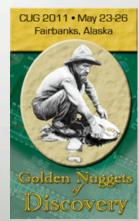


Authoring User-Defined Domain Maps in Chapel

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Cray Inc.

CUG 2011: May 24th, 2011



What is Chapel?



- A new parallel programming language
 - Design and development led by Cray Inc.
 - Started under the DARPA HPCS program
- Overall goal: Improve programmer productivity
 - Improve the programmability of parallel computers
 - Match or beat the performance of current programming models
 - Support better portability than current programming models
 - Improve the robustness of parallel codes
- A work-in-progress



Chapel's Implementation

- Being developed as open source at SourceForge
- Licensed as BSD software
- Target Architectures:
 - multicore desktops and laptops
 - commodity clusters
 - Cray architectures
 - systems from other vendors
 - (in-progress: CPU+accelerator hybrids)



Chapel's High-Level Themes

General Parallel Programming

• "any parallel algorithm on any parallel hardware"

Multiresolution Parallel Programming

- high-level features for convenience/simplicity
- low-level features for greater control

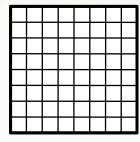
Control over Locality/Affinity of Data and Tasks

for scalability



Sample Computation: Sum-of-Squares

```
config const n = computeProblemSize();
const D = [1..n, 1..n];
```



var A, B: [D] real;

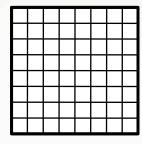
sumOfSquares

const sumOfSquares = + reduce (A**2 + B**2);



Sample Computation: Sum-of-Squares

```
config const n = computeProblemSize();
const D = [1..n, 1..n];
```



var A, B: [D] real;

sumOfSquares

D

**2

A

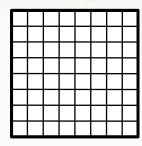
B

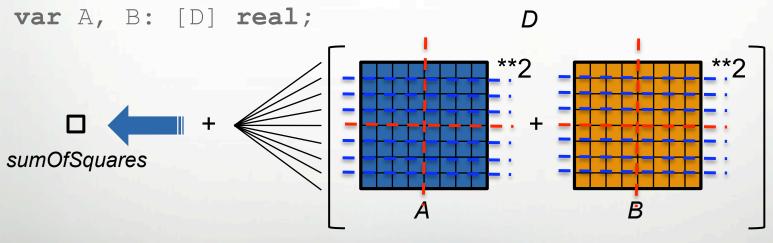
const sumOfSquares = + reduce (A**2 + B**2);



Sample Computation: Sum-of-Squares

```
config const n = computeProblemSize();
const D = [1..n, 1..n] dmapped ...;
```





const sumOfSquares = + reduce (A**2 + B**2);



Sum-of-Squares Implementation

```
config const n = computeProblemSize();
const D = [1..n, 1..n];
var A, B: [D] real;

const sumOfSquares = + reduce (A**2 + B**2);
```

How is this global-view computation implemented in practice?

ZPL: Block-distributed arrays, serial on-node computation (inflexible)

HPF: Not particularly well-defined ("trust the compiler")

Chapel: Very flexible and well-defined via domain maps (stay tuned)

Outline



- ✓ Background and Motivation
- Chapel Background:
 - Locales
 - Domains, Arrays, and Domain Maps
- Implementing Domain Maps
- Wrap-up

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The Locale Type

Definition

- Abstract unit of target architecture
- Supports reasoning about locality
- Capable of running tasks and storing variables
 - i.e., has processors and memory

Properties

- a locale's tasks have ~uniform access to local vars
- Other locale's vars are accessible, but at a price

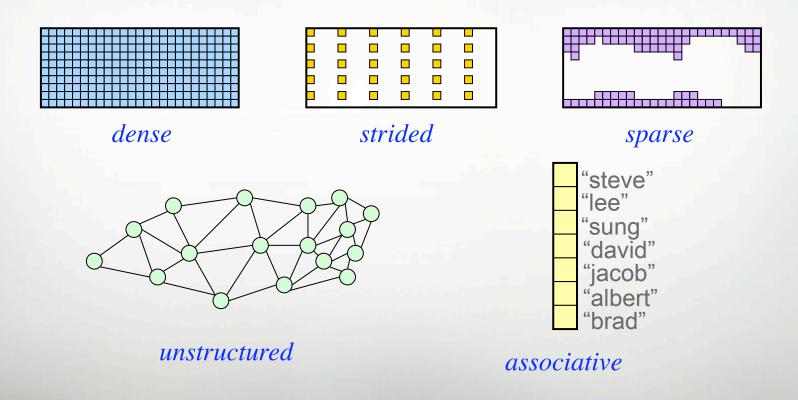
Locale Examples

- A multi-core processor
- An SMP node



Chapel Domain/Array Types

Chapel supports several types of domains and arrays:





