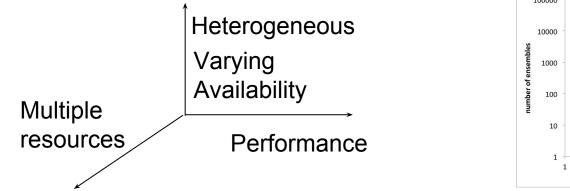
Executing dynamic heterogeneous workloads on Blue Waters with RADICAL-Pilot

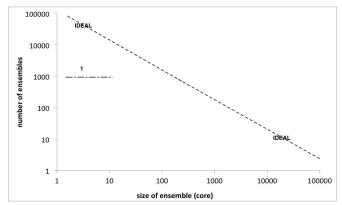
Mark Santcroos^{*}, Ralph Castain[†], Andre Merzky^{*}, Iain Bethune[‡] and Shantenu Jha^{*} * School of Electrical and Computer Engineering, Rutgers University, New Brunswick, New Jersey, USA [†] Intel Corporation, USA [‡] EPCC, The University of Edinburgh, Edinburgh, UK

Research in Advanced DIstributed Cyberinfrastructure & Applications Laboratory (RADICAL) Rutgers University http://radical.rutgers.edu http://radical-cybertools.github.io

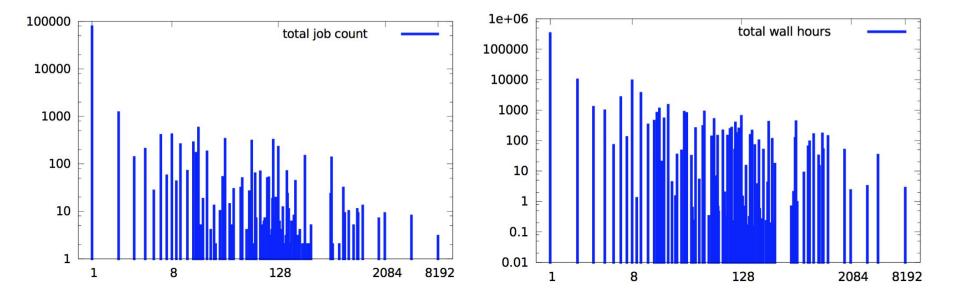
Extreme Scale "Task-level Parallelism" on HPDC

- Problems in computational science *naturally* amenable to "task level" parallelism computing
- Beyond HTC vs HPC
- Given access to X cores/nodes slice/dice or distribute as needed.
- Resources and workloads are characterised by a range of properties:





Blue Waters Job Size Distribution



Requirements / goals

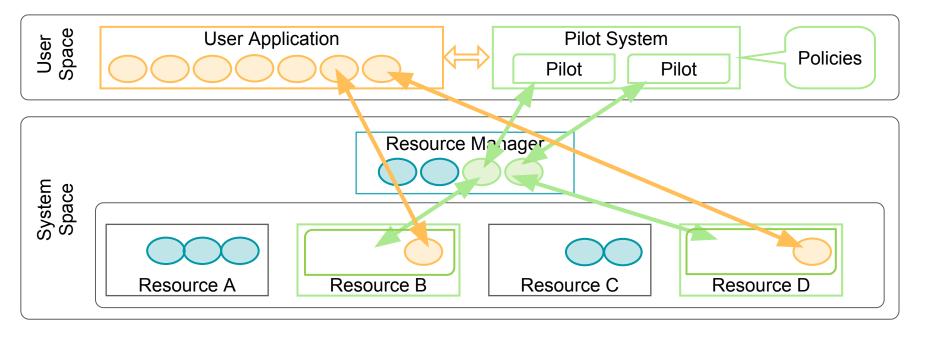
- Workload with heterogeneous tasks
 - Varying core count
 - Varying application
 - MPI / non-MPI
- Dynamic workload with workload unknown in advance
 - Task N+1 depends on task N
- Control over concurrency of tasks
 - Might be loosely coupled (e.g. replica exchange)
- ~10k concurrent tasks

