

# What You Need to Know About KNL

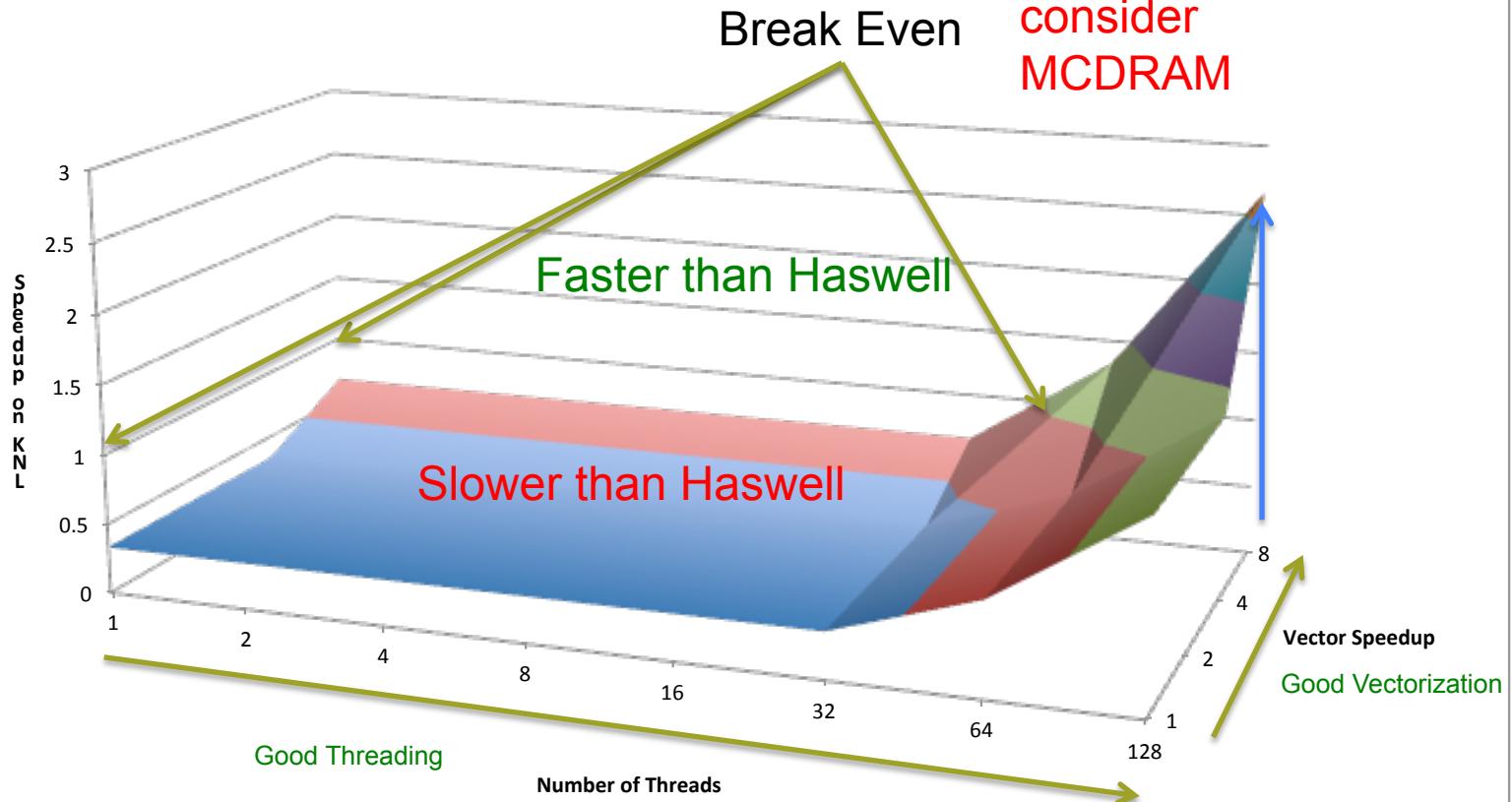
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Trinity  
Member of CTO Office

# Setting expectations, when will KNL out-preform state-of-the-art Xeon



- Xeon is faster
  - Haswell 2.3 GHZ
  - KNL 1.4 GHZ
- KNL has more cores
  - Haswell has 32 cores (2-16 core sockets)
  - **KNL has 68 cores**
- KNL has more threads
  - Haswell allows 64 threads – 32x2 Hyperthreads
  - **KNL allows 272 threads – 68x4 Hyperthreads**
- KNL has longer SIMD unit
  - Haswell has SIMD unit of length 4
  - **KNL has SIMD unit of length 8**
- **KNL has High Bandwidth Memory – 4-5 times Xeon bandwidth**
- **KNL has less low level cache**

## KNL versus Haswell Simple Amdahl's Law



# Important Architectural Data for HPC chips



Arch	MIMD Procs	SIMD DP Width	Close Memory (Cache) KB/core	HBM GB	HBM GB/sec	Main Memory GB	Main Memory GB/sec
K20X	14	32	90	6	250	-	
K40	15	32	96	12	288	-	
K80(/2)	15	32	96	12	240	-	
P100	56	32	73	16	720	-	
Ivybridge	8	2	2887	-	-	128	100
Haswell	16	4	2887	-	-	256	120
KNL (MCDRAM)	68	8	557 (25264MB)	16	480 (329)	384	80

# Important Data for Haswell and KNL

	Latency Haswell Nanosec (clocks)	Haswell Size /core	Latency KNL Nanosec (clocks)	KNL Size /core	Bandwidth Haswell GB/ sec Stream Triad	Bandwidth KNL GB/sec Stream Triad
L1 Cache	2.7 (6)	32KB	2.9 (4)	32KB		
L2 Cache	8.11 (19)	256KB	13.6 (19)	512KB		
MCDRAM (Cache) (Level 3)	24.35 (56)	2.5MB	173.5 (243)	242MB		329
MCDRAM (Flat)			174.2 (244)	242MB		486
DDR (Flat)			151.3 (212)		102	90
DDR (Cache)						59

# Most Important Feature of KNL is MCDRAM



- Most of our applications are memory bandwidth limited.
- MCDRAM presents a smaller high bandwidth memory that can be used to address memory bandwidth issues
  - When used as a separate address space (FLAT) it has the highest Triad Stream transfer rate of 480 GB/sec and a slightly higher latency of 174 nanoseconds than DDR memory
  - When used as cache (CACHE) it has a lower Triad Stream transfer rate of 329 GB/sec and about the same latency MCDRAM Flat



# MCDRAM as Cache

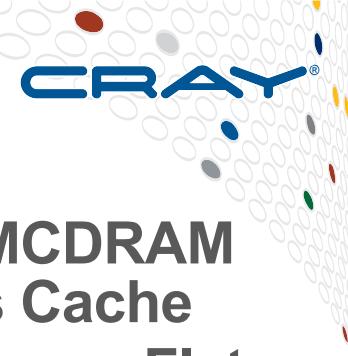
- Latency is much longer than typical caches
- Bandwidth is excellent albeit slower than MCDRAM as separate address space (480 versus 325).
- Only way to transfer pages to MCDRAM without going through low level caches is to use it as Cache
- MCDRAM as Cache has one major potential hazard. It is direct mapped and some applications may experience significant degradation in performance. This degradation is caused by some nodes of a large parallel run experiencing cache conflicts and introducing load imbalance.



# MCDRAM as Flat

- **Latency is approximately the same as DDR latency – little slower**
- **Bandwidth is excellent – five times faster than DDR.**
- **Two considerations**
  - When application and data require less than 16 GBytes
    - Easy to use - `numctl --membind = 1`
  - When application and data require more than 16 GBytes
    - Very difficult to get to run better than using MCDRAM as Cache
    - All Flat means that low level cache is as poor as accelerators
    - 50% or 75% Flat should work better; however, identifying which arrays to put into Flat is a complex problem
    - Using Flat as a scratch pad, swapping data in and out ruins contents of low level
      - Still being considered by some application developers
- **Using MCDRAM as flat can be used to increase the total memory on the node – when memory size is more important than maximum performance.**

# What you need to know about MCDRAM



- Very few applications (if any) have shown that using MCDRAM as Flat gives an improvement over using MCDRAM as Cache
- If your application plus data uses less than 16 Gbytes, use Flat with job launch command numactl --membind=1
- If your application plus data uses more than 16 Gbytes, use Cache
- If you use Cache – check to see if your application is susceptible to issues caused by Direct Mapped Cache
  - Run application with DDR (numactl --membind=0) only up to 100 to 1000 nodes
  - Run application with Cache up to 100 to 1000 nodes
  - If DDR only scales better than Cache then Cache has issues with cache conflicts



# Closer Look at MCDRAM

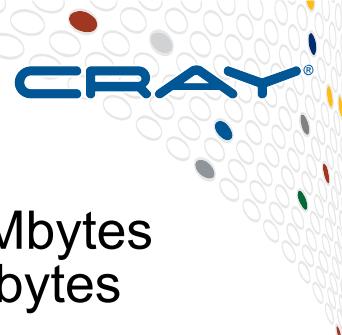
Presentation by Nathan Wichmann



# Different Triads

- **Triad**
  - $a(j) = b(j) + \text{scalar} * c(j)$
  - Simple STREAM Triad with 3 stride 1 arrays
- **Triad – Stride -1 Gather**
  - $a(j) = b(j) + \text{scalar} * c(\text{index1}(j))$
  - Triad with one arrays loaded using index array, but index is still a stride-1 pattern
- **Triad – Random Gather**
  - $a(j) = b(j) + \text{scalar} * c(\text{indexr}(j))$
  - Triad with one arrays loaded using index array, but index is still a **random** pattern
- **Triad – Full Random Gather / Scatter**
  - $a(\text{indexr}(j)) = b(\text{indexr}(j)) + \text{scalar} * c(\text{indexr}(j))$
  - Full random pattern on all three arrays

# Array Sizes and footprints



- **Each of the arrays are 2M elements (15.26 Mbytes)**
  - Simple Triad uses 3 arrays and thus touches a total of ~46 Mbytes
  - Triads that use index use 4 arrays and thus a total of ~61 Mbytes
- **Each core works on its own copy for a total footprint of 3.1-4.1 Gbytes per node**
  - No L3 cache reuse on Haswell
  - But should be able to easily reside inside of 50% of MCDRAM, no matter the mode
- **Ideal problem to showcase the advantage of MCDRAM**
  - High bandwidth test that fits in MCDRAM but cannot be blocked for L3 on XEON

# Different KNL modes



- **Quad – 0% cache**
  - Also known as “flat mode”
  - All of MCDRAM must be explicitly targeted
- **Quad – 100% cache**
  - All of MCDRAM is a direct mapped cache
  - MCDRAM is not available for capacity
  - Don’t have to do anything to the codes
- **Quad – 50% cache**
  - Also known as “hybrid mode”
  - 50% of MCDRAM is a direct mapped cache
  - 50% of MCDRAM must be explicitly targeted
  - Pragmatic compromise?

# Fastmem directives

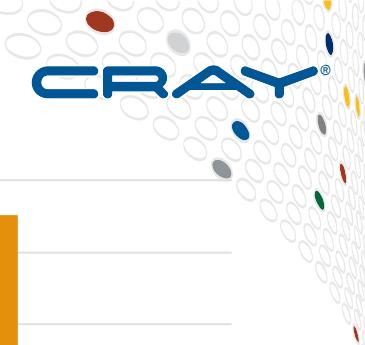


- Add the !DIR\$ ATTRIBUTES FASTMEM :: X directive when you declare the arrays.

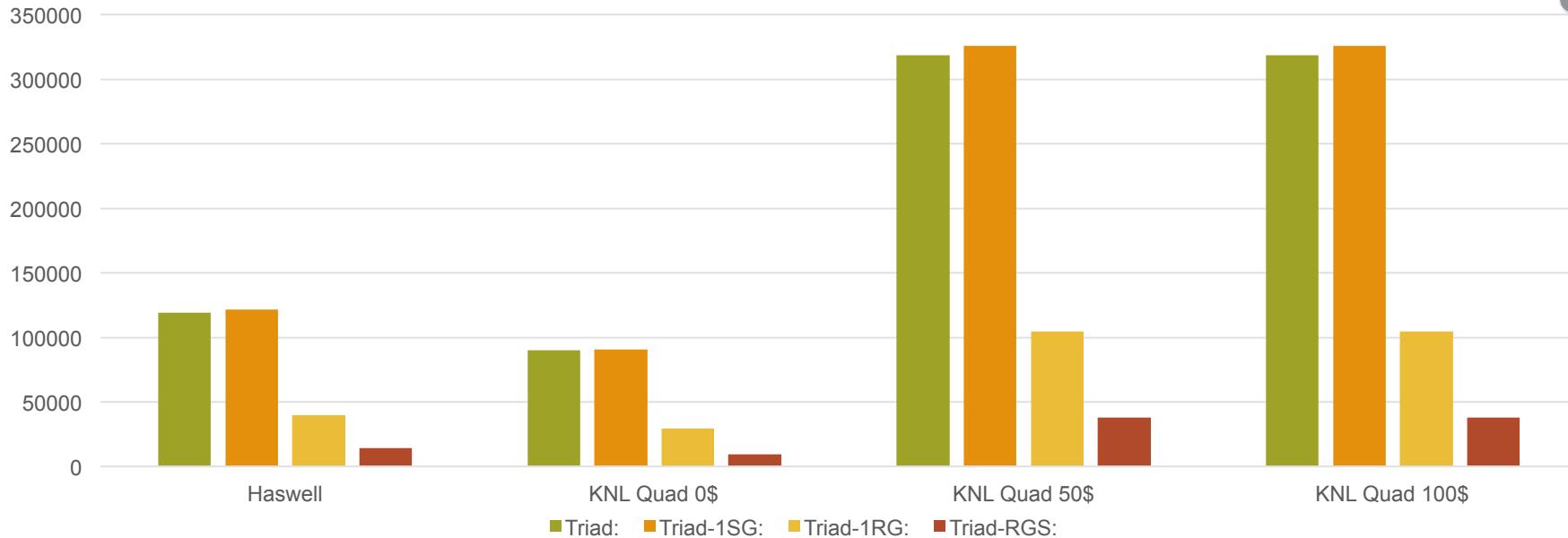
```
module swap PrgEnv-cray PrgEnv-intel
#Have to find a memkind library until it gets added to CLE
(which will be soon).
module use /cray/css/users/kjt/opt/modulefiles
module load memkind
```

```
module mods
  real(kind=8), allocatable :: a(:),b(:),c(:),d(:),x(:)
!dir$ attributes fastmem :: a,b,c,d,x
end module mods
```