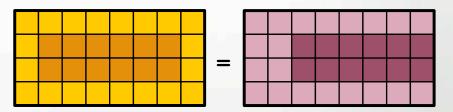
Chapel Domain/Array Operations

Whole-Array Operations; Parallel and Serial Iteration

A = forall (i,j) in D do (i +
$$j/10.0$$
);

| 1.1 | 1.2 | 1.3 | 1.4 | 1.5 | 1.6 | 1.7 | 1.8 |
|-----|-----|-----|-----|-----|-----|-----|-----|
| 2.1 | 2.2 | 2.3 | 2.4 | 2.5 | 2.6 | 2.7 | 2.8 |
| 3.1 | 3.2 | 3.3 | 3.4 | 3.5 | 3.6 | 3.7 | 3.8 |
| 4.1 | 4.2 | 4.3 | 4.4 | 4.5 | 4.6 | 4.7 | 4.8 |

Array Slicing; Domain Algebra



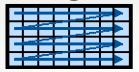
 And several other operations: indexing, reallocation, domain set operations, scalar function promotion, ...



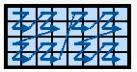
Data Parallelism: Implementation Qs

Q1: How are arrays laid out in memory?

Are regular arrays laid out in row- or column-major order? Or...?







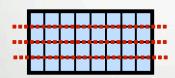


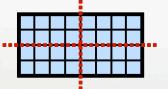
...?

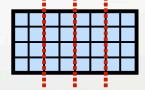
What data structure is used to store sparse arrays? (COO, CSR, ...?)

Q2: How are data parallel operators implemented?

- How many tasks?
- How is the iteration space divided between the tasks?









...?



Data Parallelism: Implementation Qs

Q3: How are arrays distributed between locales?

- Completely local to one locale? Or distributed?
- If distributed... In a blocked manner? cyclically? block-cyclically? recursively bisected? dynamically rebalanced? ...?

Q4: What architectural features will be used?

- Can/Will the computation be executed using CPUs? GPUs? both?
- What memory type(s) is the array stored in? CPU? GPU? texture? ...?

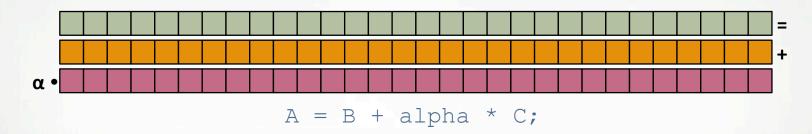
A1: In Chapel, any of these could be the correct answer

A2: Chapel's *domain maps* are designed to give the user full control over such decisions

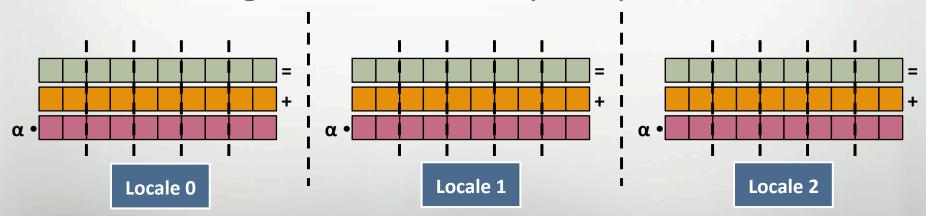




Domain maps are "recipes" that instruct the compiler how to map the global view of a computation...



...to the target locales' memory and processors:







Domain Maps: "recipes for implementing parallel/ distributed arrays and domains"

They define data storage:

- Mapping of domain indices and array elements to locales
- Layout of arrays and index sets in each locale's memory

...as well as operations:

- random access, iteration, slicing, reindexing, rank change, ...
- the Chapel compiler generates calls to these methods to implement the user's array operations



Domain Maps: Layouts and Distributions

Domain Maps fall into two major categories:

layouts: target a single locale

- (that is, a desktop machine or multicore node)
- examples: row- and column-major order, tilings, compressed sparse row

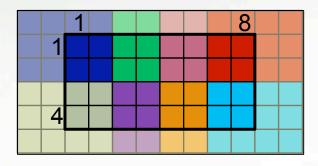
distributions: target distinct locales

- (that is a distributed memory cluster or supercomputer)
- examples: Block, Cyclic, Block-Cyclic, Recursive Bisection, ...



