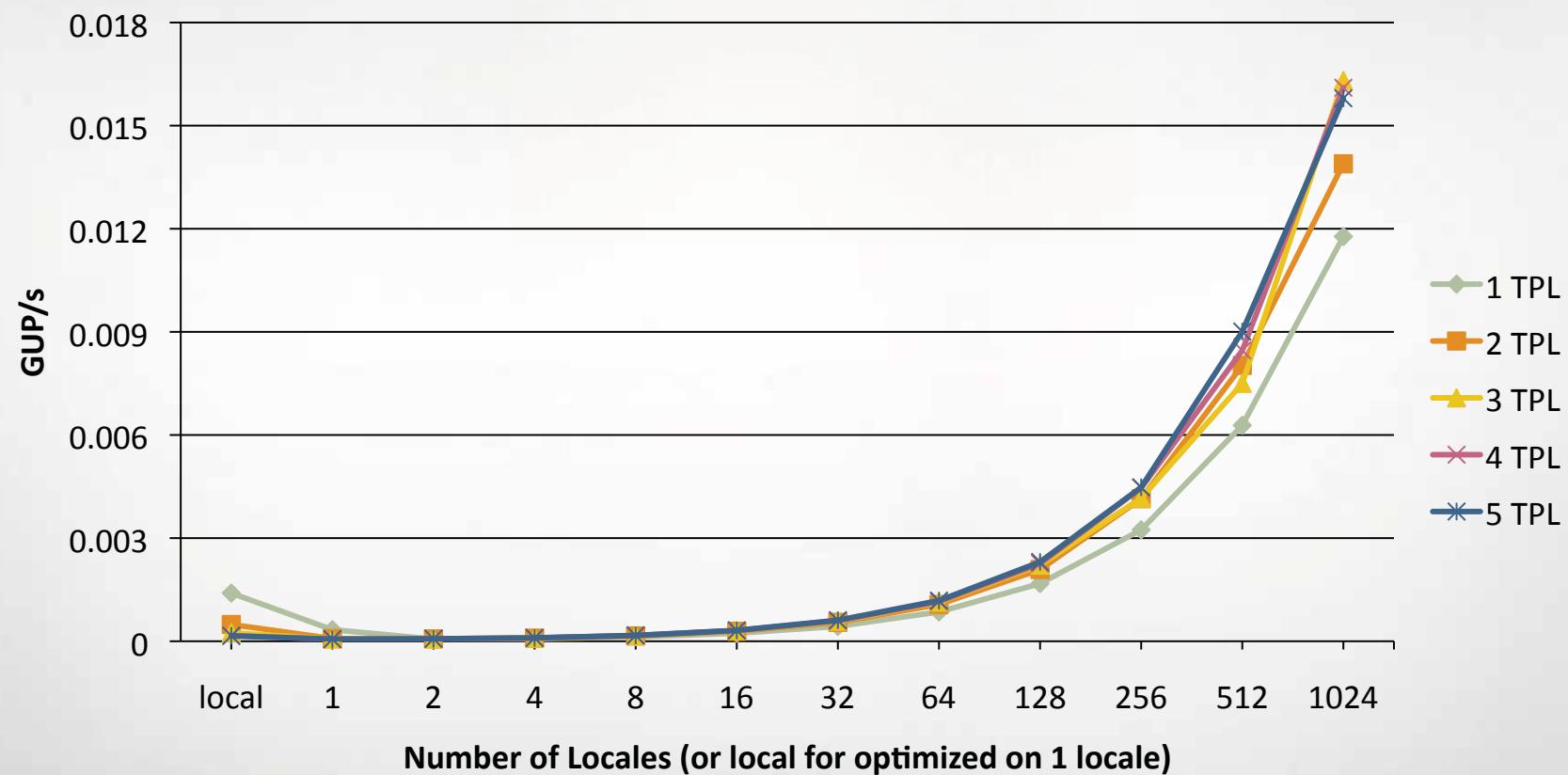
```
config const m = ..., N_U = ..., tpl = ...;
const TableDist = new Block(1,uint(64),[0..m-1],tpl),
      UpdateDist = new Block(1,uint(64),[0..N_U-1],tpl),
      TableSpace: domain(1,uint(64))
                  distributed TableDist = [0..m-1],
      Updates: domain(1,uint(64))
                  distributed UpdateDist = [0..N_U-1];
var T: [TableSpace] uint(64);
forall (i,r) in (Updates,RAStream()) do
  on T.domain.dist.ind2loc( $r \& (m-1)$ ) do
    T( $r \& (m-1)$ )  $\hat{=}$   $r$ ;
```

