



Production I/O Characterization on the Cray XE6

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Motivation

- I/O behavior plays a key role in productivity for largescale HPC systems
 - Our target: Hopper, a 153,216 core Cray XE6 at NERSC
- Challenges in understanding I/O behavior
 - Hundreds of users across a wide spectrum of science
 - Applications vary in data volume, I/O strategy, and access methods
 - How can we consistently characterize production I/O behavior across applications?
 - How do we quickly identify applications that could most benefit from additional tuning assistance?



Approach

- 1. Adapt the Darshan I/O characterization tool for use in the Cray environment
 - Tune to reflect Hopper system characteristics
 - Integrate transparently for maximum coverage
- 2. Evaluate Darshan for production deployment
 - Measure overhead at scale for multiple workloads
- 3. Deploy Darshan
 - Store characterization data for post-processing and exploration
 - Provide immediate feedback to users
- 4. Develop tools that leverage Darshan data
 - Rapid feedback to power users
 - Metrics to automatically flag jobs that exhibit unusual I/O behavior for administrators

Darshan background

Darshan is an open source, application-level instrumentation library that uses link-time instrumentation for static executables and LD_PRELOAD for dynamic executables.

What does it record?

- Counters, histograms, and strategically chosen timestamps related to I/O activity
- POSIX, POSIX stream, MPI-IO, and limited HDF5 and PNetCDF functions
- Access patterns, access sizes, I/O time, alignment, datatypes, etc.
- Builds on characterization ideas from Kotz and Nieuwejaar Charisma study
- Does not record a complete trace of I/O operations

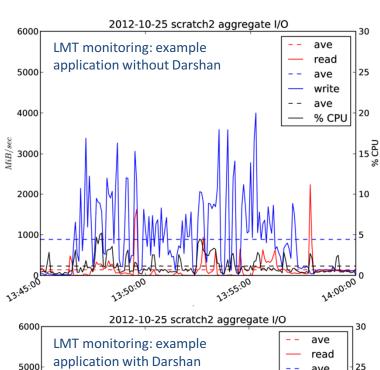
How does it store results?

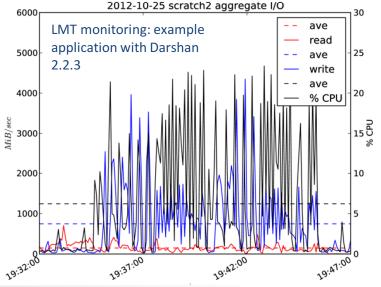
- Minimal overhead during execution
- Reduction, compression, and persistent storage at MPI_Finalize() time
- Produces a single, compact log for each instrumented job



Tuning example: shared-file instrumentation

- Initial performance experiments with Darshan 2.2.3 exhibited performance degradation for shared-file applications
- LMT monitoring of servers revealed that the raw I/O throughput was comparable, but CPU usage was much higher
- Source of overhead: issuing stat() calls on each rank to collect additional information at open() time
- Concurrent stat() at 12,288 processes can already add 2 seconds to file open time
- Worked to develop alternative instrumentation



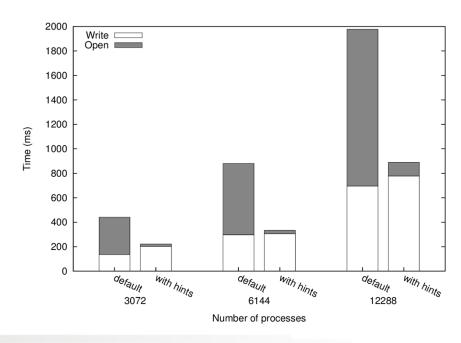






Tuning example: writing log files efficiently

- Darshan writes all results to a unified log file at shutdown after custom reduction and gzip compression
- Final results are typically quite small: hundreds of bytes per process, sometimes even less
- Regardless of how the application performed I/O, Darshan itself uses
 MPI-IO collectives internally to write results
- Portable and efficient: leverages aggregation and stripe alignment
- We improved MPI-IO collective performance even further with hints:
 - romio_no_indep_rw
 - cb_nodes=4
- Limit the number of processes that actually open the output file
- Drastically reduces the cost of writing data at larger process counts





System integration challenges

- Cray environment:
 - Multiple compilers (Cray, PGI, Pathscale, Intel, GNU)
 - Static linking by default
 - Unified cc, ftn, and CC compiler scripts
- Our requirements:
 - Support as many configurations as reasonably possible
 - Enable and disable via software module
 - Transparent for users (no need for different compilers or link options)
- We experimented with multiple deployment methods during this study
- Our plan moving forward is to use the PE_PRODUCT_LIST mechanism
 - This is a set of environment variables that can be used by software modules to specify additional linker options for the Cray compiler scripts