Sample Distributions: Block and Cyclic

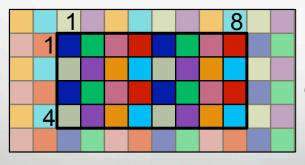
var Dom = [1..4, 1..8] dmapped Block([1..4, 1..8]);



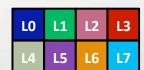
distributed to



var Dom = [1..4, 1..8] dmapped Cyclic(startIdx=(1,1));

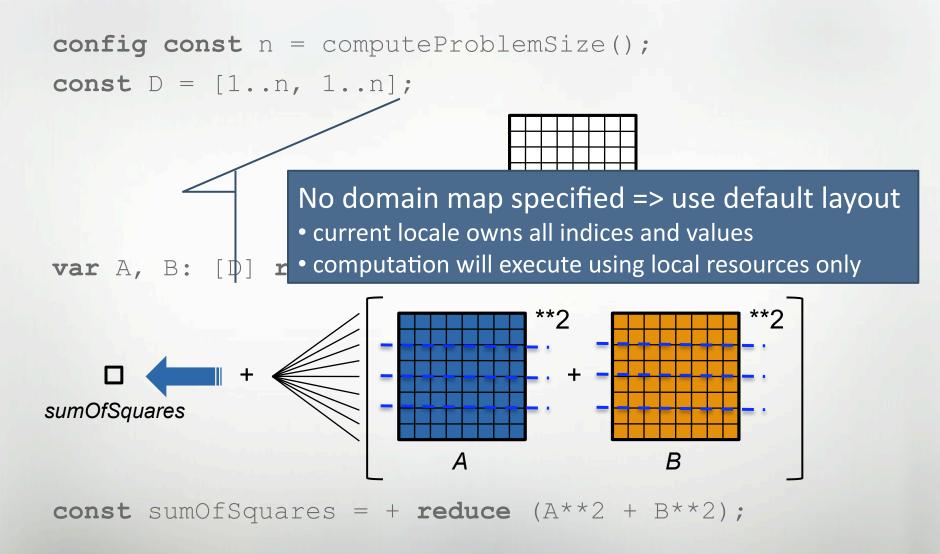


distributed to



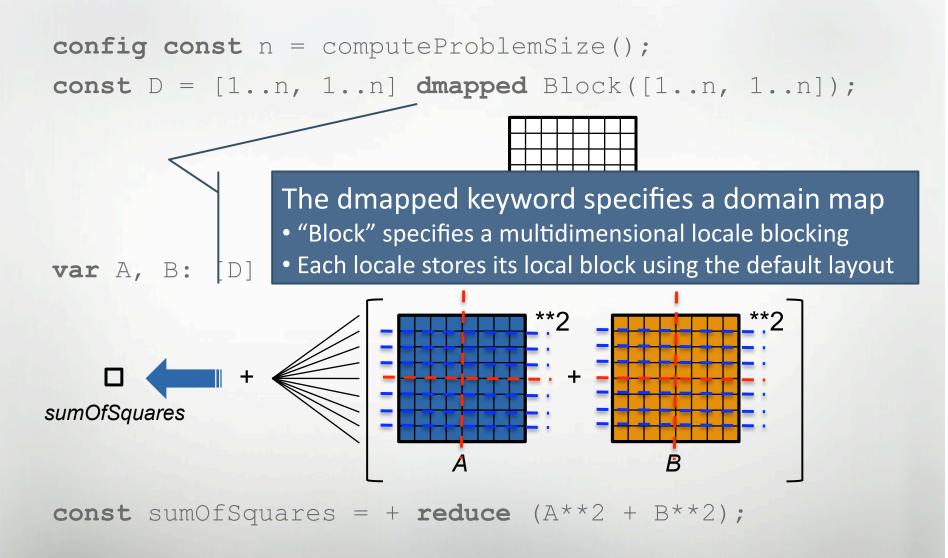


Sample Computation: Local Sum-of-Squares



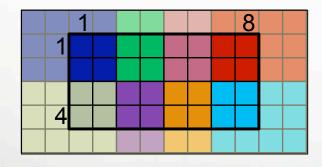


Sample Computation: Distributed Sum-of-Squares





The Complete Block class constructor

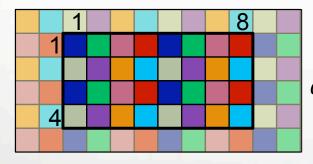


distributed to

| LO | L1 | L2 | L3 |
|----|----|----|-----------|
| L4 | L5 | L6 | L7 |



The Cyclic class constructor



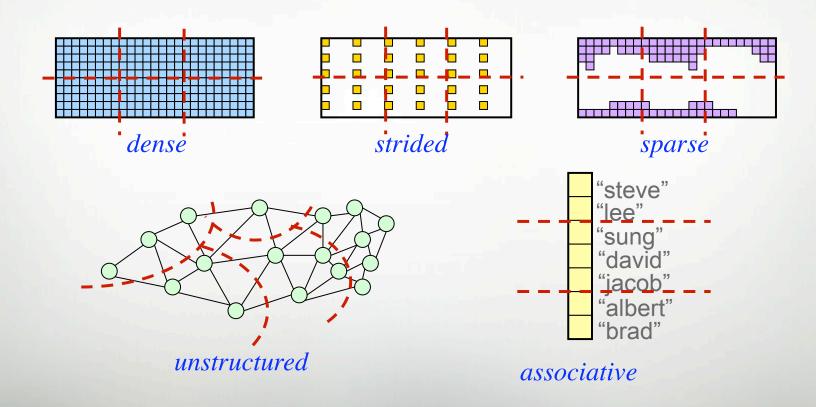
distributed to

| LO | L1 | L2 | L3 |
|----|----|----|----|
| L4 | L5 | L6 | L7 |



All Domain Types Support Domain Maps

All Chapel domain types support domain maps







- ✓ Background and Motivation
- ✓ Domains, Arrays, and Domain Maps
- Implementing Domain Maps
 - Philosophy
 - Implementing Layouts
 - Implementing Distributions
- Wrap-up



Chapel's Domain Map Philosophy

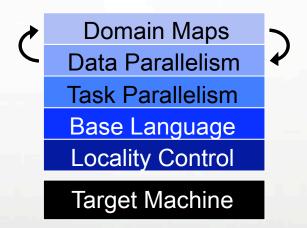
- 1. Chapel provides a library of standard domain maps
 - to support common array implementations effortlessly
- 2. Advanced users can write their own domain maps in Chapel
 - to cope with shortcomings in our standard library
- 3. Chapel's standard layouts and distributions will be written using the same user-defined domain map framework