Call ind2loc method directly

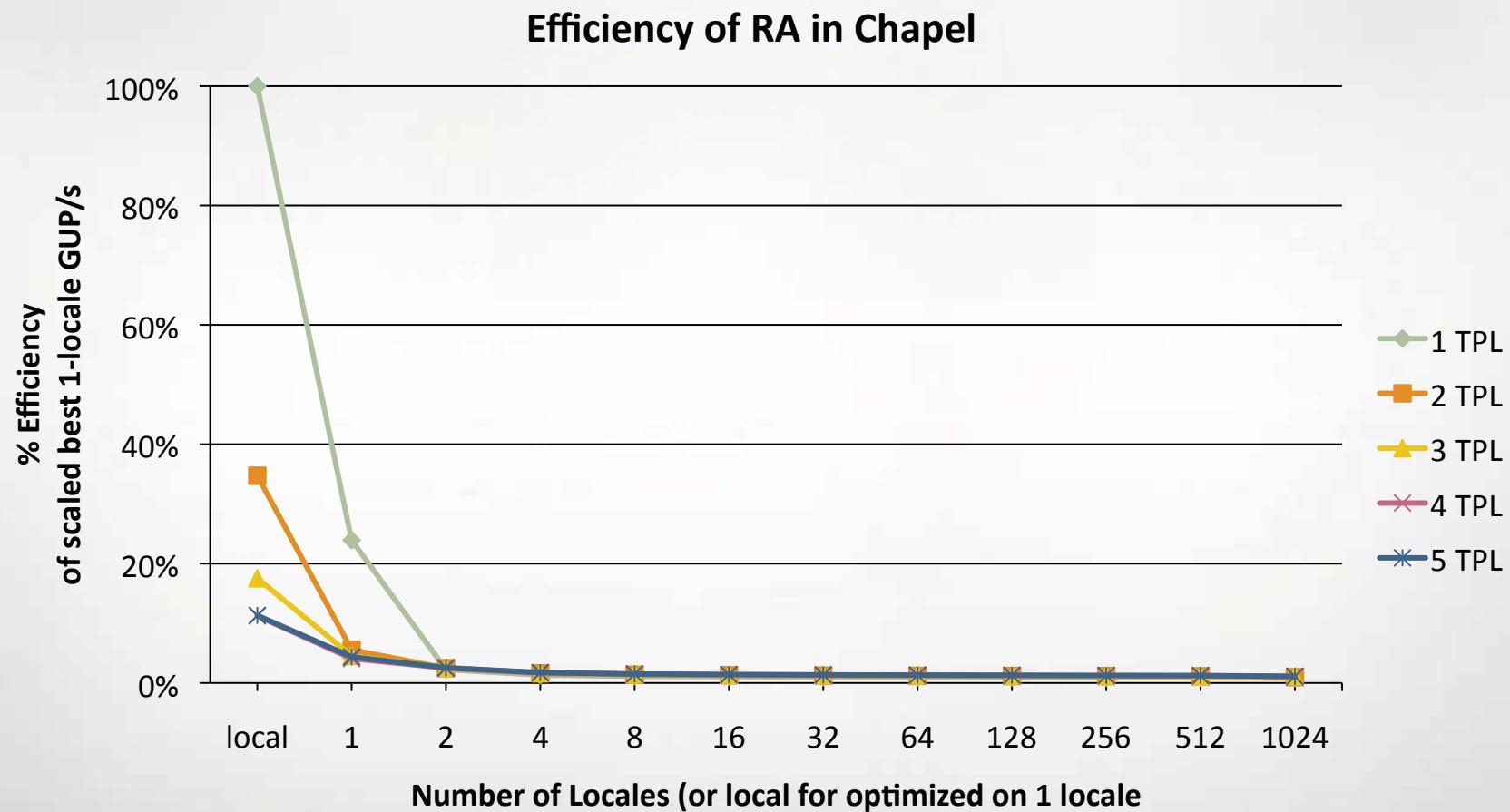
# HPCC RA Performance

CRAY  
THE SUPERCOMPUTER COMPANY

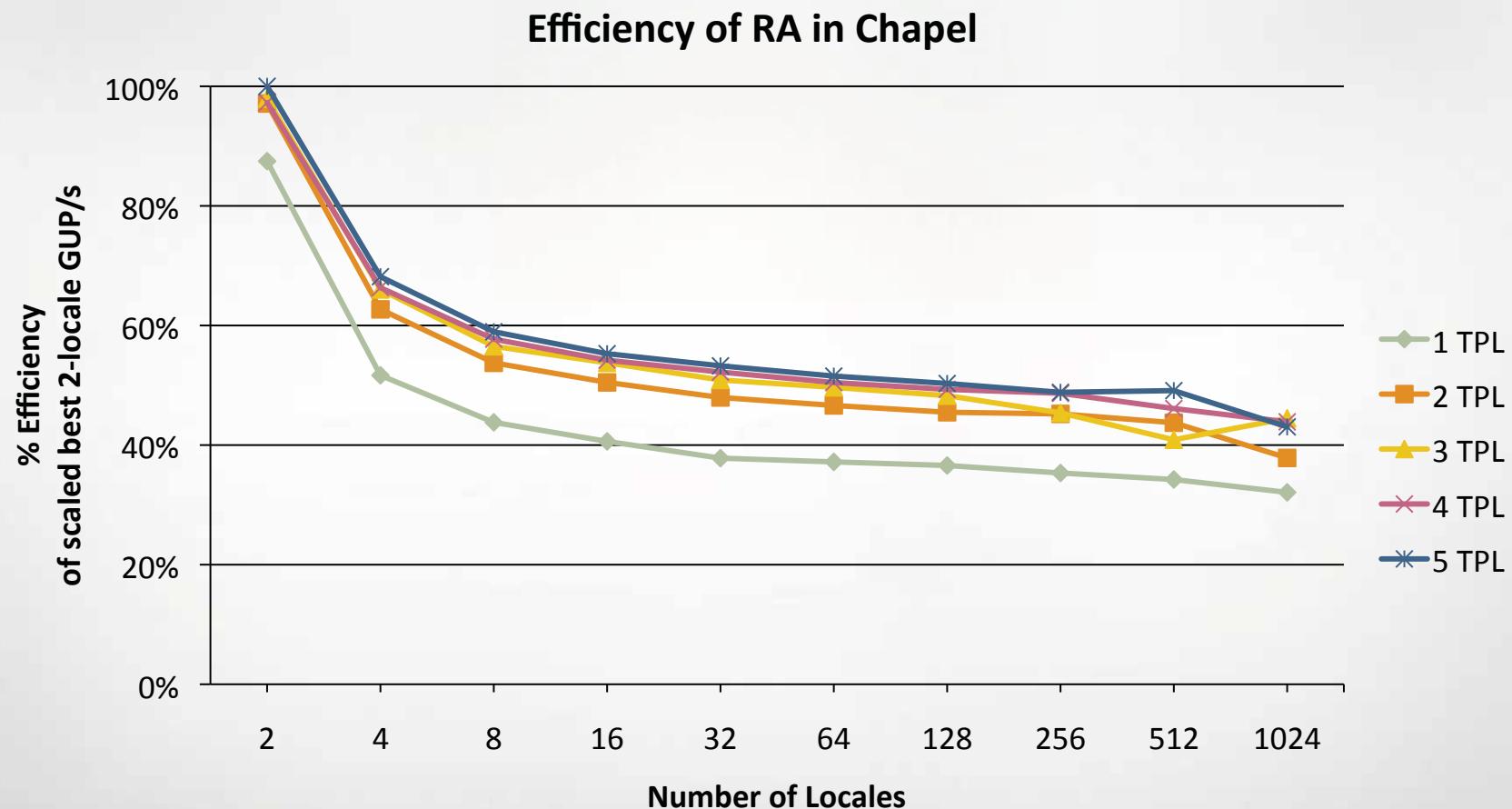
Performance of RA in Chapel



# HPCC RA Efficiency I



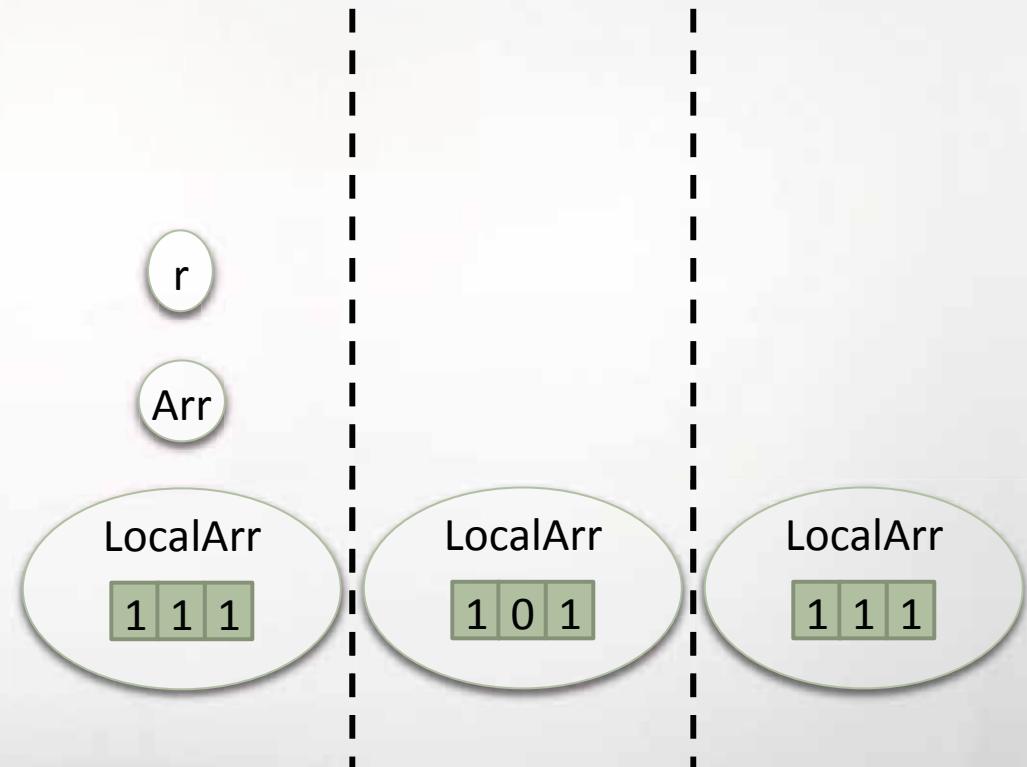
# HPCC RA Efficiency II



# Optimization: Remote Value Forwarding

## Simple example

```
var Arr: [Dom] int;  
var r: int;  
on Locales(1) {  
    Arr(r) ^= r;  
}
```



## Outline

- What is Chapel?
- Chapel's Parallel Programming Model
- HPCC STREAM Triad in Chapel
- HPCC RA in Chapel
- Summary and Future Work

## Summary

The global-view programming model is easy to use.

- Shorter, more concise code
- Separation of concerns (partitioning)
- Easy to change data distributions

Distributions implement the global-view model.

- Flexible mechanism for experimentation
- Implementation of distributions is in Chapel

## Future Work

- Optimizations
  - Within the compiler
  - Within the runtime
  - Within the distributions
- Complete implementation of Block distribution
- Implement new distributions
  - Cyclic, BlockCyclic, RecursiveBisection
- Experiment with variations of STREAM and RA

# Questions?