# No arrays in explicit fastmem



No fastmem directives



- Haswell is faster than 0% cache
- But 50% and 100% cache modes work well to capture available reuse

COMPUTE

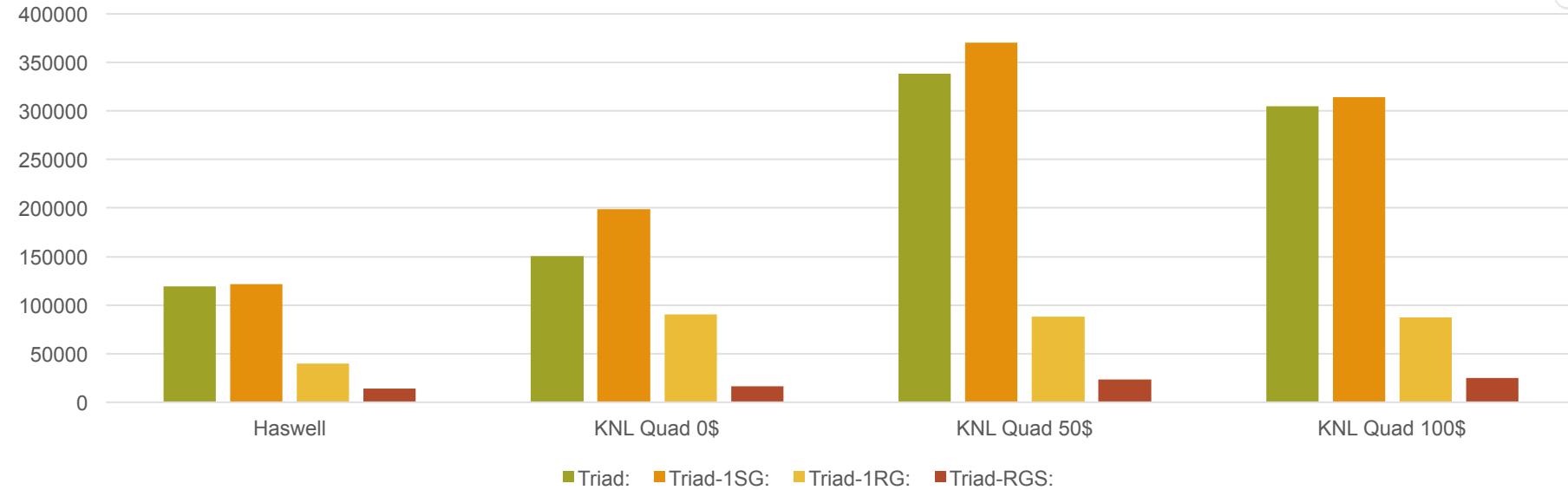
STORE

ANALYZE

# Most arrays in explicit fastmem



Fastmem b, index\*, c

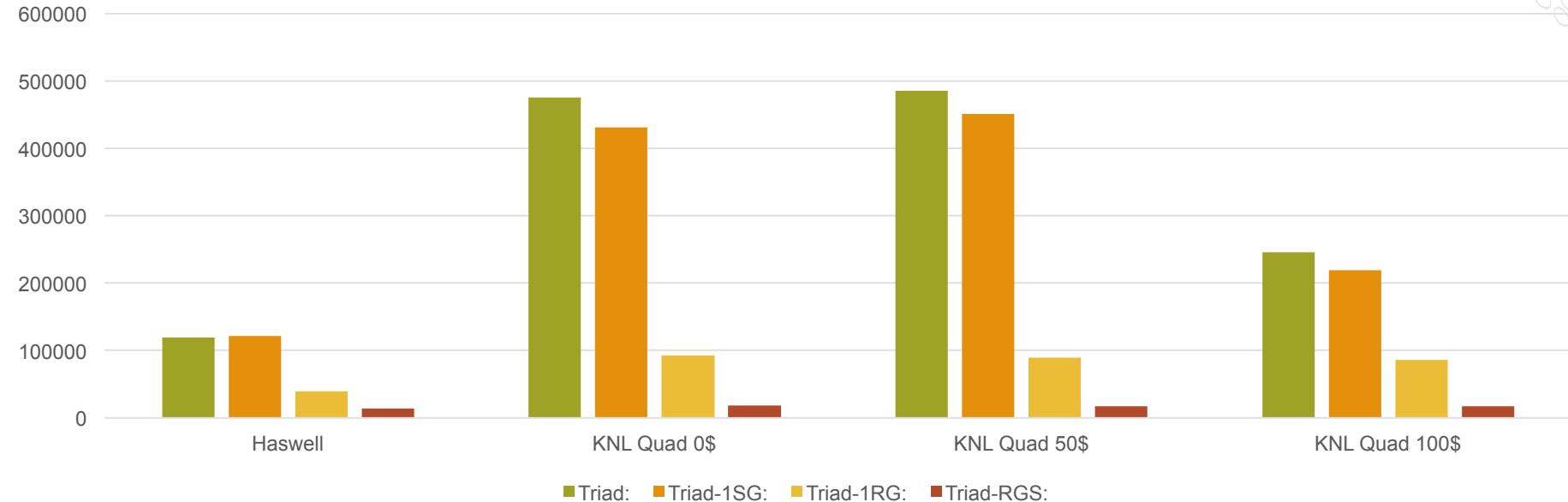


- Even with most of the arrays in fastmem, 0% still falls short of cache for stride-1 cases

# ALL arrays in explicit fastmem



Fastmem b, index\*, c, a



- **0% mode finally does very well with all arrays in fastmem**
  - But 50% cache easily keeps up because its fastmem can also hold the data
- **100% cache does poorly if combined with fastmem directive**
  - Why?

COMPUTE

STORE

ANALYZE



# Let's prematurely jump to some observations and conclusions

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C O M P U T E

| S T O R E

| A N A L Y Z E

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# 0% cache mode



- Maximizes amount of explicit space
- Highest performance if you can fit \*all\* of your data
- But performance can quickly suffer if you miss some MCDRAM reuse opportunity
  - Amdahl's Law
- Conclusion
  - If your full data set can fit inside of 16 Gbytes per node, just use this mode

# 100% cache mode



- **Maximizes amount of cache space**
- **Does not require any directives**
- **It seem to capture reuse**
- **Conclusion**
  - A great starting point!
    - You might not be able to do any better than this

# 50% cache mode

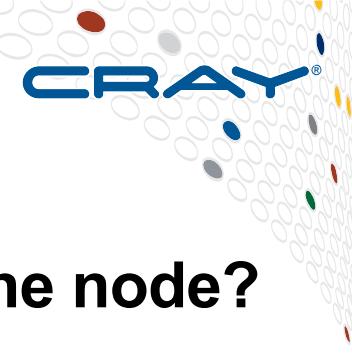


- **Cuts cache space by half**
  - How often will that matter?
- **Captures a lot of reuse with no work**
- **Allows for some incremental upside**
- **0% cache will beat this only if 100% of bandwidth sensitive data fits in 16 GB but does not fit into 8GB**
- **Conclusion**
  - Gives good performance in a number of situations, but allows for one to experiment with fastmem directives
  - The mode has not been adequately tested yet



**That sounds good,  
but wait...  
just one problem**

# Single node vs multi-node

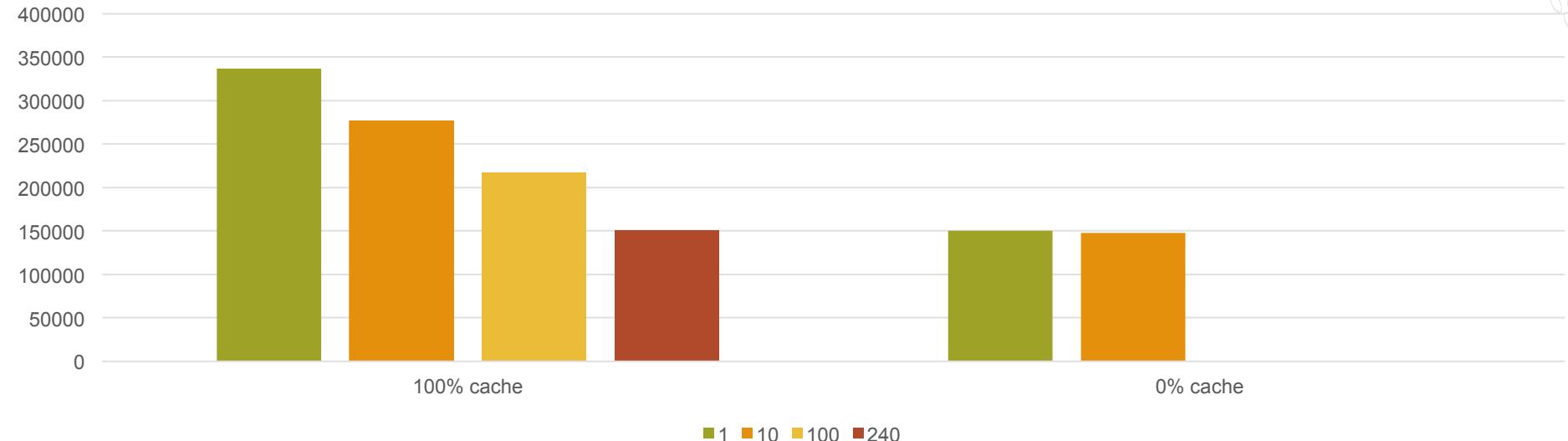


- All of these tests were run on just a single node
- Does the picture change if we run on more than one node?

# Increase fastmem for 100% cache mode



Performance per node when using multiple nodes  
Simple Triad - Fastmem b, index\*, c



- The performance of cache drops significantly when more nodes are used!
- Why?

COMPUTE

STORE

ANALYZE

# Basic picture of memory and cache



MEM							
MEM							
MEM							
MEM							
MEM							
MEM							
\$	\$	\$	\$	\$	\$	\$	\$

- There are 6 memory locations that map to the same location in cache
  - Assumes a 96 GB DRAM config. Larger memory have more
- Cache is direct mapped. I.E. there is only 1 “way”

# Start filling pages in memory



					PE3's A		
						PE5's A	
	PE4's A						
			PE2's A				
PE1's A							
\$PE1's A	\$PE4's A	\$	\$PE2's A	\$	\$PE3's A	\$PE5's A	\$

- OS starts to place pages that contain the variable “A” for each PE as those PEs reach the allocation statement
  - Each page placement is more or less “random” based on when various PEs arrive at the allocate, and the order in which the pages were free, perhaps even by a previous program

# Conflicts are bound to happen



			PE6's A				
					PE3's A		
						PE5's A	
	PE4's A						
			PE2's A				
PE1's A							
\$PE1's A	\$PE4's A	\$	\$ conflict	\$	\$PE3's A	\$PE5's A	\$

- If two PEs A are aliased to the same cache location, then thrashing may occur if those PEs are both reuse A at “about” the same time

COMPUTE

STORE

ANALYZE

# Cache trashing will occur, but how often?



- **Good news: The cache is really large, so there is a relatively low probability of an aliasing conflict occurring**
  - Conflict probability is also a function of size of the data being reused
  - This is why we often don't observe this on single node runs
- **Bad news: All of these are NON-ZERO probabilities, and the dice are rolled on every node in the job**
  - As the number of nodes in the job increases, the probability and aliasing problems will occur approaches 100%
    - If one node sees it in one part of the code, all nodes will see it via synchronization
    - Furthermore the slow PEs can be different for different parts of the code
  - Hyperthreading only makes the problem worse as more threads are putting more pressure on the cache levels
  - OpenMP may lower the probability in some cases, but since different threads will hit different memory sections, the problem will still occur
  - This is not limited to “high-bandwidth” arrays, but even those arrays that may get good L2 cache hit rates

# Direct mapped cache causing scaling problems



- If your performance is impacted by the effectiveness of the MCDRAM cache, you may experience scaling problems
  - This will likely show up as communication, but will be because of synchronization, not bandwidth or latency constraints
- Normal profiling may not point to the offending compute region
  - Only a few PEs might be slow, and thus that signal could be drowned out by the other PEs in the job
- We may be able to limit this problem by running in 50% cache and putting some arrays into fastmem
  - But this remains a theory and so far has not worked out in at least one case

# Conclusions



- Flat mode is good if your entire data set can fit into 16GB
- Flat mode seems likely to be unforgiving if some bandwidth data does not fit
- Cache mode seems to capture reuse well
  - And it requires no work on the part of the benchmarker
- But is highly susceptible to thrashing if important data aliases to the same location
  - This becomes more and more likely as node counts increase
- Each mode has serious issues and seem likely to cause performance problems as one tries to use more nodes or run larger problems