 - to avoid a performance cliff between "built-in" and user-defined domain maps
- 4. Domain maps should only affect implementation and performance, not semantics
 - to support switching between domain maps effortlessly



Chapel's Multiresolution Design

Multiresolution Design: Support multiple tiers of features

- higher levels for programmability, productivity
- lower levels for greater degrees of control
- build the higher-level concepts in terms of the lower
 Chapel language concepts



- separate concerns appropriately for clean design
 - yet permit the user to intermix the layers arbitrarily



Domain Maps and Multiresolution

- Domain Maps are implemented using Chapel
- They are considered Chapel's highest-level feature
- As such they are implemented using lower-level Chapel concepts:
 - base language: classes, iterators, type inference, generic types to organize and simplify code
 - task parallelism: to implement parallel operations
 - locality control: locales and on-clauses to map to hardware
 - data parallelism: other domains and arrays for local storage







Domain Map

Represents: a domain

map value

Generic w.r.t.: index type

State: the domain map's representation

Typical Size: Θ(1)

Domain

Represents: a domain

Generic w.r.t.: index type

State: representation of index set

Typical Size: $\Theta(1) \rightarrow \Theta(numIndices)$

Array

Represents: an array

Generic w.r.t.: index type, element type

State: array elements

Typical Size: Θ(numIndices)