(Why not) batch queue jobs

- Low throughput
 - Every job needs to queue
 - Breaks especially in dynamic workload situations
- No control over concurrency
- Limit on total concurrency
- Maximum of one task per node
- Job arrays are too inflexible (nor available on BW)
- Too many flavours

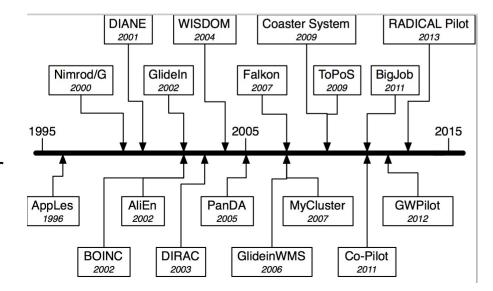
Pilot Abstraction

Working definition: A system that generalizes a placeholder to allow application-level control over acquired resources.



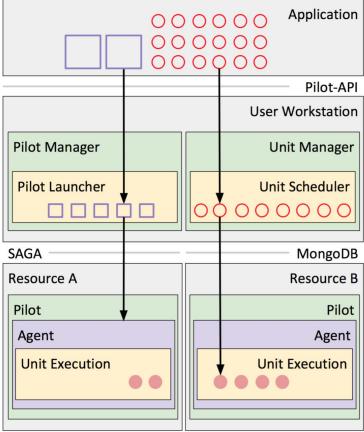
Advantages of Pilot-Abstraction

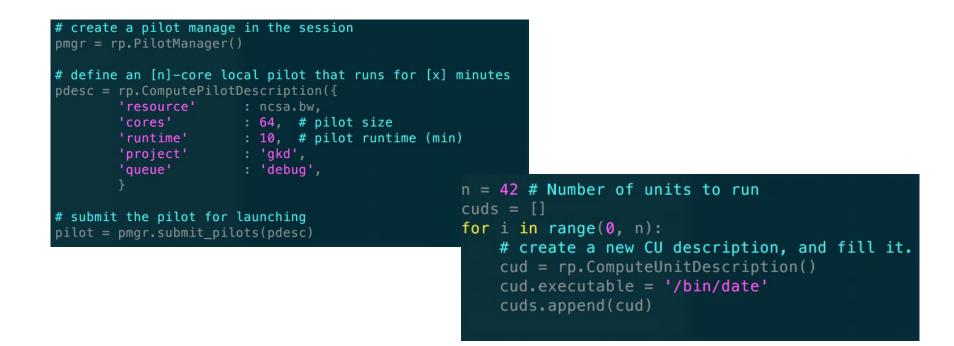
- Decouples workload management from resource management
- Flexible Resource Management
 - Enables the fine-grained "slicing and dicing" of resources
 - Tighter temporal control and other advantages of application-level Scheduling (avoid limitations of system-level scheduling)
- Build higher-level frameworks without explicit resource management



RADICAL-Pilot Overview

- Programmable interface, arguably unique:
 - Well defined state models for pilots and units.
- Supports research whilst supporting production scalable science:
 - Pluggable components; introspection.
- Portability and Interoperability:
 - Works on Crays, most known clusters,
 XSEDE resources, OSG, and Amazon EC2.
 - Modular pilot agent for different architectures.
- Scalable:
 - Agent, communication, throughput.





create a unit manager, submit units, and wait for their completion
umgr = rp.UnitManager(session=session)
umgr.add_pilots(pilot)
umgr.submit_units(cuds)
umgr.wait_units()

Agent Architecture

• Components:

Enact state transitions for Units

• State Updater: Communicate with client library and DB

• Scheduler:

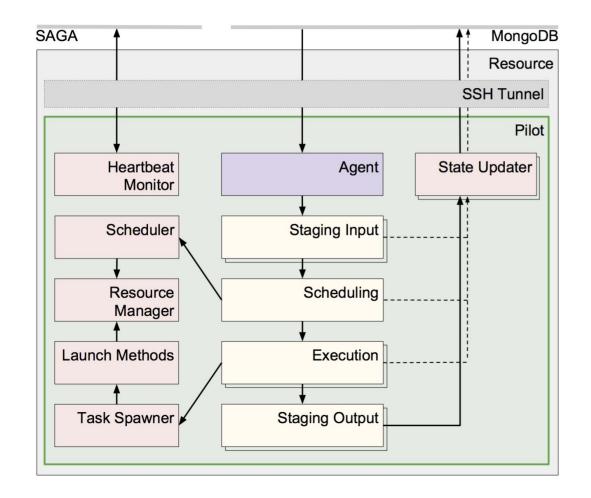
Maps Units onto compute nodes

- Resource Manager: Interfaces with batch queuing system, e.g. PBS, SLURM, etc.
- Launch Methods:

Constructs command line, e.g. APRUN, SSH, ORTE, MPIRUN

• Task Spawner:

Executes tasks on compute nodes



(Why not) RADICAL-Pilot + APRUN

- RP Agent runs on MOM node
- Uses aprun to launch tasks onto the worker nodes

- Low throughput (ALPS not designed for short/small tasks)
- Limit on total concurrency (1000 aprun instances)
- Maximum of one task per node

(Why not) RADICAL-Pilot + CCM

- Bootstrapper runs on MOM node
- Bootstrapper creates "cluster"
- Uses ccmrun to launch RP Agent into the "cluster"

• Not universally available

RADICAL-Pilot + ORTE-CLI (a bit better)

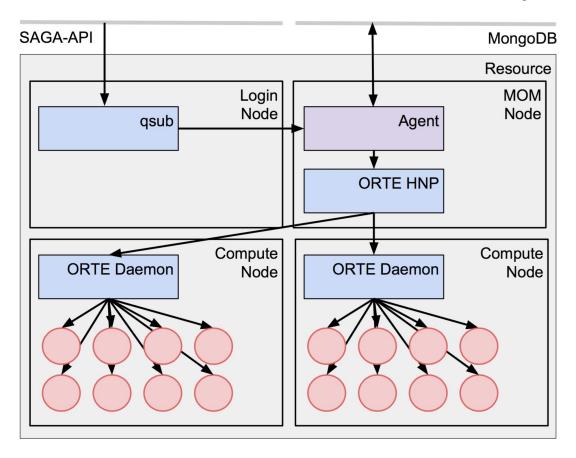
• ORTE: Open RunTime Environment

- Isolated layer used by Open MPI to coordinate task layout
- Runs a set of daemons over compute nodes
- No ALPS concurrency limits
- Supports multiple tasks per node
- orte-submit is CLI which submits tasks to those daemons
 - 'sub-agent' on compute node that executes these
 - Limited by fork/exec behavior
 - Limited by open sockets/file descriptors
 - Limited by file system interactions

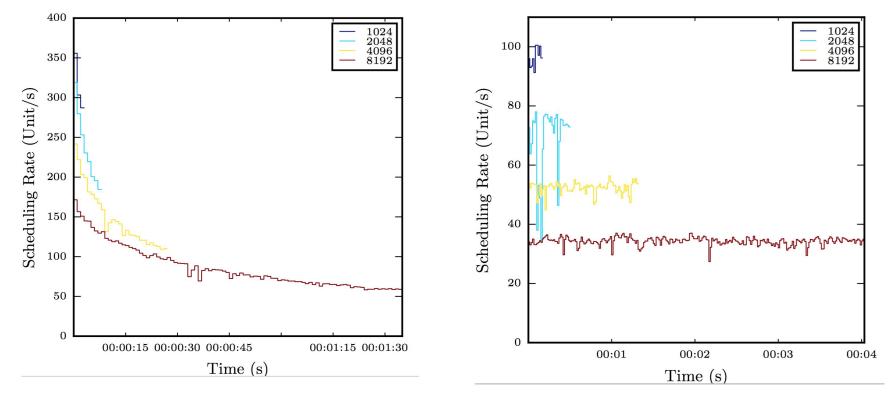
RADICAL-Pilot + ORTE-LIB (much better)