# Next most important feature of KNL

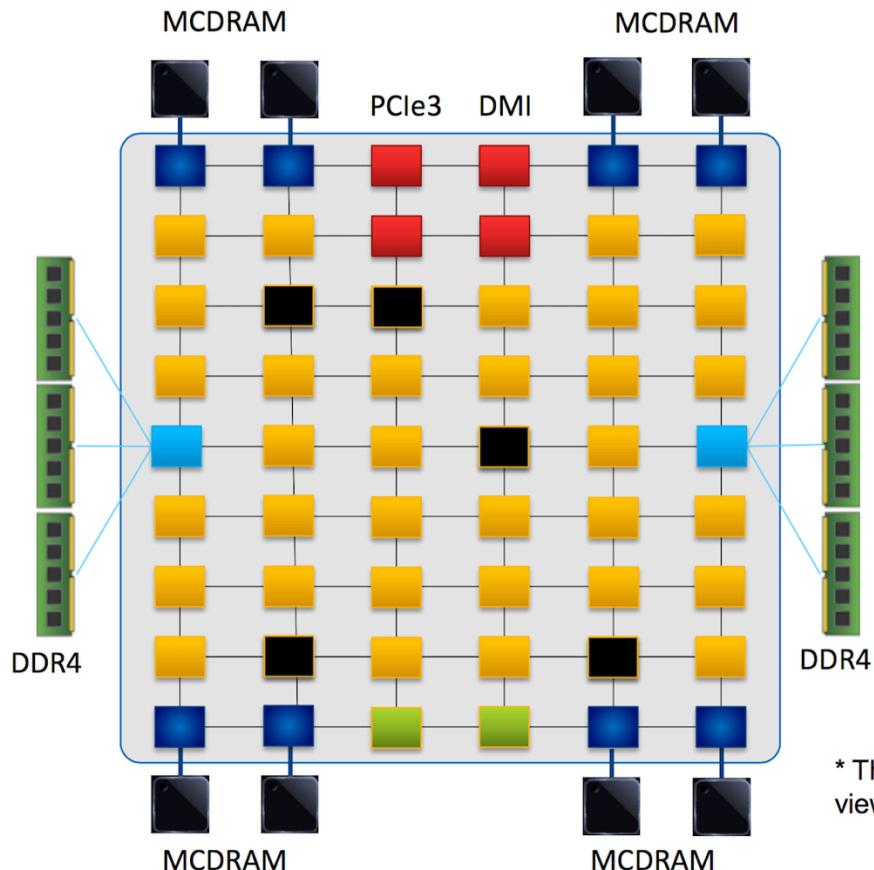
- **Vectorization**
  - KNL can generate up to 16 results each clock cycle
    - Eight Adds and Eight Multiplies
  - KNL has predicated execution
    - Has a mask register to handle IF conditions
  - KNL has some hardware assist for gather/scatter
    - Still vectorizing loops with indirect addressing does not give good performance
  - Vector operations on operands aligned on cache boundaries perform better than misaligned operands.
- **Vectorization will not perform as expected if the loop is memory bandwidth limited**
  - Cache utilization is more important than vectorization



# Next most important feature of KNL

- **Lots of hardware threads on the node**
  - The KNL node is a 2-D mesh of 36 tiles
    - Each tile has two cores
      - Each core allows four hardware threads
    - On current systems two of the tiles are disabled, resulting in 34 usable tiles
  - Each tile has
    - Two cores with Level 1 cache
    - Shared Level 2 cache
    - Cache Homing Agent (CHA) which handles cache coherency across the node

# Tile To NUMA Node Mapping Concepts

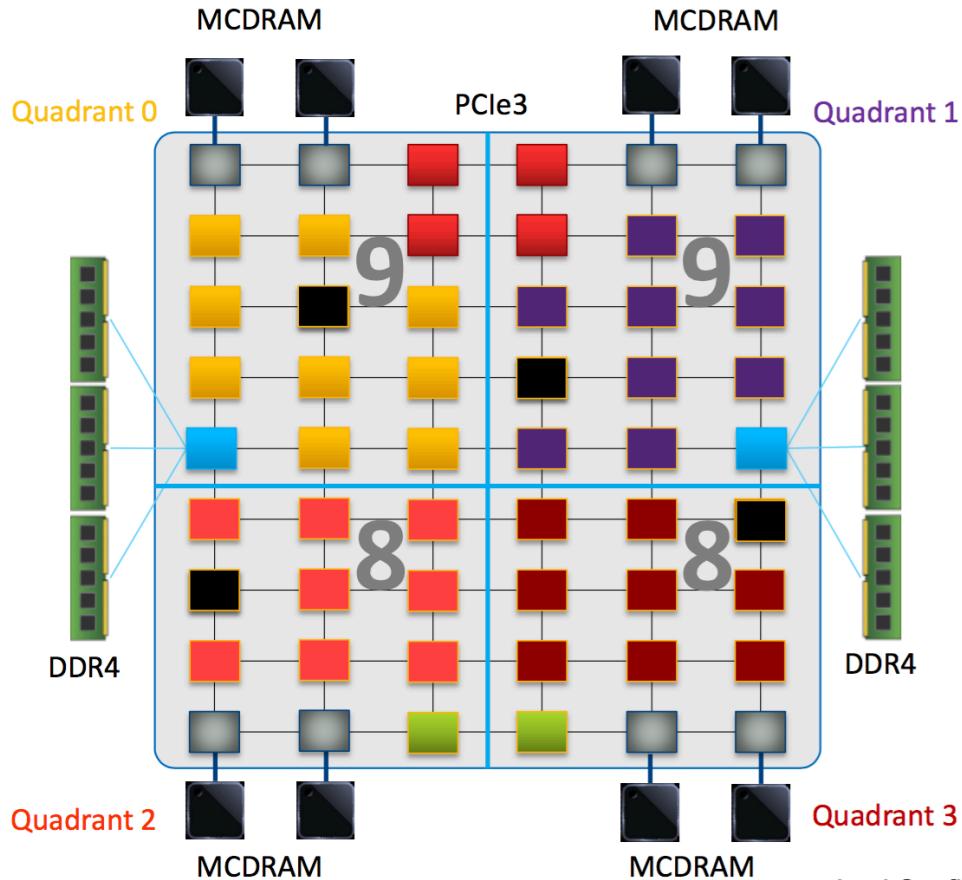


- KNL has 38 processing tiles per chip
- 68 core part will have 4\* tiles disabled at the factory
- ApicID is hardwired into each tile
- Disabled tiles will vary from part to part
- As a result, ApicIDs will vary from part to part

- Processing Tile**
- MCDram Controller**
- DDR4 Controller**
- PCIe, DMI Interfaces**
- Other**
- Factory disabled**

\* This slide has animation showing the random location of 4 disabled tiles. When viewed in .pdf format it looks like 5 tiles have been disabled.

# 68 Core SNC4 Tile Mapping Example – Intel BIOS



- Like SNC2, it is up to the BIOS to define and implement the tile assignments
- There are no hard wired circuits within KNL as to how tiles are mapped to NUMA nodes
- Best configuration is to have 9 tiles per quadrant on the top and 8 tiles on the bottom
- This is the preferred model that will be implemented by Intel BIOS

# Options for using all the hardware threads



## ● All - MPI

- Works surprisingly well; however, typically MPI does not run well on hyper-threads – maybe 64 MPI tasks on the node is best; however, I have seen 128 MPI tasks on the node run better than 64. Should be tried for those applications that have no shared memory threading before spending a lot of work on application.

## ● All – Shared Memory threads

- VERY Few applications are threaded well enough to utilize the KNL node in this way (TBD)

## ● Hybrid MPI/Shared Memory threads

- Best approach for those applications that already have shared memory threading in the application
- Need to find “Sweet Spot”

# What you need to know about the KNL node



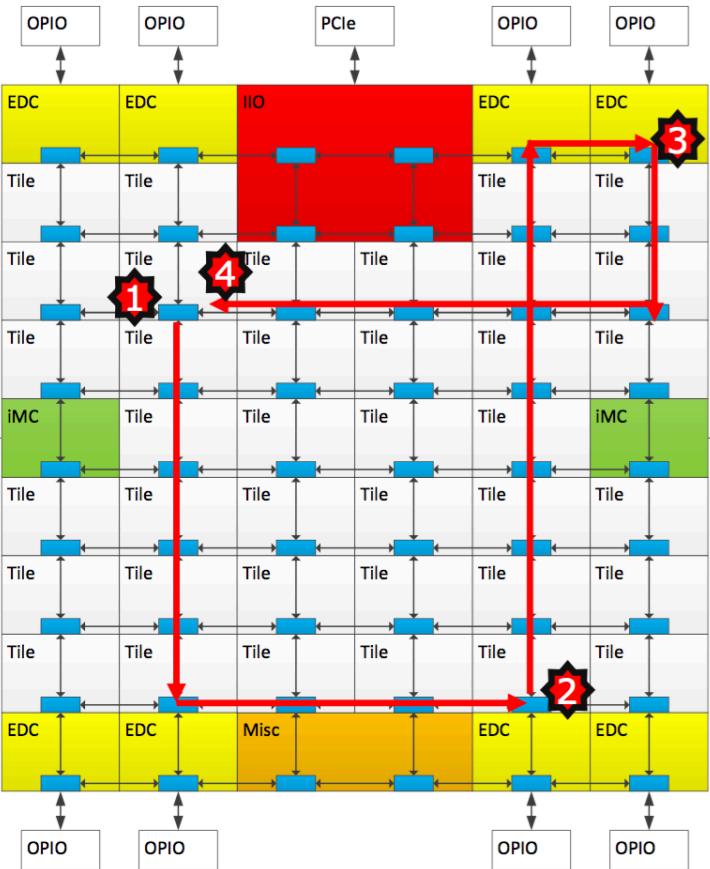
- **KNL has significant NUMA issues on the node which introduces inefficiencies in shared memory threading**
  - High level OpenMP with first touch are a necessity for good scaling
- **The tile is the closest thing to a UMA region – should we at least use 34 MPI tasks on the node with 2,4 or 8 threads per MPI task?**
- **A Quadrant is the next memory region when using SNC4**
  - NUMA issues within the Quadrant

# Investigate Clustering Modes



- **Three Major cluster modes being used**
  - Quad
  - SNC4
  - SNC2
- **Couple more that are not used**
  - All to All
  - Hemisphere

# Cluster mode: All-to-All



**Address uniformly  
hashed across all  
distributed directories**

## Typical Read L2 miss

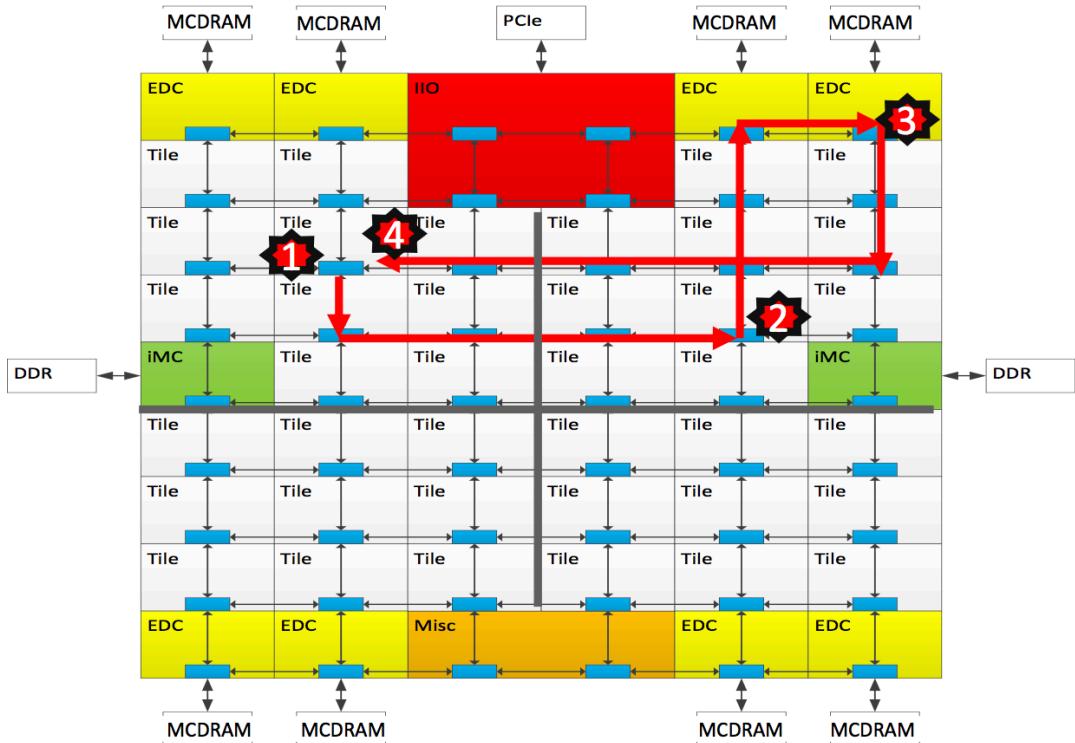
1. L2 miss encountered
2. Send request to the distributed directory
3. Miss in the directory. Forward to memory
4. Memory sends the data to initial requestor

# Cluster – All to All



- Has lowest bandwidth and highest latency of all clustering modes
- Do not know of any application using this today

# Cluster Mode: Quadrant



Chip divided into four virtual Quadrants

Address hashed to a Directory in the same quadrant as the Memory

Affinity between the Directory and Memory

Lower latency and higher BW than all-to-all. SW Transparent.

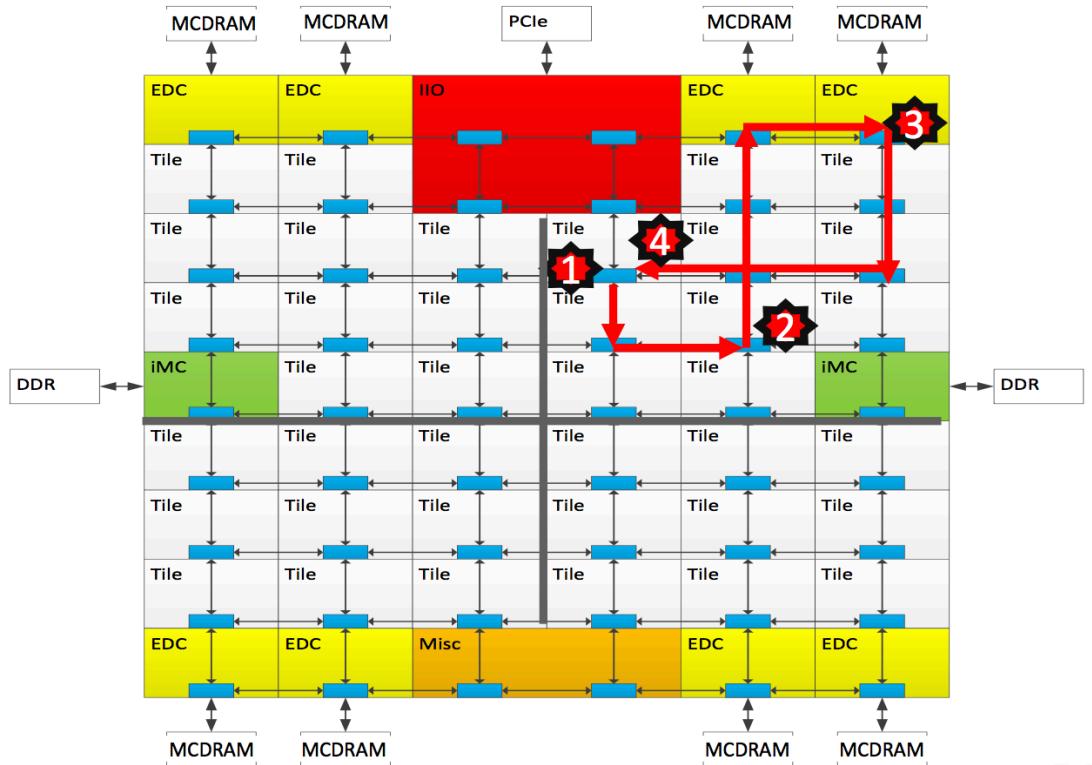
- 1) L2 miss, 2) Directory access, 3) Memory access, 4) Data return

# Cluster – Quad



- Allows a collection of cores within a quadrant to get all of the memory bandwidth
  - Could be important for a bandwidth sensitive, load imbalanced mostly-MPI application.
    - Mostly MPI = greater than/or equal to 16 MPI tasks per node
- Introduces additional NUMA affects when threading across more than one tile

# Cluster Mode: Sub-NUMA Clustering (SNC)



Each Quadrant (Cluster) exposed as a separate NUMA domain to OS.