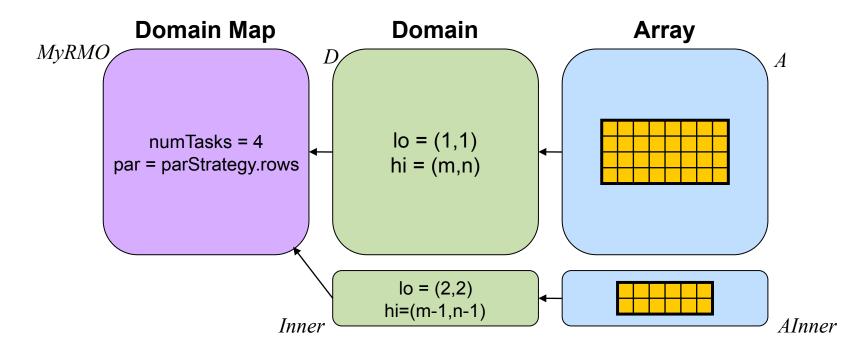
Chapel Declarations and Resulting Descriptors

```
const myDomMap = new dmap(DomMapName(args));
const D1 = [1..10] dmapped MyDomMap,
      D2 = [1..20] dmapped MyDomMap;
var A1, B1: [D1] real,
                                                    A1
   A2, B2: [D2] string,
   C2: [D2] complex;
                                                    B1
                               D1
                                                    A2
        myDomMap
                                                    B2
                               D2
```



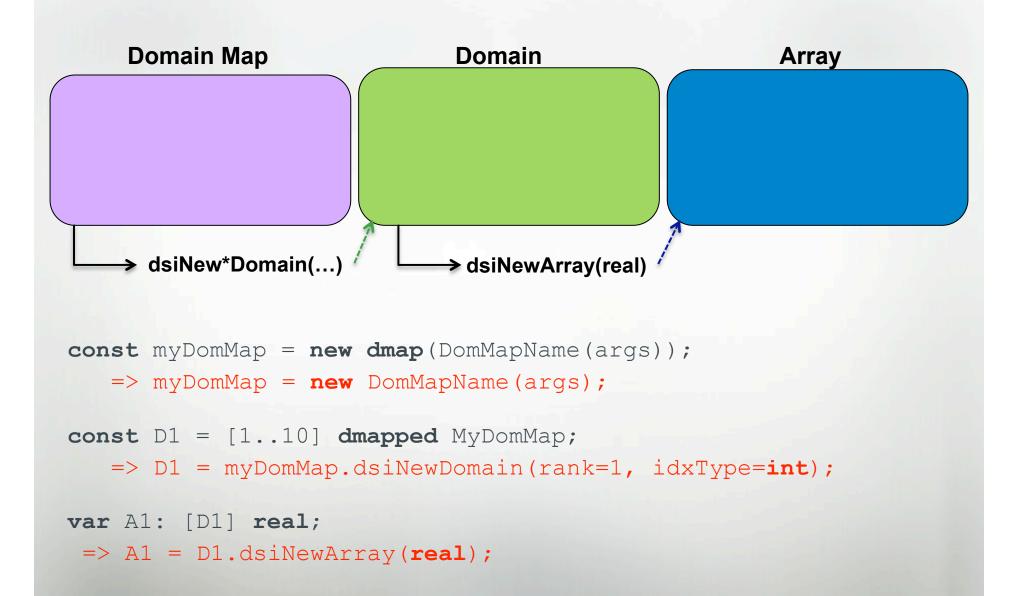
Sample Layout Descriptors

AInner: [Inner] real;