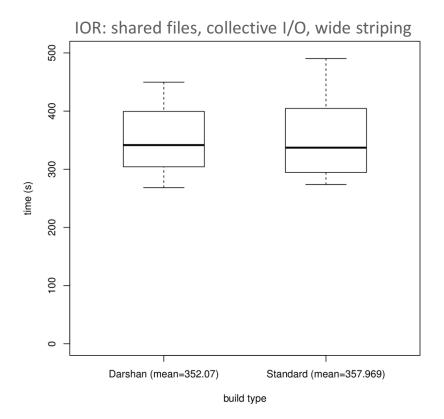
Evaluation: end-to-end overhead

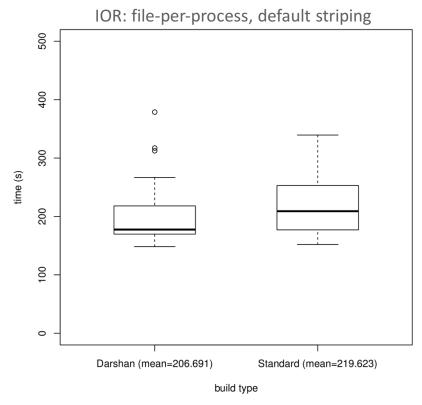
After adapting Darshan to the Cray XE6 environment, our next goal was to evaluate the impact of Darshan on large-scale applications to verify suitability for production deployment.

- First experiment: measure end-to-end run time of IOR benchmark writing and reading 1.5 TiB of data at 12,288 processes
 - Includes all Darshan startup, instrumentation, and shutdown costs
 - Measured by timing the "aprun" command
 - Evaluate both shared-file and file-per-process examples
 - IOR configured to use MPI-IO; Darshan instrumented both the MPI-IO and POSIX API calls
- Individual runs are susceptible to variance
- Gathered 20 independent samples for each test case over a four day period and analyzed the results



Evaluation: end-to-end overhead





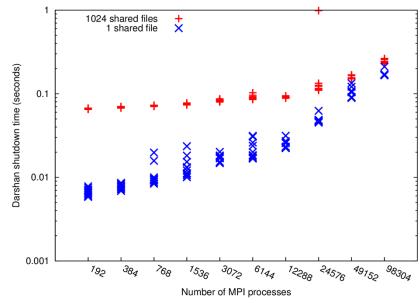
- Box plots of 20 samples each indicate no clear Darshan overhead at 12,288 processes relative to normal system variance
- Insufficient evidence for difference in mean run time (t-test)
- Other observations: variance is significant, and file-per-process access patterns perform better on Hopper at this scale

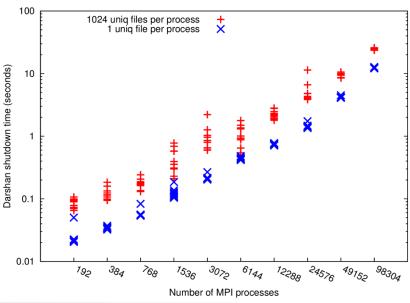
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Evaluation: Darshan shutdown costs

- End-to-end Darshan overhead is obscured by I/O variance, but we can measure internal Darshan mechanisms directly using microbenchmarks
- These examples show the total time required by Darshan to reduce, compress, and write log files at MPI_Finalize() time
- Average one-time cost for 98,304 processes:
 - 1 shared file: 0.17 seconds
 - 1024 shared files: 0.24 seconds
 - 1 file-per-process: 12.3 seconds
 - 1024 files-per-process: 24.8 seconds
- Largest scenario: emulates application opening 100,663,296 unique files
- Darshan falls back to less granular instrumentation if memory threshold is exceeded by opening too many files



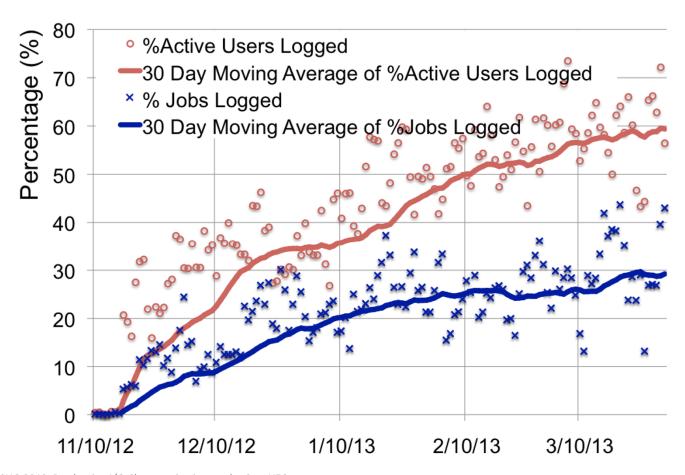






Deployment: coverage as of March 2013

- Percentage of active users and jobs instrumented per day since initial Hopper deployment
- Nearing an average of 60% and 30%, respectively, by end of March 2013