Looks analogous to 4-Socket Xeon

Affinity between Tile, Directory and Memory

Local communication. Lowest latency of all modes.

SW needs to NUMA optimize to get benefit.

- 1) L2 miss, 2) Directory access, 3) Memory access, 4) Data return

# Cluster – SNC2 – SNC4



- Introduces 2 or 4 socket boundaries to node.
  - Cons
    - Latency and Bandwidth between sockets reduced
    - Collection of cores can only access  $\frac{1}{4}$  memory bandwidth and  $\frac{1}{4}$  of DDR and  $\frac{1}{4}$  MCDRAM efficiently
  - Pros
    - Assists in improving locality when threading across multiple tiles, when `OMP_NUM_THREADS > 8`
    - Has the best latency to memory of all the clustering modes

# Simple Rules for Clustering

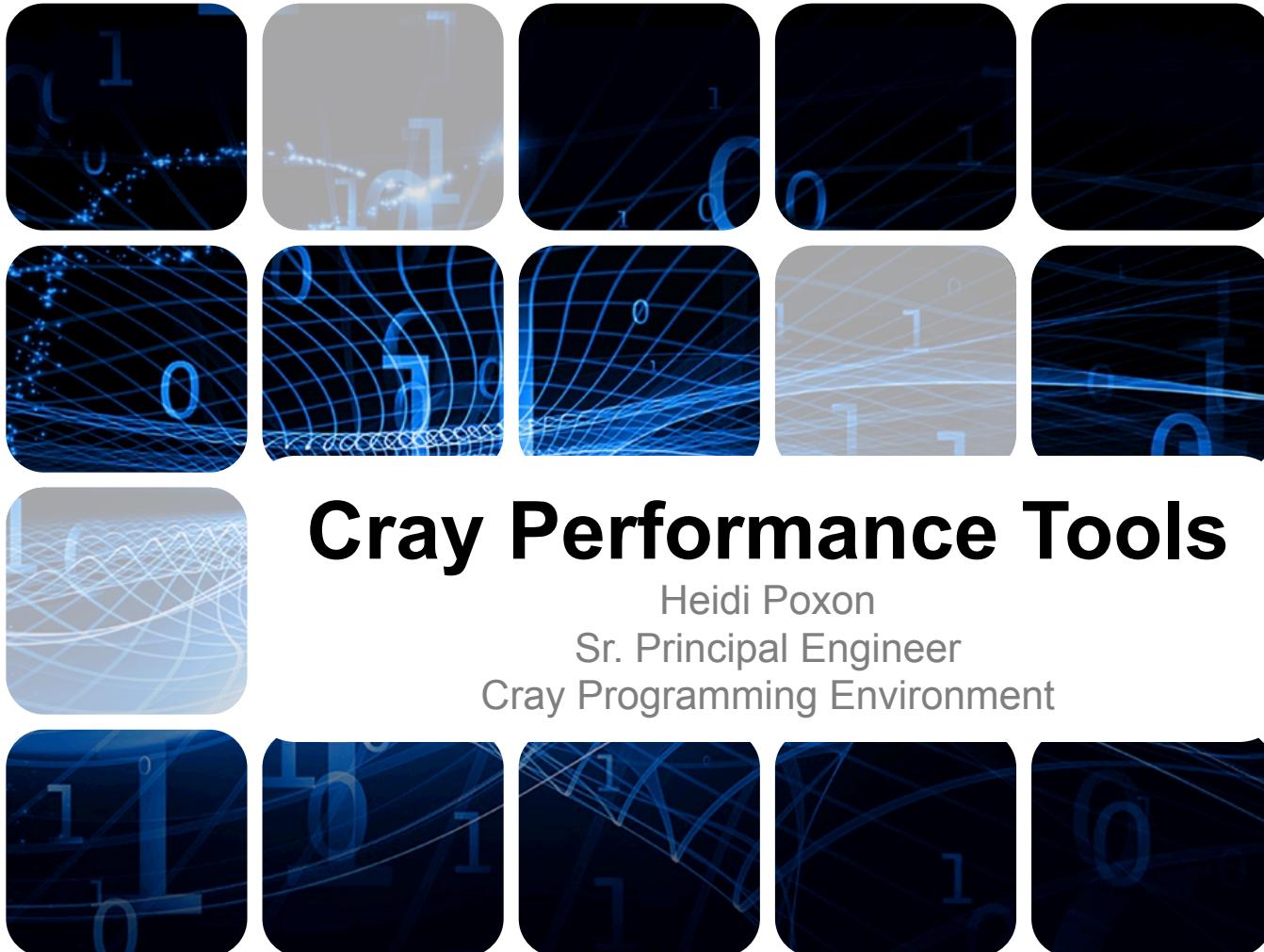
- **If the number of MPI tasks is greater than or equal to 32**
  - If your application is memory bandwidth limited, Quadrant would be best
  - If your application is latency bound, SNC4 should be best
- **If the number of MPI tasks is less than or equal to 16**
  - To reduce the NUMA effects SNC2 or SNC4 should be used
  - If your application is memory bandwidth limited, Quadrant would be best
- **Most of the time the differences in clustering modes will be less than 10% - will be some outliers.**
- **Find out what best for you and stick with it**
- **Nightmare for operational system because re-provisioning takes a long long time (Today 20-25 minutes)**

# A simple flowchart for examining an application with Perftools



# Before We get started

- Many of the tools discussed will work with the Intel compiler; however, we highly recommend that you try to build with the Cray compiler, since the following only works with the Cray compiler
  - Memory Analysis tools
  - Loop profiling
  - Reveal
  - Profiling of library groups

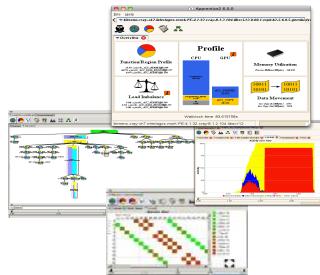


# Cray Performance Tools

Heidi Poxon  
Sr. Principal Engineer  
Cray Programming Environment

# Using Cray Performance Tools

Load modules to access software



Build and instrument program

Run program and view results



# Two Modes of Use

- **CrayPat-lite** for novice users, or convenience
- **CrayPat** for in-depth performance investigation and tuning assistance
- Both offer:
  - Whole program analysis across many nodes
  - Indication of causes of problems
  - Suggestions of modifications for performance improvement

# “Lite” Mode

Load performance tools instrumentation module

```
$ module load perftools-lite
```

Build program  
(no modification to makefile)

```
$ make
```

Run program  
(no modification to batch script)

```
$ aprun a.out
```



a.out (instrumented program)

Condensed report to stdout  
a.out\*.rpt (same as stdout)  
a.out\*.ap2  
files

# Example CrayPat-lite Output



```
#####
#                                #
#      CrayPat-lite Performance Statistics    #
#                                #
#####
```

CrayPat/X: Version 6.3.2.461 Revision 56930ff 02/01/16 15:31:33  
Experiment: lite lite/sample\_profile  
Number of PEs (MPI ranks): 64  
Numbers of PEs per Node: 32 PEs on each of 2 Nodes  
Numbers of Threads per PE: 1  
Number of Cores per Socket: 16  
Execution start time: Tue Feb 2 18:53:50 2016  
System name and speed: kay 2301 MHz (approx)

Avg Process Time: 64.38 secs  
High Memory: 1,563 MBytes 24.43 MBytes per PE  
MFLOPS: Not supported (see observation below)  
I/O Read Rate: 48.514130 MBytes/sec  
I/O Write Rate: 22.281350 MBytes/sec  
Avg CPU Energy: 41,820 joules 20,910 joules per node  
Avg CPU Power: 649.53 watts 324.77 watts per node

Table 1: Profile by Function Group and Function (top 10 functions shown)

Samp%	Samp	Imb.	Imb.	Group
		Samp	Samp%	Function
PE=HIDE				
100.0%	6,156.2	--	--	Total
66.3%	4,082.5	--	--	USER
11.8%	729.2	48.8	6.4%	mult_su3_mat_vec_sum_4dir
10.2%	629.4	49.6	7.4%	mult_adj_su3_mat_4vec
6.1%	377.1	28.9	7.2%	mult_su3_nn
5.9%	365.4	42.6	10.6%	mult_su3_na
5.4%	329.4	37.6	10.4%	scalar_mult_add_lathwvec_proj
3.8%	232.9	39.1	14.6%	mult_su3_sitelink_lathwvec
25.3%	1,557.0	--	--	MPI
12.8%	789.3	163.7	17.5%	MPI_Wait
6.7%	411.9	74.1	15.5%	MPI_Isend
4.9%	300.2	95.8	24.6%	MPI_Allreduce
5.9%	365.4	44.6	11.1%	STRING
5.9%	365.4	44.6	11.1%	memcpy

# Identify High Time Consuming Areas

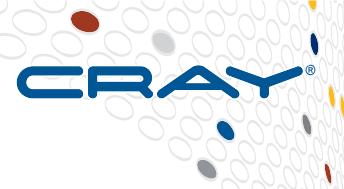


Table 2: Profile by Group, Function, and Line

Samp%	Samp	Imb.	Imb.	Group
		Samp	Samp%	Function
				Source
				Line
				PE=HIDE
100.0%	55,605.7	--	--	Total
<hr/>				
56.5%	31,412.8	--	--	USER
<hr/>				
19.7%	10,944.1	--	--	create_boundary\$boundary_
3				source.omp_removed/test/compile/boundary.f90
<hr/>				
4    7.8%	4,355.9	175.1	3.9%	line.265
4    1.8%	977.0	98.0	9.1%	line.268
4    1.1%	617.0	94.0	13.2%	line.273
4    2.0%	1,133.6	101.4	8.2%	line.549
4    1.1%	590.0	66.0	10.1%	line.557
<hr/>				
10.7%	5,937.8	--	--	get_block\$blocks_
3				source.omp_removed/test/compile/blocks.f90
<hr/>				
4    4.1%	2,305.7	145.3	5.9%	line.221
4    1.0%	569.5	77.5	12.0%	line.243
4    2.9%	1,610.1	134.9	7.7%	line.246

COMPUTE

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ANALYZE

# Guidance: How Can I Learn More?



MPI utilization:

The time spent processing MPI communications is relatively high. Functions and callsites responsible for consuming the most time can be found in the table generated by `pat_report -O callers+src` (within the MPI group).

# Guidance: Reduce Shared Resource Contention



## Metric-Based Rank Order:

When the use of a shared resource like memory bandwidth is unbalanced across nodes, total execution time may be reduced with a rank order that improves the balance.

A file named `MPICH_RANK_ORDER.USER_Time` was generated along with this report and contains usage instructions and the custom rank order from the following table.

Rank Order	Node Metric	Reduction in Imb.	Maximum Value	Average Value
Current	15.46%		1.134e+03	9.588e+02
Custom	1.46%	<b>14.202%</b>	9.731e+02	9.588e+02



# More In-depth Analysis and Bottleneck Detection

---

C O M P U T E

| S T O R E

| A N A L Y Z E

# Types of Perftools pat\_builds

## ● Sampling

- module load perftools-base perftools
- make <executable>
- pat\_build -O apa <executable>
  - Produces <executable>+pat
- Execute <executable>+pat
  - Produces <filename>s.xf
- pat\_report <filename>s.xf > profile\_samp
  - Produces <filename>.apa

## ● Trace

- pat\_build -O <filename>.apa
  - Produces <executable>+apa
- Execute <executable>+apa
  - Produces <filename>t.xf
- pat\_report <filename>t.xf > profile\_apa

# How to Use CrayPat

- **Make sure the following modules are loaded:**

- \$ module load perftools-base perftools

- **2 instrumentation examples:**

- \$ pat\_build my\_program
  - \$ pat\_build -u -g mpi my\_program

Same sampling experiment is used when perftools-lite module is loaded

- **Run application**

- \$ aprun -n ... my\_program+pat

- **Create report**

- \$ pat\_report my\_program.xf > my\_report

# Predefined Trace Wrappers (-g tracegroup)



- blas Basic Linear Algebra subprograms
  - caf Co-Array Fortran (Cray CCE compiler only)
  - hdf5 manages extremely large data collection
  - heap dynamic heap
  - io includes stdio and sysio groups
  - lapack Linear Algebra Package
  - math ANSI math
  - mpi MPI
  - omp OpenMP API
  - pthreads POSIX threads
  - shmem SHMEM
  - sysio I/O system calls
  - system system calls
  - upc Unified Parallel C (Cray CCE compiler only)
- Some may not work  
with Intel Compiler

For a full list, please see **pat\_build(1)** man page

# Control Data Collection with Runtime Options



- Runtime controlled through **PAT\_RT\_XXX** environment variables
- See **intro\_cravpat(1)** man page
- Examples of control
  - Enable full trace
  - Change number of data files created
  - Enable collection of CPU, network or power counter events
  - Enable tracing filters to control trace file size (max threads, max call stack depth, etc.)