- All the same as ORTE-CLI, but
 - Uses library calls instead of orte-submit processes
 - No central fork/exec limits
 - Shared network socket
 - (Hardly) no central file system interactions

RADICAL-Pilot + ORTE on Cray



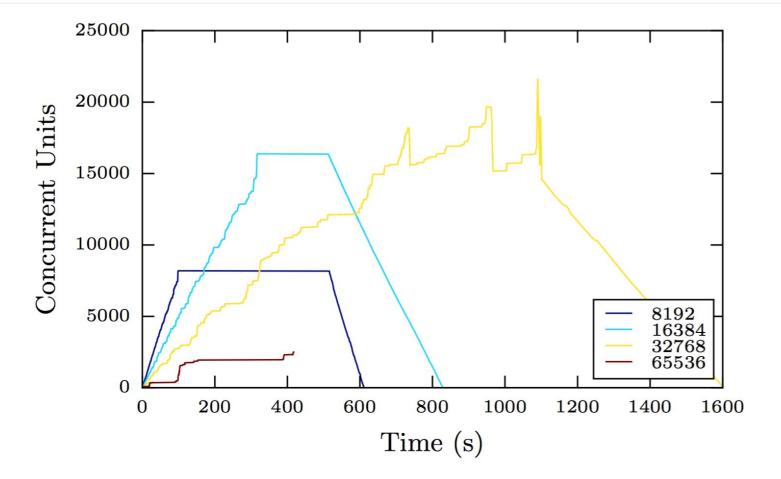
Micro Benchmark: Scheduler



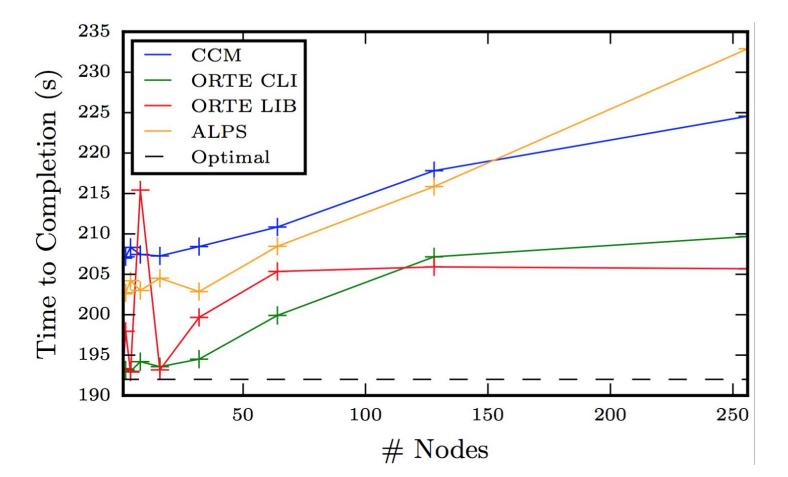
Scheduling only

Scheduling and unscheduling

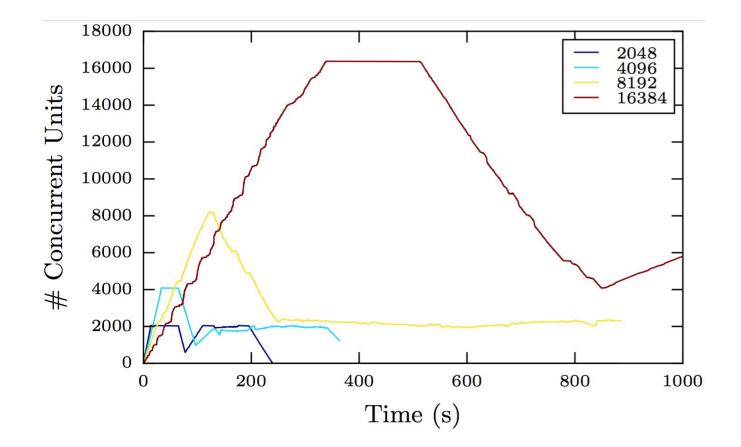
Micro Benchmark: Executor Scaling



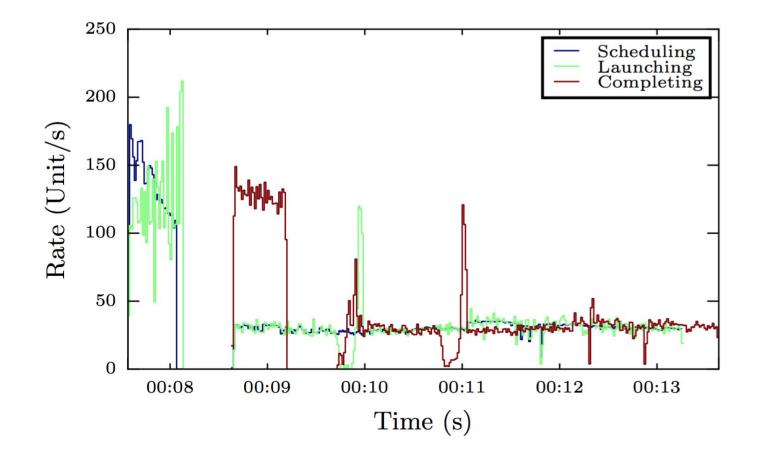
Agent Performance: Full Node Tasks (3 x 64s)



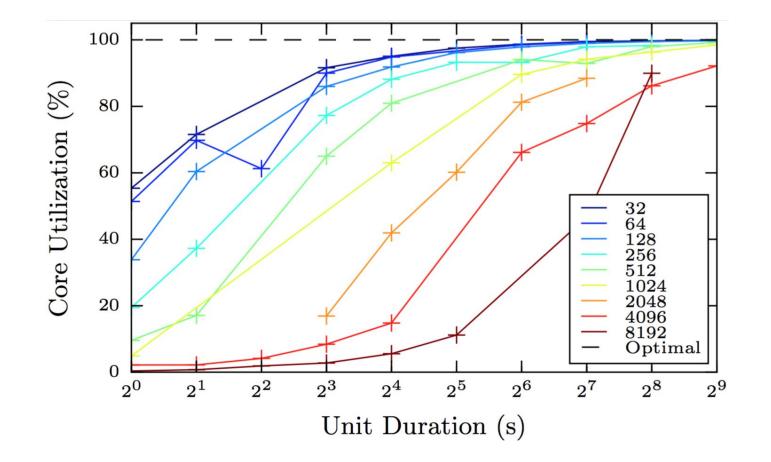
Agent Performance: Concurrent Units (3x)



Agent Performance: Turnaround (3 x 4k x 64s)



Agent Performance: Resource Utilization



Conclusion

- There is no "one size fits all" in HPC
- With general tools extend functionality of Cray HPC systems
- Achieved 16k concurrent tasks
- Launch rate of ~100 tasks / second
- Efficiency large dependent on task count and duration
- Cray specific PMI excludes running Cray MPI linked applications

Future work

• RADICAL-Pilot

- Bulks all the way
- Agent scheduler overhaul
- Topology aware task placement
- Heterogenous node scheduling (I.e. GPU)

• ORTE

- Fabrics-based inter-ORTE communication
- Optimize ORTE communication topology

References

- RADICAL-Pilot: Scalable Execution of Heterogeneous and Dynamic Workloads on Supercomputers
 - <u>http://arxiv.org/abs/1512.08194</u>
- A Comprehensive Perspective on the Pilot-Job Systems
 - <u>http://arxiv.org/abs/1508.04180</u>
- RADICAL-Cybertools overview
 - http://radical-cybertools.github.io/
- RADICAL-Pilot Github
 - <u>https://github.com/radical-cybertools/radical.pilot</u>
- RADICAL-Pilot Documentation
 - <u>http://radicalpilot.readthedocs.org/</u>

Micro Benchmark: Exec Rate + Concurrency (1x4k)