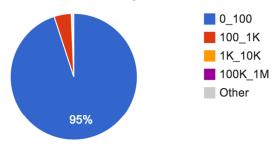
Deployment: user experience

- The NERSC web portal allows users to interact with their jobs and allocations
- This screenshot shows I/O summary information automatically generated from Darshan logs for completed jobs
- Provides rapid (within a few minutes) initial feedback on I/O behavior

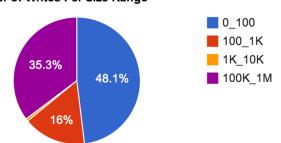
IO Summary from Darshan

Start	End	Wallclock (secs)		MB Written		Estimated Percent Time Spent in I/O
04-05 16:04:05	04-05 16:06:51	166	590.3	597.6	355.52	2.01%





Number of Writes Per Size Range





Metrics: redundant read traffic

- Administrators can also filter logs using metrics designed to automatically identify applications that may benefit from tuning assistance
- We explored three example metrics that can be quickly computed from Darshan log data
- First example is redundant read traffic: applications that read more bytes of data from the file system than were present in the file
- Even with caching effects, this type of job can cause disruptive I/O network traffic through redundant file system transfers
- Candidates for aggregation or collective I/O

Summary of matching jobs:

Redundant read threshold	> 1 TiB
Total jobs analyzed	261,890
Jobs matching heuristic	671
Unique users matching heuristic	37
Largest single-job redundant read volume	547 TiB

Top example

Scale: 6,138 processes

Run time: 6.5 hours

Avg. I/O time per process: 27 minutes

Metric: Read 548.6 TiB of data from a 1.2
 TiB collection of read-only files

 Used 8 KiB read operations and generated 457 X redundant read traffic



Metrics: metadata time ratio

- Percentage of cumulative I/O time spent performing metadata operations such as open(), close(), stat(), and seek()
- Close() cost can be misleading due to write-behind cache flushing, but metadata ratio is often a key indicator of inefficient file organization
- Most relevant for jobs that performed a significant amount of I/O
- Candidates for coalescing files and eliminating extra metadata calls like stat() where possible

Summary of matching jobs:

Thresholds	meta_time / nprocs > 30 s
	$nprocs \ge 192$
	metadata_ratio $\geq 25\%$
Total jobs analyzed	261,890
Jobs matching heuristic	252
Unique users matching heuristic	45
Largest single-job metadata ratio	> 99%

Top example

Scale: 40,960 processes

Run time: 229 seconds

Max. I/O time per process: 103 seconds

Metric: 99% of I/O time in metadata operations

 Generated 200,000+ files using 600,000+ write() and 600,000+ stat() calls



Metrics: small writes to shared files

- Small writes can contribute to poor performance as a result of poor file system stripe alignment, but there are many factors to consider:
 - Writes to non-shared files may be cached aggressively
 - Collective writes are normally optimized by MPI-IO
 - Throughput can be influenced by additional factors beyond write size
- We searched for jobs that wrote less than 1 MiB per operation to shared files without using any collective operations
- Candidates for collective I/O or batching/buffering of write operations

Summary of matching jobs:

Thresholds	> 100 million small writes
	0 collective writes
Total jobs analyzed	261,890
Jobs matching heuristic	220
Unique users matching heuristic	11
Largest single-job small write count	5.7 billion

Top example

Scale: 128 processes

Run time: 30 minutes

Max. I/O time per process: 12 minutes

- Metric: issued 5.7 billion write operations, each less than 100 bytes in size, to shared files
- Averaged just over 1 MiB/s per process during shared write phase



Summary and conclusions

Lessons learned while adapting Darshan for the Cray environment:

- Adapting system tools to the Cray XE6 environment requires consideration of file system characteristics and environment configuration
- "Minimal overhead" is a moving target! New platforms bring new challenges.
- I/O variance remains a significant factor in I/O performance
- Darshan performance is suitable for full-time production deployment

Deployment experience:

- It is possible to perform application-level I/O characterization of full-scale production workloads
- Darshan has been successfully deployed on the Hopper system at NERSC
 - Rolled out in stages from November 2012 to February 2013
 - 30% of jobs and compute hours instrumented by March 2013
- Deployment impact:
 - Immediate feedback to users on I/O behavior
 - Ability to automatically scan jobs using simple metrics to identify applications that need further assistance



Future work

- Continue to improve Darshan
 - Use system-specific mechanisms for characterization when available
 - Improve Cray environment integration
 - Modularization
- Continue to use Darshan data to solve I/O problems
 - Use metrics to identify and help the users who need it most
 - More rigorous analysis methods
 - Improve feedback for users
 - Broader workload studies



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