# Performance Counters Overview



- Cray supports raw counters, derived metrics and thresholds for:
  - Processor (core and uncore)
  - Network
  - Accelerator
  - Power
- Predefined groups
  - Groups together counters for experiments
- See *hwpc*, *nwpc*, *accpc*, and *rapl* man pages

# How to Get List of Events for a Processor



- Run the following utility on a compute node:
  - `papi_native_avail`
- To collect performance counters
  - Set `PAT_RT_PERFCTR` environment variable to list of events or group prior to execution

# MPI Sync Time



- Measure load imbalance in programs instrumented to trace MPI functions to determine if MPI ranks arrive at collectives together
- Separates potential load imbalance from data transfer
- Sync times reported by default if MPI functions traced  
(pat\_build -g mpi)

```
Module load perftools-base  
Module load perftools  
make clean  
make  
pat_build -O apa exe  
(sampling experiment)  
aprun exe+pat  
pat_report <file>.xt
```



Next chart

Computation  
Highest

Communication  
Highest

Synchronization  
Highest

A  
Examine sampling  
exp with line #  
Table 3

B  
Examine  
trace exp  
with  
pat\_build -O <file.apa>  
Table 5

C  
Examine  
sampling exp  
for high % Imb sample

# If it doesn't Vectorize – fix it

Examine sampling  
exp with line #  
Table 3

```

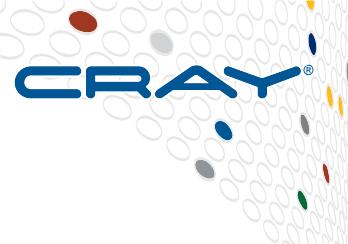
63. + 1----< do l = lmin, lmax
64. + 1 2--< do n = 1, 12
65. 1 2      pmold(l) = pmid(l)
66. 1 2      wlft (l) = 1.0 + gamfac1*(pmid(l) - plft(l)) * plfti(l)
67. 1 2      wrgh (l) = 1.0 + gamfac1*(pmid(l) - prgh(l)) * prghi(l)
68. 1 2      wlft (l) = clft(l) * sqrt(wlft(l))
69. 1 2      wrgh (l) = crgh(l) * sqrt(wrgh(l))
70. 1 2      zlft (l) = 4.0 * vlf(l) * wlft(l) * wlft(l)
71. 1 2      zrgh (l) = 4.0 * vrgh(l) * wrgh(l) * wrgh(l)
72. 1 2      zlft (l) = -zlft(l) * wlft(l)/(zlft(l) - gamfac2*(pmid(l) - plft(l)))
73. 1 2      zrgh (l) = zrgh(l) * wrgh(l)/(zrgh(l) - gamfac2*(pmid(l) - prgh(l)))
74. 1 2      umidl(l) = ulft(l) - (pmid(l) - plft(l)) / wlft(l)
75. 1 2      umidr(l) = urgh(l) + (pmid(l) - prgh(l)) / wrgh(l)
76. 1 2      pmid (l) = pmid(l) + (umidr(l) - umidl(l))*(zlft(l) * zrgh(l)) / (zrgh(l) - zlft(l))
77. 1 2      pmid (l) = max(smallp,pmid(l))
78. 1 2      if (abs(pmid(l)-pmold(l))/pmid(l) < tol ) exit
79. 1 2--> enddo
80. 1----> enddo

```

A	====						
	10.8%	376.9	--	--	riemann_		
3					riemann.f90		
	-----						
4	1.4%	47.4	32.6	41.0%	line.77		
4	3.9%	135.8	28.2	17.3%	line.78		

ftn-6254 ftn: VECTOR RIEMANN, File = riemann.f90, Line = 64

A loop starting at line 64 was not vectorized because a recurrence was found  
on "pmid" at line 77.



# If it doesn't Vectorize – fix it

```
63. + 1----< do l = lmin, lmax
64. + 1 2--< do n = 1, 12
65. 1 2      pmold(l) = pmid(l)
66. 1 2      wlft (l) = 1.0 + gamfac1*(pmid(l) - plft(l)) * plfti(l)
67. 1 2      wrgh (l) = 1.0 + gamfac1*(pmid(l) - prgh(l)) * prghi(l)
68. 1 2      wlft (l) = clft(l) * sqrt(wlft(l))
69. 1 2      wrgh (l) = crgh(l) * sqrt(wrgh(l))
70. 1 2      zlft (l) = 4.0 * vlft(l) * wlft(l) * wlft(l)
71. 1 2      zrgh (l) = 4.0 * vrgh(l) * wrgh(l) * wrgh(l)
72. 1 2      zlft (l) = -zlft(l) * wlft(l)/(zlft(l) - gamfac2*(pmid(l) - plft(l)))
73. 1 2      zrgh (l) = zrgh(l) * wrgh(l)/(zrgh(l) - gamfac2*(pmid(l) -
                                         prgh(l)))
74. 1 2      umidl(l) = ulft(l) - (pmid(l) - plft(l)) / wlft(l)
75. 1 2      umidr(l) = urgh(l) + (pmid(l) - prgh(l)) / wrgh(l)
76. 1 2      pmid (l) = pmid(l) + (umidr(l) - umidl(l))*(zlft(l) * zrgh(l)) /
                           (zrgh(l) - zlft(l))
77. 1 2      pmid (l) = max(smallp,pmid(l))
78. 1 2      if (abs(pmid(l)-pmold(l))/pmid(l) < tol ) exit
79. 1 2--> enddo
80. 1----> enddo

62. A----> converged =.F.
63. + 1----< do n = 1, 12
64. 1 Vr2--< do l = lmin, lmax
65. 1 Vr2      if(.not.converged(l))then
66. 1 Vr2          pmold(l) = pmid(l)
67. 1 Vr2          wlft (l) = 1.0 + gamfac1*(pmid(l) - plft(l)) * plfti(l)
68. 1 Vr2          wrgh (l) = 1.0 + gamfac1*(pmid(l) - prgh(l)) * prghi(l)
69. 1 Vr2          wlft (l) = clft(l) * sqrt(wlft(l))
70. 1 Vr2          wrgh (l) = crgh(l) * sqrt(wrgh(l))
71. 1 Vr2          zlft (l) = 4.0 * vlft(l) * wlft(l) * wlft(l)
72. 1 Vr2          zrgh (l) = 4.0 * vrgh(l) * wrgh(l) * wrgh(l)
73. 1 Vr2          zlft (l) = -zlft(l) * wlft(l)/(zlft(l) - gamfac2*(pmid(l) - plft(l)))
74. 1 Vr2          zrgh (l) = zrgh(l) * wrgh(l)/(zrgh(l) - gamfac2*(pmid(l) - prgh(l)))
75. 1 Vr2          umidl(l) = ulft(l) - (pmid(l) - plft(l)) / wlft(l)
76. 1 Vr2          umidr(l) = urgh(l) + (pmid(l) - prgh(l)) / wrgh(l)
77. 1 Vr2          pmid (l) = pmid(l) + (umidr(l) - umidl(l))*(zlft(l) * zrgh(l)) / &
                               (zrgh(l)-zlft(l))
78. 1 Vr2          pmid (l) = max(smallp,pmid(l))
79. 1 Vr2          if (abs(pmid(l)-pmold(l))/pmid(l) < tol ) then
80. 1 Vr2              converged(l) = .T.
81. 1 Vr2              endif
82. 1 Vr2          endif
83. 1 Vr2      enddo
84. 1 Vr2--> enddo
85. + 1      if(all(converged(lmin:lmax)))exit
86. 1-----> enddo
```

# If it Vectorizes, check memory utilization

44.	V---< do n = nmin-1, nmax							
45.	V ar(n) = a(n) + para(n,1)*diffa(n) + para(n,2)*da(n+1) + para(n,3)*da(n)		4	0.1%	3.4	5.6	62.0%	line.44
48.	V al(n+1) = ar(n)		4	2.5%	88.5	27.5	23.8%	line.45
49.	V---> enddo		4	0.2%	7.1	8.9	56.0%	line.48
53.	fV---< do n = nmin, nmax		4	0.0%	0.1	1.9	94.9%	line.49
54.	fV onemfl= 1.0 - flat(n)		4	0.3%	9.6	10.4	52.3%	line.53
55.	fV ar(n) = flat(n) * a(n) + onemfl * ar(n)		4	0.3%	10.1	10.9	51.9%	line.54
56.	fV al(n) = flat(n) * a(n) + onemfl * al(n)		4	0.9%	31.3	14.7	32.0%	line.55
57.	fV---> enddo		4	0.2%	7.3	11.7	61.6%	line.56
67.	f---< do n = nmin, nmax		4	0.0%	0.2	1.8	90.6%	line.57
68.	f deltaa(n) = ar(n) - al(n)		4	0.4%	13.6	10.4	43.5%	line.68
69.	f a6(n) = 6. * (a(n) - .5 * (al(n) + ar(n)))		4	0.8%	27.2	13.8	33.9%	line.69
70.	f scrch1(n) = (ar(n) - a(n)) * (a(n)-al(n))		4	0.4%	15.4	11.6	43.2%	line.70
71.	f scrch2(n) = deltaa(n) * deltaa(n)		4	0.4%	13.4	10.6	44.3%	line.71
72.	f scrch3(n) = deltaa(n) * a6(n)		4	0.4%	14.0	12.0	46.2%	line.72
73.	f---> enddo		4	0.1%	4.8	8.2	63.5%	line.75
74.			4	0.3%	10.2	9.8	49.1%	line.76
75.	Vr2--< do n = nmin, nmax		4	0.6%	19.4	14.6	43.1%	line.77
76.	Vr2 if(scrch1(n) <= 0.0) then		4	0.5%	17.9	11.1	38.4%	line.78
77.	Vr2 ar(n) = a(n)		4	1.2%	40.6	16.4	28.8%	line.80
78.	Vr2 al(n) = a(n)		4	1.8%	63.9	24.1	27.5%	line.81
79.	Vr2 endif		4	0.0%	0.2	2.8	92.7%	line.82
80.	Vr2 if(scrch2(n) < +scrch3(n)) al(n) = 3. * a(n) - 2. * ar(n)		4	0.1%	4.3	7.7	64.4%	line.84
81.	Vr2 if(scrch2(n) < -scrch3(n)) ar(n) = 3. * a(n) - 2. * al(n)		4	0.5%	16.0	12.0	43.2%	line.85
82.	Vr2--> enddo		4	0.6%	21.3	13.7	39.4%	line.86
83.			4	0.0%	0.0	1.0	95.7%	line.87
84.	Vr2--< do n = nmin, nmax							
85.	Vr2 deltaa(n)= ar(n) - al(n)							
86.	Vr2 a6(n) = 6. * (a(n) - .5 * (al(n) + ar(n)))							
87.	Vr2--> enddo							



# If it Vectorizes, check memory utilization

```
44. V---< do n = nmin-1, nmax
45. V      ar(n) = a(n) + para(n,1)*diffa(n) + para(n,2)*da(n+1) + para(n,3)*da(n)
46. V      al(n+1) = ar(n)
47. V---> enddo
48. fV---< do n = nmin, nmax
49. fV      onemfl= 1.0 - flat(n)
50. fV      ar(n) = flat(n) * a(n) + onemfl * ar(n)
51. fV      al(n) = flat(n) * a(n) + onemfl * al(n)
52. fV---> enddo
53. f----< do n = nmin, nmax
54. f      deltaa(n) = ar(n) - al(n)
55. f      a6(n) = 6. * (a(n) - .5 * (al(n) + ar(n)))
56. f      scrch1(n) = (ar(n) - a(n)) * (a(n)-al(n))
57. f      scrch2(n) = deltaa(n) * deltaa(n)
58. f      scrch3(n) = deltaa(n) * a6(n)
59. f----> enddo
60.
61. Vr2--< do n = nmin, nmax
62. Vr2      if(scrch1(n) <= 0.0) then
63. Vr2          ar(n) = a(n)
64. Vr2          al(n) = a(n)
65. Vr2      endif
66. Vr2      if(scrch2(n) < +scrch3(n)) al(n) = 3. * a(n) - 2. * ar(n)
67. Vr2      if(scrch2(n) < -scrch3(n)) ar(n) = 3. * a(n) - 2. * al(n)
68. Vr2--> enddo
69.
70. Vr2--< do n = nmin, nmax
71. Vr2      deltaa(n)= ar(n) - al(n)
72. Vr2      a6(n) = 6. * (a(n) - .5 * (al(n) + ar(n)))
73. Vr2--> enddo
```

Remove unnecessary arrays  
and fuse loops

```
52. Vr2--< do n = nmin, nmax
53. Vr2      onemfl= 1.0 - flat(n)
54. Vr2      ar(n) = flat(n) * a(n) + onemfl * ar(n)
55. Vr2      al(n) = flat(n) * a(n) + onemfl * al(n)
56. Vr2      deltaa(n) = ar(n) - al(n)
57. Vr2      a6(n) = 6. * (a(n) - .5 * (al(n) + ar(n)))
58. Vr2      scrch1s = (ar(n) - a(n)) * (a(n)-al(n))
59. Vr2      scrch2s = deltaa(n) * deltaa(n)
60. Vr2      scrch3s = deltaa(n) * a6(n)
61. Vr2      if(scrch1s <= 0.0) then
62. Vr2          ar(n) = a(n)
63. Vr2          al(n) = a(n)
64. Vr2      endif
65. Vr2      if(scrch2s < +scrch3s) al(n) = 3. * a(n) - 2. * ar(n)
66. Vr2      if(scrch2s < -scrch3s) ar(n) = 3. * a(n) - 2. * al(n)
67. Vr2      deltaa(n)= ar(n) - al(n)
68. Vr2      a6(n) = 6. * (a(n) - .5 * (al(n) + ar(n)))
69. Vr2--> enddo
```

Examine  
trace exp  
with  
`pat_build -O <file.apa>`  
Table 5

Table 5: MPI Message Stats by Caller (limited entries shown)