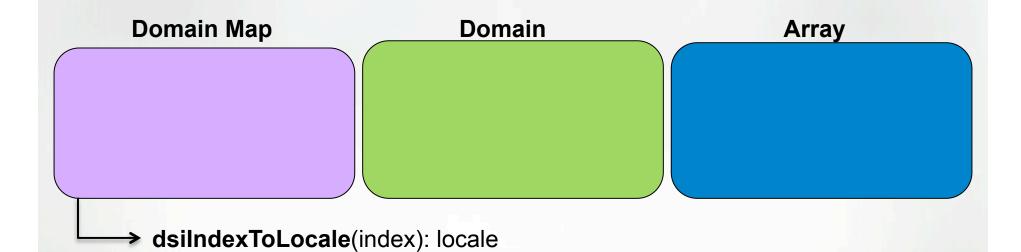






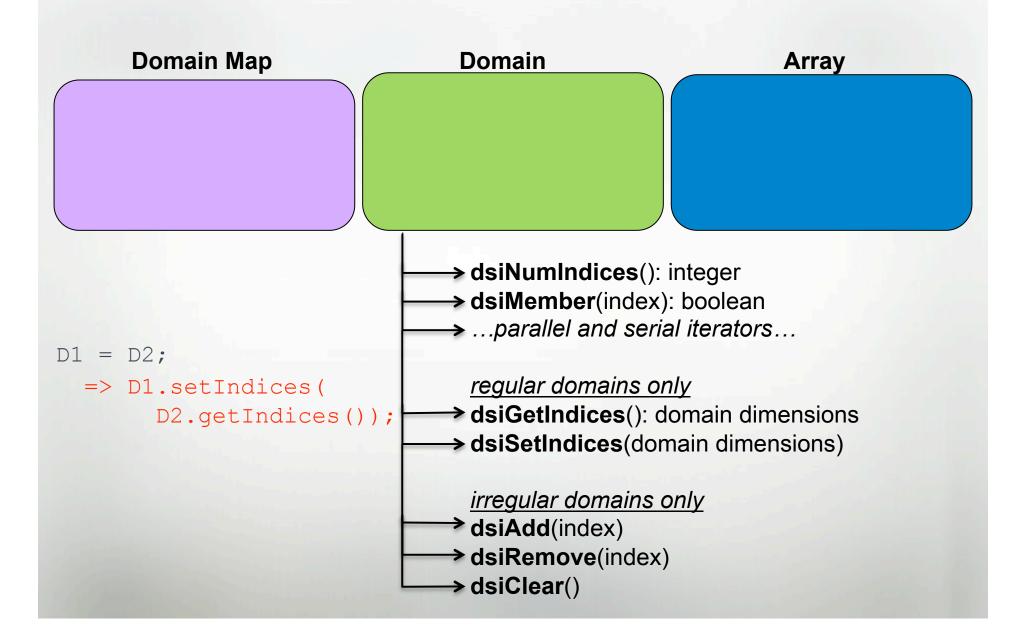




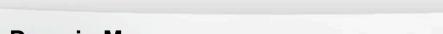


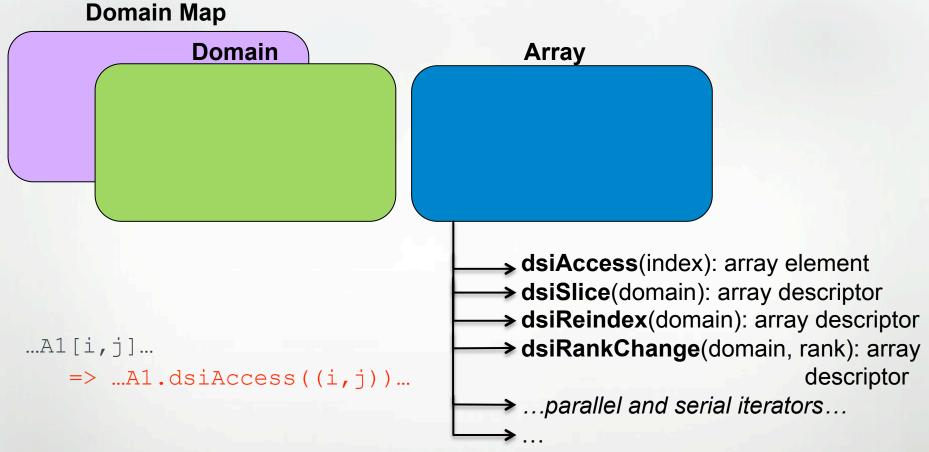
```
...myDomMap.indexToLocale((i,j))...
=> myDomMap.indexToLocale((i,j))
```













Distribution Descriptors (One Approach)

Domain Map

Role: Similar to layout's domain map descriptor

Size: $\Theta(1) \rightarrow$ Θ(#locales)

Role: Similar to layout's domain descriptor, but no Θ(#indices) storage

Size: $\Theta(1) \rightarrow$ Θ(#locales)

Domain **Array**

Role: Similar to layout's array descriptor, but data is moved to local descriptors

Size: $\Theta(1) \rightarrow$ Θ(#locales)

Local

Global

one instance

per object

(logically)

one instance per locale per object (typically)

Role: Stores localespecific domain map parameters

Size: $\Theta(???)$

Role: Stores locale's subset of domain's index set

Size: $\Theta(1) \rightarrow$ Θ(#indices / #locales)

Role: Stores locale's subset of array's elements

Size:

Θ(#indices / #locales)

Compiler only knows about global descriptors so local are just a specific type of state; interface is identical to layouts



Sample Distribution Descriptors

Global

one instance per object (logically)