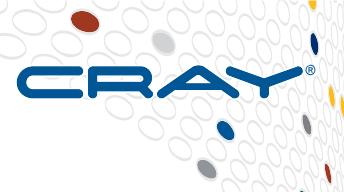
MPI	MPI Msg	Bytes	MPI Msg	MsgSz	64KiB<=	Function
Msg		Count	<16	MsgSz	Caller	
Bytes%		Count	<1MiB	PE=[mmm]		
		Count				
100.0%	3,774,873,804.0	4,851.0	51.0	4,800.0	Total	
100.0%	3,774,873,600.0	4,800.0	0.0	4,800.0	<code>mpi_alltoall_</code>	
66.7%	2,516,582,400.0	3,200.0	0.0	3,200.0	<code>sweezy_</code>	
4	66.7%	2,516,582,400.0	3,200.0	0.0	3,200.0	<code>pe.0</code>
4	66.7%	2,516,582,400.0	3,200.0	0.0	3,200.0	<code>pe.128</code>
4	66.7%	2,516,582,400.0	3,200.0	0.0	3,200.0	<code>pe.255</code>
	33.3%	1,258,291,200.0	1,600.0	0.0	1,600.0	<code>sweepz_</code>
4	33.3%	1,258,291,200.0	1,600.0	0.0	1,600.0	<code>pe.0</code>
4	33.3%	1,258,291,200.0	1,600.0	0.0	1,600.0	<code>pe.128</code>
4	33.3%	1,258,291,200.0	1,600.0	0.0	1,600.0	<code>pe.255</code>

B

====	15.5%	9.610204	--   --   361.0	MPI
====	15.5%	9.607925	0.229909   2.3%	300.0  mpi_alltoall_

22. + call MPI\_ALLTOALL(send1,Ya2abuff\_size,VH1\_DATATYPE,recv1,Ya2abuff\_size,  
VH1\_DATATYPE,MPI\_COMM\_ROW,mpierr)

# If Computation is load-imbalance – make sure computation is optimized



```
trace exp  
with  
pat_build -O <file.apa>  
Table 2
```

C	10.7%	6.636565	--	--	353.0	MPI_SYNC
	4.7%	2.888334	2.880515	99.7%	300.0	mpi_alltoall_(sync)
	4.2%	2.605623	2.605590	100.0%	1.0	mpi_init_(sync)
	1.8%	1.120161	1.111971	99.3%	51.0	mpi_allreduce_(sync)

	41.8%	5.536452	2.880515	99.7%	mpi_alltoall_(sync)
	41.8%	5.536452	--	--	pe.40
	0.1%	0.007819	--	--	pe.96

Table 1: Profile by Function Group and Function



Time%	Time	Imb.	Imb.	Calls	Group
		Time	Time%		Function
					PE=HIDE
100.0%	61.975194	--	--	4,916,166.0	Total
73.8%	45.728391	--	--	4,915,451.0	USER
19.4%	12.048791	1.192879	9.0%	307,200.0	remap_
12.8%	7.930665	1.337482	14.5%	2,764,800.0	parabola_
9.7%	6.006924	0.356937	5.6%	50.0	sweepz_
9.5%	5.884056	0.335782	5.4%	100.0	sweeppy_
6.4%	3.976012	0.495753	11.1%	307,200.0	riemann_
4.1%	2.530338	0.188650	7.0%	50.0	sweepx2_
4.0%	2.486642	0.200569	7.5%	50.0	sweepx1_
2.5%	1.523841	0.163506	9.7%	307,200.0	evolve_
2.4%	1.497595	0.305383	17.0%	614,400.0	paraset_
1.5%	0.923626	0.159765	14.8%	307,200.0	flatten_
1.4%	0.870671	0.157344	15.4%	307,200.0	states_

## Load Imbalance on unoptimized routines



Table 1: Profile by Function Group and Function

## Load Imbalance on optimized routines



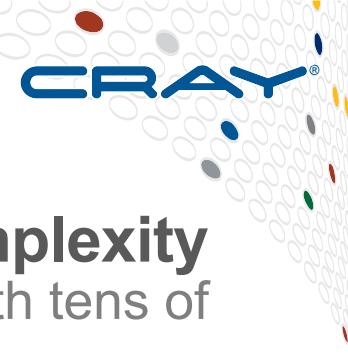
Time%	Time	Imb.	Imb.	Calls	Group
		Time	Time%		Function
					PE=HIDE
100.0%	58.936459	--	--	4,916,166.0	Total
72.3%	42.617036	--	--	4,915,451.0	USER
20.4%	12.023155	1.106771	8.5%	307,200.0	remap_
11.0%	6.476465	1.210511	15.8%	2,764,800.0	parabola_
10.1%	5.981124	0.297341	4.8%	50.0	sweepz_
10.0%	5.887939	0.295370	4.8%	100.0	sweeppy_
4.3%	2.532612	0.159728	6.0%	50.0	sweepx2_
4.2%	2.487010	0.172008	6.5%	50.0	sweepx1_
4.0%	2.360784	1.441720	38.1%	307,200.0	riemann_
2.6%	1.526218	0.139174	8.4%	307,200.0	evolve_
2.5%	1.501093	0.268596	15.2%	614,400.0	paraset_
1.6%	0.917283	0.142748	13.5%	307,200.0	flatten_
1.5%	0.871589	0.141770	14.0%	307,200.0	states_

COMPUTE

STORE

ANALYZE

# Motivation for Load Imbalance Analysis



- **Increasing system software and architecture complexity**
  - Current trend in high end computing is to have systems with tens of thousands of processors
    - This is being accentuated with multi-core processors
- **Applications have to be very well balanced In order to perform at scale on these MPP systems**
  - Efficient application scaling includes a balanced use of requested computing resources
- **Desire to minimize computing resource “waste”**
  - Identify slower paths through code
  - Identify inefficient “stalls” within an application

# Find Program Load Imbalance



Table 1: Profile by Function Group and Function

Time%	Time	Imb.	Imb.	Calls	Group
		Time	Time%		Function
					PE=HIDE
100.0%	1.957703	--	--	42,970.8	Total
60.0%	1.174021	--	--	3,602.0	USER
30.8%	0.603850	0.176924	23.0%	1,198.0	calc3_
19.2%	0.375117	0.128748	26.0%	1,200.0	calc2_
9.1%	0.178111	0.081880	32.0%	1,200.0	calc1_
36.0%	0.704928	--	--	9,613.0	MPI_SYNC
25.8%	0.505174	0.385130	76.2%	9,596.0	mpi_barrier_(sync)
10.2%	0.199537	0.199518	100.0%	1.0	mpi_init_(sync)
4.0%	0.078736	--	--	29,754.8	MPI
2.3%	0.045351	0.003531	7.3%	9,596.0	MPI_BARRIER
1.1%	0.021520	0.051295	71.6%	8,756.9	MPI_ISEND

# Sort MPI Messages by Caller



MPI	MPI Msg	MPI Msg	MsgSz	4KiB<=	Function
Msg	Bytes	Count	<16	MsgSz	Caller
Bytes%			Count	<64KiB	PE=[mmm]
100.0%	34,940,767.4	8,771.9	258.6	8,513.3	Total
100.0%	34,940,647.4	8,756.9	243.6	8,513.3	MPI_ISEND
56.2%	19,622,700.0	4,837.5	56.2	4,781.2	calc2_
3					shallow_
4	56.4%	19,718,400.0	7,200.0	2,400.0	4,800.0   pe.0
4	56.4%	19,699,200.0	4,800.0	0.0	4,800.0   pe.32
4	42.3%	14,784,000.0	4,800.0	1,200.0	3,600.0   pe.63
4	42.5%	14,851,950.0	3,693.8	75.0	3,618.8   calc1_
3					shallow_
4	56.4%	19,718,400.0	7,200.0	2,400.0	4,800.0   pe.0
4	42.3%	14,774,400.0	3,600.0	0.0	3,600.0   pe.31
4	42.3%	14,774,400.0	3,600.0	0.0	3,600.0   pe.62

COMPUTE

STORE

ANALYZE

# Analyze MPI Message Sizes



## Total

MPI Msg Bytes%	100.0%
MPI Msg Bytes	4,465,684,125.8
MPI Msg Count	13,057.0 msgs
MsgSz <16 Count	719.0 msgs
16<= MsgSz <256 Count	28.0 msgs
256<= MsgSz <4KiB Count	0.7 msgs
4KiB<= MsgSz <64KiB Count	279.8 msgs
64KiB<= MsgSz <1MiB Count	12,029.6 msgs

## MPI\_Send

MPI Msg Bytes%	100.0%
MPI Msg Bytes	4,465,680,353.8
MPI Msg Count	12,318.0 msgs
MsgSz <16 Count	8.0 msgs
16<= MsgSz <256 Count	0.0 msgs
256<= MsgSz <4KiB Count	0.7 msgs
4KiB<= MsgSz <64KiB Count	279.8 msgs
64KiB<= MsgSz <1MiB Count	12,029.6 msgs

## MPI\_Send / LAMMPS\_NS::Comm::reverse\_comm

MPI Msg Bytes%	48.6%
MPI Msg Bytes	2,171,466,150.3
MPI Msg Count	6,006.0 msgs
MsgSz <16 Count	0.0 msgs
16<= MsgSz <256 Count	0.0 msgs
256<= MsgSz <4KiB Count	0.0 msgs
4KiB<= MsgSz <64KiB Count	0.0 msgs
64KiB<= MsgSz <1MiB Count	6,006.0 msgs

## MPI\_Send / LAMMPS\_NS::Comm::reverse\_comm / LAMMPS\_NS::Verlet::run

MPI Msg Bytes%	48.6%
MPI Msg Bytes	2,169,218,110.3
MPI Msg Count	6,000.0 msgs
MsgSz <16 Count	0.0 msgs
16<= MsgSz <256 Count	0.0 msgs
256<= MsgSz <4KiB Count	0.0 msgs
4KiB<= MsgSz <64KiB Count	0.0 msgs
64KiB<= MsgSz <1MiB Count	6,000.0 msgs



# Maximize On-node Communication by Reordering MPI ranks

# When Is Rank Re-ordering Useful?

- Maximize on-node communication between MPI ranks
- Physical system topology agnostic
- Grid detection and rank re-ordering is helpful for programs with significant point-to-point communication
- Relieve on-node shared resource contention by pairing threads or processes that perform different work on the same node
  - for example: computation with off-node communication

# MPI Rank Reorder – Two Interfaces Available



- **CrayPat**

- Available with sampling or tracing
- Include `-g mpi` when instrumenting program
- Run program and let CrayPat determine if communication is dominant, detect communication pattern and suggest MPI rank order if applicable

- **grid\_order utility**

- User knows communication pattern in application and wants to quickly create a new MPI rank placement file
- Available when perftools-base module is loaded

# MPI Rank Order Observations



Table 1: Profile by Function Group and Function

Time%	Time	Imb. Time	Imb. Time%	Calls	Group Function PE=HIDE
100.0%	463.147240	--	--	21621.0	Total
<b>52.0%</b>	<b>240.974379</b>	--	--	21523.0	<b>MPI</b>
47.7%	221.142266	36.214468	14.1%	10740.0	mpi_recv
4.3%	19.829001	25.849906	56.7%	10740.0	MPI_SEND
43.3%	200.474690	--	--	32.0	USER
41.0%	189.897060	58.716197	23.6%	12.0	sweep_
1.6%	7.579876	1.899097	20.1%	12.0	source_
4.7%	21.698147	--	--	39.0	MPI_SYNC
4.3%	20.091165	20.005424	99.6%	32.0	mpi_allreduce_(sync)
0.0%	0.000024	--	--	27.0	SYSCALL

# MPI Rank Order Observations (2)



## MPI Grid Detection:

There appears to be point-to-point MPI communication in a 96 X 8 grid pattern. The 52% of the total execution time spent in MPI functions might be reduced with a rank order that maximizes communication between ranks on the same node. The effect of several rank orders is estimated below.

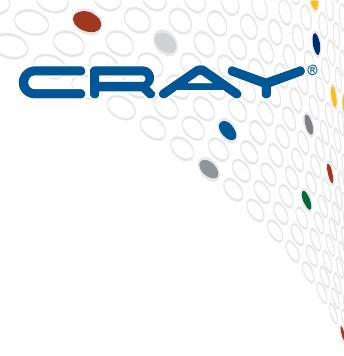
A file named MPICH\_RANK\_ORDER.Grid was generated along with this report and contains usage instructions and the Custom rank order from the following table.

Rank Order	On-Node Bytes/PE	On-Node Bytes/PE% of Total	MPICH_RANK_REORDER_METHOD
Custom	2.385e+09	95.55%	3
SMP	1.880e+09	75.30%	1
Fold	1.373e+06	0.06%	2
RoundRobin	0.000e+00	0.00%	0

# Auto-Generated MPI Rank Order File



# Using New MPI Rank Order



- Save grid\_order output to file called **MPICH\_RANK\_ORDER**
- `$ export MPICH_RANK_REORDER_METHOD=3`
- Run non-instrumented binary with and without new rank order to check overall wallclock time for performance improvements
- Can be used for all subsequent executions of **same job size**

# Summary



- Users continue to need tools to help find critical performance bottlenecks within a program
- Cray performance tools offer functionality that reduces the time investment associated with porting and tuning applications on new and existing Cray systems

# If Computation is load-imbalance – make sure to use the best MPI task mapping to nodes.



## Metric-Based Rank Order:

When the use of a shared resource like memory bandwidth is unbalanced across nodes, total execution time may be reduced with a rank order that improves the balance. The metric used here for resource usage is: USER Samp

For each node, the metric values for the ranks on that node are summed. The maximum and average value of those sums are shown below for both the current rank order and a custom rank order that seeks to reduce the maximum value.

A file named MPICH\_RANK\_ORDER.USER\_Samp was generated along with this report and contains usage instructions and the Custom rank order from the following table.

Rank Order	Node Metric	Reduction in Imb.	Maximum Value	Average Value
Current	3.56%		3.008e+06	2.901e+06
Custom	0.01%	3.557%	2.901e+06	2.901e+06

# If Communication is load-imbalance – make sure to use the best MPI task mapping to nodes.



## MPI Grid Detection:

There appears to be point-to-point MPI communication in a 8 X 8 X 8 grid pattern. The 16.4% of the total execution time spent in MPI functions might be reduced with a rank order that maximizes communication between ranks on the same node. The effect of several rank orders is estimated below.

A file named MPICH\_RANK\_ORDER.Grid was generated along with this report and contains usage instructions and the Custom rank order from the following table.

Rank Order	On-Node Bytes/PE	On-Node Bytes/PE% of Total	MPICH_RANK_REORDER_METHOD Bytes/PE
Custom	3.396e+12	74.82%	3
RoundRobin	3.026e+12	66.67%	0
SMP	2.680e+12	59.04%	1
Fold	1.700e+12	37.45%	2

# A pointer for looking at load imbalance



- When running on a large number of MPI tasks, the load imbalance is averaged over all the tasks. This can result in an obscure situation where a seemingly small load imbalance can hide a performance bottleneck.
- For example, in a run of the NIF problem we had a bad load imbalance and the culprit was
- || 0.4% | 82.6 | 195.4 | 70.8% |**hypre\_BinarySearch**
  - This was run on 128 MPI tasks (2 nodes); however, while the average load imbalance was an average of 82.6 samples (.826 seconds) some of the MPI tasks took significantly longer. By optimizing hypre-BinarySearch we were able to get a nice performance boost.

# Memory Analysis Tool



- **First load modules**

- module load perftools-base perftools-lite-hbm
- make <executable>
- Execute <executable>
- After execution you will see a <filename>.rpt and/or information attached to stdout

Will not work  
with Intel Compiler

- **The memory analysis tool can be run on KNL or Haswell**

- Better information is obtained on Haswell due to better hardware counters
  - The tool understands that the target system is KNL, so Haswell hardware counters are used to identify issues on KNL



```
#####
#          #
#      CrayPat-lite Performance Statistics      #
#          #
#####
```

CrayPat/X: Version 6.4.6 Revision 7d0d87c 02/20/17 09:52:37

Experiment: lite lite/hbm

Number of PEs (MPI ranks): 512

Numbers of PEs per Node: 32 PEs on each of 16 Nodes

Numbers of Threads per PE: 1

Number of Cores per Socket: 18

Execution start time: Thu Feb 23 09:50:22 2017

System name and speed: kay 2101 MHz (approx)

Intel broadwell CPU Family: 6 Model: 79 Stepping: 1

Will not work  
with Intel Compiler

Avg Process Time: 246.96 secs

High Memory: 596,629.9 MBytes 1,165.3 MBytes per PE

Sampling overflow event and threshold:

MEM\_LOAD\_UOPS\_RETIRED: HIT\_LFB:precise=2 400,009

Table 1: HBM Objects Sorted by Weighted Sample



Will not work  
with Intel Compiler

Object Sample	Object Sample	Object Max	Object Max Size	Object PE=HIDE	Object Location
Weight%	Count%	Count	Active	(MBytes)	
100.0%	100.0%	25,751.0	581.4	1,611.566	Total
<hr/>					
12.8%	7.2%	1,841.6	1.0	43.875	yspecies@allocate@../source/f90_files/solve/init_field.f90@line.83
11.5%	9.0%	2,320.0	1.0	141.750	q@allocate@../source/f90_files/solve/init_field.f90@line.83
11.2%	18.9%	4,854.1	1.0	0.276	\$\$_rf@local@<unknown>@line.0
10.0%	8.2%	2,107.5	1.0	131.625	diffflux@local@../source/f90_files/solve/rhsf.f90@line.101
9.9%	16.5%	4,260.7	1.0	0.276	\$\$_rb@local@<unknown>@line.0
8.3%	3.1%	804.8	1.0	131.625	grad_ys@local@../source/f90_files/solve/rhsf.f90@line.101
6.4%	2.6%	666.6	1.0	0.844	phi@allocate@../source/modules/mixfrac_m.f90@line.166
5.9%	2.8%	723.8	1.0	141.750	tmmp@local@../source/f90_files/solve/rhsf.f90@line.89
4.2%	6.8%	1,754.1	1.0	2.531	u@allocate@../source/f90_files/solve/init_field.f90@line.83
2.2%	3.8%	967.4	1.0	0.051	\$\$_dif@local@<unknown>@line.0
2.0%	3.3%	848.9	1.0	0.051	\$\$_c@local@<unknown>@line.0
1.4%	0.7%	186.2	1.0	47.250	q_err@allocate@../source/modules/erk_m.f90@line.463
1.2%	2.0%	514.7	1.0	0.051	\$\$_x@local@<unknown>@line.0
1.1%	0.4%	103.3	1.0	43.875	h_spec@local@../source/f90_files/solve/rhsf.f90@line.111
<hr/>					



Table 2: HBM Objects Sorted by Sample Count

Will not work  
with Intel Compiler

Object	Object	Object	Object	Object	Object Location
Sample	Sample	Sample	Max	Max Size	PE=HIDE
Count%	Weight%	Count	Active	(MBytes)	
100.0%	100.0%	25,751.0	581.4	1,611.566	Total
<hr/>					
18.9%	11.2%	4,854.1	1.0	0.276	\$\$_rf@local@<unknown>@line.0
16.5%	9.9%	4,260.7	1.0	0.276	\$\$_rb@local@<unknown>@line.0
9.0%	11.5%	2,320.0	1.0	141.750	q@allocate@../source/f90_files/solve/init_field.f90@line.83
8.2%	10.0%	2,107.5	1.0	131.625	diffflux@local@../source/f90_files/solve/rhsf.f90@line.101
7.2%	12.8%	1,841.6	1.0	43.875	yspecies@allocate@../source/f90_files/solve/init_field.f90@line.83
6.8%	4.2%	1,754.1	1.0	2.531	u@allocate@../source/f90_files/solve/init_field.f90@line.83
3.8%	2.2%	967.4	1.0	0.051	\$\$_dif@local@<unknown>@line.0
3.3%	2.0%	848.9	1.0	0.051	\$\$_c@local@<unknown>@line.0
3.1%	8.3%	804.8	1.0	131.625	grad_ys@local@../source/f90_files/solve/rhsf.f90@line.101
2.8%	5.9%	723.8	1.0	141.750	tmmpp@local@../source/f90_files/solve/rhsf.f90@line.89
2.6%	6.4%	666.6	1.0	0.844	phi@allocate@../source/modules/mixfrac_m.f90@line.166
2.1%	0.7%	549.4	1.0	0.021	sqrtq@local@../source/modules/computeCoefficients_r.f90@line.122
2.0%	1.2%	514.7	1.0	0.051	\$\$_x@local@<unknown>@line.0
1.6%	0.9%	403.0	1.0	0.051	\$\$_y@local@<unknown>@line.0
1.3%	0.7%	334.3	1.0	0.051	\$\$_xs@local@<unknown>@line.0
1.1%	0.3%	283.6	1.0	0.051	\$\$_rmcwrlk1@local@<unknown>@line.0
1.0%	0.9%	268.9	1.0	43.875	diffusion@local@../source/f90_files/solve/rhsf.f90@line.147
1.0%	0.6%	246.3	1.0	0.051	\$\$_xxwt@local@<unknown>@line.0
<hr/>					



Table 4: Profile by Function

Samp%	Samp	Imb.	Imb.	MEM_LOAD_UOPS_RETIRED	RESOURCE_STALLS	Group
		Samp	Samp%	:HIT_LFB:precise=2	:ALL	Function=[MAX10]
						PE=HIDE
100.0%	26,420.2	--	--	10,569,018,266	407,403,254,579	Total
-----						
96.9%	25,594.0	--	--	10,238,474,892	385,090,314,646	USER
45.7%	12,078.5	309.5	2.5%	4,831,543,082	111,983,149,693	reaction_rate_vec_.LOOP@li.347
13.1%	3,460.0	430.0	11.1%	1,384,031,140	24,253,285,440	rhsf_.LOOP@li.924
6.1%	1,610.0	116.0	6.7%	644,008,240	14,874,097,115	computecoefficients_r_.LOOP@li.277
4.9%	1,291.7	70.3	5.2%	516,699,907	6,342,301,352	computecoefficients_r_.LOOP@li.157
4.2%	1,115.9	33.1	2.9%	446,392,856	46,282,065,332	rhsf_.LOOP@li.626
4.0%	1,064.6	152.4	12.6%	425,838,487	30,080,185,807	derivative_x_bnds_buff_r_
3.5%	920.1	18.9	2.0%	368,067,656	24,815,310,694	calc_primary_vars_.LOOP@li.52
3.0%	793.4	17.6	2.2%	317,354,015	27,640,713,430	integrate_.LOOP@li.121
2.2%	582.5	58.5	9.2%	232,991,961	12,088,490,382	rhsf_.LOOP@li.1742
=====						
3.1%	817.5	--	--	327,059,702	21,968,897,774	ETC
-----						
2.1%	544.6	124.4	18.6%	217,925,216	19,260,243,388	_cray_mpi_memcpy_snb
=====						

Will not work  
with Intel Compiler

Table 1: Profile by Function



Samp%	Samp	Imb.	Imb.	Group
		Samp	Samp%	Function
				PE=HIDE
100.0%	56,875.0	--	--	Total
<hr/>				
76.2%	43,356.4	--	--	USER
<hr/>				
23.4%	13,307.7	912.3	6.4%	reaction_rate_vec_.LOOP@li.349
9.2%	5,238.2	943.8	15.3%	rhsf_.LOOP@li.626
5.5%	3,131.3	212.7	6.4%	rhsf_.LOOP@li.1643
5.4%	3,045.9	99.1	3.2%	computecoefficients_r_.LOOP@li.277
5.3%	3,012.7	319.3	9.6%	rhsf_.LOOP@li.1687
4.6%	2,606.5	252.5	8.8%	calc_primary_vars_.LOOP@li.49
3.7%	2,110.6	1,192.4	36.2%	rhsf_.LOOP@li.924
2.5%	1,433.7	110.3	7.2%	computecoefficients_r_.LOOP@li.157
2.4%	1,375.5	402.5	22.7%	rhsf_.LOOP@li.454
2.2%	1,231.2	689.8	36.0%	integrate_.LOOP@li.121
2.1%	1,201.5	348.5	22.5%	derivative_x_bnds_buff_r_
1.5%	863.5	205.5	19.3%	rhsf_.LOOP@li.1742
<hr/>				
18.8%	10,696.5	--	--	MPI
<hr/>				
12.4%	7,049.7	3,691.3	34.4%	MPI_WAITANY
5.1%	2,876.8	3,602.2	55.7%	mpi_waitall
<hr/>				
4.9%	2,799.9	--	--	ETC
<hr/>				
3.2%	1,840.0	174.0	8.7%	cray_ALOG10
1.5%	874.3	90.7	9.4%	_EXP_Z

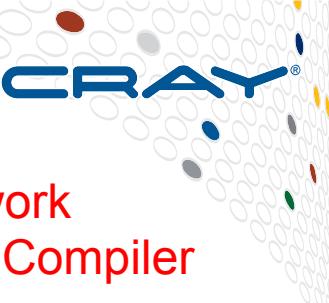
Will not work  
with Intel Compiler

Table 5: Profile by Group, Function, and Line



Will not work  
with Intel Compiler

Samp%	Samp	Imb.	Imb.	MEM_LOAD_UOPS_RETIRIED	RESOURCE_STALLS	Group
	Samp	Samp%		:HIT_LFB:precise=2	:ALL	Function=[MAX10]
						Source
						Line
						PE=HIDE
100.0%	26,420.2	--	--	10,569,018,266	407,403,254,579	Total
-----						
96.9%	25,594.0	--	--	10,238,474,892	385,090,314,646	USER
-----						
45.7%	12,078.5	--	--	4,831,543,082	111,983,149,693	reaction_rate_vec_.LOOP@li.347
3						getrates.f
-----						
4    1.0%	261.3	50.7	16.3%	104,503,914	2,421,756,549	line.486
4    1.1%	299.3	65.7	18.0%	119,716,756	2,733,588,335	line.3479
4    1.0%	274.5	155.5	36.2%	109,794,658	2,521,794,348	line.5001
=====						
13.1%	3,460.0	--	--	1,384,031,140	24,253,285,440	rhsf_.LOOP@li.924
3						rhsf.f90
-----						
4    5.5%	1,443.5	211.5	12.8%	577,416,117	10,080,883,003	line.933
4    3.7%	971.9	172.1	15.1%	388,757,966	6,818,669,249	line.938
4    4.0%	1,044.5	200.5	16.1%	417,807,057	7,352,846,349	line.943
=====						
6.1%	1,610.0	--	--	644,008,240	14,874,097,115	computecoefficients_r_.LOOP@li.277
3						computeCoefficients_r.f90
4  1.0%	271.2	54.8	16.9%	108,464,159	2,504,937,154	line.374
4.9%	1,291.7	--	--	516,699,907	6,342,301,352	computecoefficients_r_.LOOP@li.157



# Excerpt from getrates.f

Will not work  
with Intel Compiler

3478.	M m V	C	oh
3479.	M m V		DDOT=+RF(I,3)+RB(I,4)+RF(I,5)+RF(I,6)+RB(I,7)+RB(I,8)+RB(I,9)
3480.	M m V		*+RB(I,9)+RF(I,10)+RF(I,12)+RB(I,15)+RF(I,19)+RB(I,20)+RB(I,21)
3481.	M m V		*+RB(I,22)+RB(I,24)+RF(I,30)+RB(I,32)+RF(I,33)+RB(I,40)+RB(I,45)
3482.	M m V		*+RB(I,45)+RF(I,47)+RF(I,49)+RF(I,49)+RB(I,50)+RF(I,55)+RB(I,57)
3483.	M m V		*+RF(I,58)+RB(I,59)+RF(I,61)+RB(I,64)+RB(I,68)+RF(I,70)+RF(I,71)
3484.	M m V		*+RF(I,73)+RF(I,82)+RF(I,90)+RB(I,92)+RF(I,102)+RF(I,104)+RF(I,105)
3485.	M m V		*+RF(I,106)+RB(I,112)+RF(I,128)+RB(I,136)+RF(I,138)+RF(I,139)
3486.	M m V		*+RB(I,152)+RB(I,156)+RF(I,160)+RF(I,167)+RB(I,169)+RB(I,176)
3487.	M m V		*+RF(I,178)+RB(I,182)+RF(I,184)+RF(I,188)+RF(I,189)+RF(I,190)
3488.	M m V		*+RB(I,197)+RF(I,202)+RF(I,208)+RF(I,211)+RF(I,221)+RB(I,222)
3489.	M m V		*+RF(I,224)+RB(I,226)+RF(I,232)+RF(I,236)+RB(I,250)+RF(I,252)
3490.	M m V		*+RF(I,253)+RF(I,282)+RF(I,283)