Domain Map

boundingBox =
[1..4, 1..8]

targetLocales =

L0 L1 L2 L3

L4 L5 L6 L7

Domain

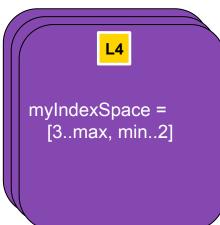
indexSet = [1..4, 1..8]

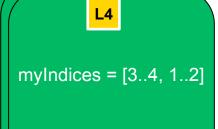
Array

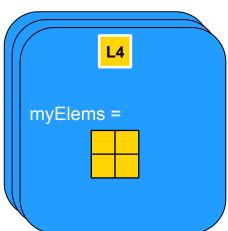
__

Local

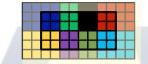
one instance per node per object (typically)







var Dom= [1..4, 1..8] dmapped Block(boundingBox=[1..4, 1..8]);









Sample Distribution Descriptors

Global

one instance per object (logically)

Domain Map

boundingBox =
[1..4, 1..8]

targetLocales =

L0 L1 L2 L3

L4 L5 L6 L7

Domain

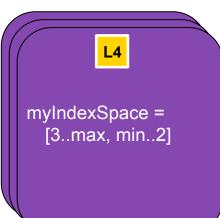
indexSet = [2..3, 2..7]

Array

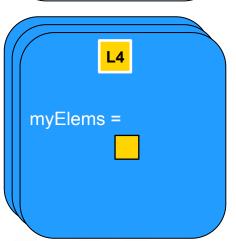
-

Local

one instance per node per object (typically)







```
var Dom= [1..4, 1..8] dmapped Block(boundingBox=[1..4, 1..8]);
```

var Inner =
$$Dom[2..3, 2..7]$$
;







Non-Required Descriptor Interfaces

Optional Interfaces

- Do not need to be supplied for correctness
- But supplying them may permit optimizations
- Examples:
 - privatization of global descriptors
 - communication optimizations: stencils, reductions/broadcasts, remaps

User Interfaces

- Add new user methods to domains, arrays
- Not known to the compiler
- Break plug-and-play nature of distributions





- ✓ Background and Motivation
- ✓ Domains, Arrays, and Domain Maps
- ✓ Implementing Domain Maps
- Wrap-up



Domain Maps: Status

- All Chapel domains and arrays implemented using this framework
 - Full-featured Block, Cyclic, and Replicated distributions
 - COO and CSR Sparse layouts
 - Open addressing quadratic probing Associative layout
 - Block-Cyclic, Dimensional, and Distributed Associative distributions underway
- Initial performance/scaling results promising, but more work remains
- Adding documentation for authoring domain maps

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Future Directions

- More advanced uses of domain maps:
 - CPU+GPU cluster programming
 - Dynamic load balancing
 - Resilient computation
 - in situ interoperability
 - Out-of-core computations



Summary

- Chapel's domain maps are a promising language concept
 - permit better control over -- and ability to reason about -parallel array semantics than in previous languages
 - separate specification of an algorithm from its implementation details
 - support a separation of roles:
 - parallel expert writes domain maps
 - parallel-aware computational scientist uses them



For More Information on Domain Maps

- HotPAR'10 paper: User-Defined Distributions and Layouts in Chapel: Philosophy and Framework
- This CUG'11 paper
- In the Chapel release...
 - Technical notes detailing the domain map interface for programmers:
 \$CHPL_HOME/doc/technotes/README.dsi
 - Browse current domain maps:

```
$CHPL_HOME/modules/dists/*.chpl
layouts/*.chpl
internal/Default*.chpl
```



For More Information on Chapel

- Chapel Home Page (papers, presentations, tutorials):
 http://chapel.cray.com
- Chapel Project Page (releases, source, mailing lists):
 http://sourceforge.net/projects/chapel/
- General Questions/Info: <u>chapel info@cray.com</u> (or chapel-users mailing list)

Our Team



• Cray:













Brad Chamberlain

Sung-Eun Choi

Greg Titus

Lee Prokowich

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External Collaborators:











Albert Sidelnik

Jonathan Turner

Srinivas Sridharan

You? Your Student?

• Interns:













Jonathan Claridge Hannah Hemmaplardh

Andy Stone

Jim Dinan

Rob Bocchino

Mack Joyner



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http://sourceforge.net/projects/chapel/