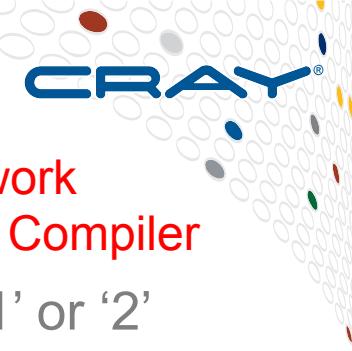


# Excerpt from rhsf.f90

```
924. + M m-----<    SPECIESX: do n=1,n_spec-1
925.     M m                         #ifdef GPU
926.     M D                         !$acc loop vector private(i, ml, mu,itmp)
927.     M m                         #endiff
928. + M m 2-----<    do i = 1, nx*ny*nz, ms
929.     M m 2                         ml = i
930.     M m 2                         mu = min(i+ms-1, nx*ny*nz)
931.     M m 2                         if (vary_in_x==1) then
932.     M m 2                         itmp = 15 + n
933. + M m 2 r4-----<>          tmmp(ml:mu,1,1,itmp) = -q(ml:mu,1,1,5+n)*u(ml:mu,1,1,1) - diffFlux(ml:mu,1,1,
934.     M m 2
935.     M m 2                         endif
936.     M m 2                         if (vary_in_y==1) then
937.     M m 2                         itmp = 15 + (n_spec-1) + n
938. + M m 2 r4-----<>          tmmp(ml:mu,1,1,itmp) = -q(ml:mu,1,1,5+n)*u(ml:mu,1,1,2) - diffFlux(ml:mu,1,1,
939.     M m 2
940.     M m 2                         endif
941.     M m 2                         if (vary_in_z==1) then
942.     M m 2                         itmp = 15 + 2*(n_spec-1) + n
943. + M m 2 r4-----<>          tmmp(ml:mu,1,1,itmp) = -q(ml:mu,1,1,5+n)*u(ml:mu,1,1,3) - diffFlux(ml:mu,1,1,
944.     M m 2                         endif
945.     M m 2----->      enddo
946.     M m----->      enddo SPECIESX
```

Will not work  
with Intel Compiler

# Tips When Running on KNL with Directives



- **Check for MCDRAM allocations** Will not work  
with Intel Compiler
  - Set CCE's CRAYMEM\_DEBUG environment variable to '1' or '2'
  - Look for posix\_memalign(allocator-"bandwidth" size=...) messages
- **Try running in split mode (25% cache) to reduce effect of TLB misses since there is no L3 cache**
  - If no performance gain over 100%, try flat
- **Don't allocate local variables within a routine to MCDRAM**
  - Convert to global first

# MCDRAM Configuration Information



Will not work  
with Intel Compiler

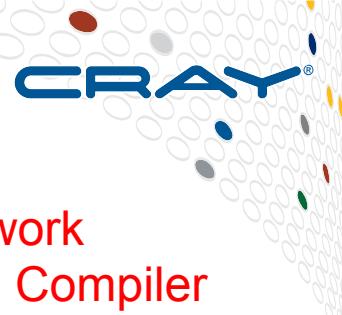
```
CrayPat/X: Version 6.4.2.36 Revision 8374f24 08/08/16 14:59:22
Experiment:           lite lite/sample_profile
Number of PEs (MPI ranks): 2,048
Numbers of PEs per Node:      32 PEs on each of 64 Nodes
Numbers of Threads per PE:    1
Number of Cores per Socket:   512 PEs on sockets with 34 Cores
                             1,536 PEs on sockets with 68 Cores
```

...

```
MCDRAM: 7.2 GHz, 16 GiB available as snc2, cache (100% cache) for 512 PEs
MCDRAM: 7.2 GHz, 16 GiB available as quad, cache (100% cache) for 1536 PEs
```

```
Avg Process Time:      3,251 secs
High Memory:          9,651,837.1 MBytes      4,712.8 MBytes per PE
I/O Read Rate:        23.645208 MBytes/sec
I/O Write Rate:       6.836539 MBytes/sec
Avg CPU Energy:       43,899,476 joules      685,929 joules per node
Avg CPU Power:        13,504 watts         211.00 watts per node
```

# Per-numanode Memory High Water Mark



- \$ module load perftools-lite Will not work  
with Intel Compiler
- Build and Run program:
  - 1) \$ aprun -m1800m -n 2048 **-N 16** -d 1 ./cpmd.x
  - 2) \$ aprun -m1800m -n 2048 **-N 16** -d 1 \  
**numactl --preferred=1** ./cpmd.x
  - 3) \$ aprun -m1800m -n 2048 **-N 64** -d 1 \  
**numactl --preferred=1** ./cpmd.x
- Run **pat\_report** to get memory high water mark breakdown
  - \$ **pat\_report -O himem** cpmd.x.\*ap2 > rpt

# 1) Flat Mode, No Allocation

MCDRAM: 7.2 GHz, 16 GiB available as quad, flat ( 0% cache)

Process	HiMem	Numanode
HiMem	Numa	PE=ALL
(MBytes)	Node 0	
		(MBytes)
827.3	827.3	Total
<hr/>		
827.3	827.3	numanode.0
<hr/>		
870.5	870.5	pe.1904
870.0	870.0	pe.2000
869.8	869.8	pe.1776
868.6	868.6	pe.1008
868.1	868.1	pe.1264

Will not work  
with Intel Compiler

16 ranks per core

No allocation to MCDRAM, only to  
DRAM so one Numanode is used

## 2) Flat Mode, Full Allocation

MCDRAM: 7.2 GHz, 16 GiB available as quad, flat ( 0% cache)

Will not work  
with Intel Compiler

Process	HiMem	HiMem	Numanode
HiMem	Numa	Numa	PE=ALL
(MBytes)	Node 0	Node 1	
	(MBytes)	(MBytes)	
827.7	0.0	827.7	Total
-----			
827.7	0.0	827.7	numanode.0
-----			
970.7	0.0	970.6	pe.18
970.4	0.0	970.4	pe.60
970.3	0.0	970.3	pe.242
970.2	0.0	970.2	pe.351
969.8	0.0	969.8	pe.562
870.1	0.0	870.1	pe.1904

16 ranks per core

With numactl --preferred=1, all allocations end up on MCDRAM Numanode 1. No allocations end up on DRAM Numanode 0.

### 3) Flat Mode, Partial Allocation

MCDRAM: 7.2 GHz, 16 GiB available as quad, flat ( 0% cache)

Process	HiMem	HiMem	Numanode
HiMem	Numa	Numa	PE=ALL
(MBytes)	Node 0	Node 1	
	(MBytes)	(MBytes)	
786.9	534.5	252.4	Total
786.9	534.5	252.4	numanode.0
794.3	538.9	255.4	pe.0
791.3	537.6	253.8	pe.64
791.2	537.3	253.9	pe.128
791.1	537.3	253.8	pe.192
791.0	537.1	253.9	pe.256
790.7	537.4	253.4	pe.136
790.7	538.0	252.8	pe.55

Will not work  
with Intel Compiler

Using 64 cores on a node creates less available space per MPI rank.

With numactl --preferred=1, after MCDRAM Numanode 1 is full, data is allocated to DRAM Numanode 0.

# MCDRAM Allocation Assistance Recap



Will not work  
with Intel Compiler

- Cray Tools track requests to memory and evaluate the bandwidth contribution of objects within a program
- Helpful for memory-intensive programs that cannot fit within MCDRAM
- Reduces time investment associated with selectively allocating data into KNL's MCDRAM
- The result is performance portable code
  - CCE memory allocation directives are ignored on X86 processors



Will not work  
with Intel Compiler

Cray Compiler Optimization Feedback

OpenMP Assistance

MCDRAM Allocation Assistance

# Reveal Overview



- Reduce effort associated with adding OpenMP to MPI programs
- Produce performance portable code
- Get insight into optimizations performed by the Cray compiler
- Use as a first step to parallelize loops that will target GPUs
- Track requests to memory and evaluate the bandwidth contribution of objects within a program

Will not work  
with Intel Compiler

# When to Move to a Hybrid Programming Model



- **When code is network bound**
  - Increased MPI collective and point-to-point wait times
- **When MPI starts leveling off**
  - Too much memory used, even if on-node shared communication is available
  - As the number of MPI ranks increases, more off-node communication can result, creating a network injection issue
- **When contention of shared resources increases**

# Approach to Adding Parallelism

## 1. Identify key high-level loops

- Determine where to add additional levels of parallelism

## 2. Perform parallel analysis and scoping

- Split loop work among threads

Will not work  
with Intel Compiler

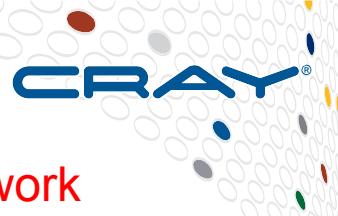
## 3. Add OpenMP layer of parallelism

- Insert OpenMP directives

## 4. Analyze performance for further optimization, specifically vectorization of innermost loops

- We want a performance-portable application at the end

# The Problem – How Do I Parallelize This Loop?



- How do I know this is a good loop to parallelize?
- What prevents me from parallelizing this loop?
- Can I get help building a directive?

Will not work  
with Intel Compiler

```
subroutine sweepz
...
do j = 1, js
  do i = 1, isz
    radius = zxc(i+mpey*isz)
    theta  = zyc(j+mpey*js)
    do m = 1, npez
      do k = 1, ks
        n = k + ks*(m-1) + 6
        r(n) = recv3(1,j,k,i,m)
        p(n) = recv3(2,j,k,i,m)
        u(n) = recv3(5,j,k,i,m)
        v(n) = recv3(3,j,k,i,m)
        w(n) = recv3(4,j,k,i,m)
        f(n) = recv3(6,j,k,i,m)
      enddo
    enddo
    ...
    call ppmlr
    do k = 1, kmax
      n = k + 6
      za (n) = zza(k)
      dz (n) = zdz(k)
      xa0(n) = zza(k)
      dx0(n) = zdz(k)
      e  (n) = p(n)/(r(n)*gamm)+0.5 &
                *(u(n)**2+v(n)**2+w(n)**2)
    enddo
    call ppmlr
...
  enddo
enddo
```

```
subroutine ppmlr

call boundary
call flatten
call paraset(nmin-4, nmax+5, para, dx, xa)

call parabola(nmin-4,nmax+4,para,p,dp,p6,pl,flat)
call parabola(nmin-4,nmax+4, para,r,dr,r6,rl,flat)
call parabola(nmin-4,nmax+4,para,u,du,u6,ul,flat)

call states(pl,ul,rl,p6,u6,r6,dp,du,dr,plft,ulft,&
            rlft,prgh,urgh,rrgh)
call riemann(nmin-3,nmax+4,gam,prgh,urgh,rrgh,&
            plft,ulft,rlft pmid umid)
call evolve(umid, pmid) ← contains more calls

call remap ← contains more calls

call volume(nmin,nmax,ngeom, radius,xa,dx,dvol)

call remap ← contains more calls

return
end
```

COMPUTE

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# Loop Work Estimates



*Gather loop statistics using the Cray performance tools and CCE to determine which loops have the most work*

- Helps identify high-level serial loops to parallelize
  - Based on runtime analysis, approximates how much work exists within a loop

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with Intel Compiler

# Collect Loop Work Estimates



- Point to Cray compiler
  - \$ module load PrgEnv-cray
- Set up `perftools` loop work estimates experiment
  - \$ module load `perftools-lite-loops`
- Build program (make)
- Run program to get loop work estimates in file with .ap2 suffix

# Example Loop Statistics

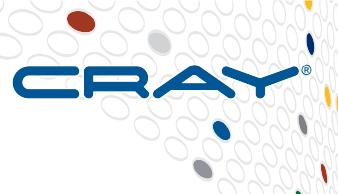


Table 2: Loop Stats by Function (from -hprofile\_generate)

Loop Incl Time	Loop Hit	Loop Trips Avg	Loop Trips Min	Loop Trips Max	Function=/.LOOP[.] PE=HIDE
Total					Will not work with Intel Compiler
8.995914	100	25	0	25	sweepy_.LOOP.1.li.33
8.995604	2500	25	0	25	sweepy_.LOOP.2.li.34
8.894750	50	25	0	25	sweepz_.LOOP.05.li.49
8.894637	1250	25	0	25	sweepz_.LOOP.06.li.50
4.420629	50	25	0	25	sweepx2_.LOOP.1.li.29
4.420536	1250	25	0	25	sweepx2_.LOOP.2.li.30
4.387534	50	25	0	25	sweepx1_.LOOP.1.li.29
4.387457	1250	25	0	25	sweepx1_.LOOP.2.li.30
2.523214	187500	107	0	107	riemann_.LOOP.2.li.63
1.541299	20062500	12	0	12	riemann_.LOOP.3.li.64
0.863656	1687500	104	0	108	parabola_.LOOP.6.li.67

# Create Program Library and Launch Reveal



- Disable loop work estimates program instrumentation
  - \$ module unload perftools-lite-loops
- Rebuild program with CCE's program library
  - Will not work with Intel Compiler
  - Add `-h pl=/full_path/program.pl` to program's Makefile
  - \$ make clean
  - \$ make
- Launch Reveal
  - \$ reveal /full\_path/program.pl loop\_estimates.ap2

# View Source and Optimization Information



Reveal

File Edit View Help

vphone.pl

Navigation

Loop Performance

- 4.0423 SWEEPX2@32
- 3.8576 SWEEPZ@51
- 3.8573 SWEEPZ@52
- 2.2068 RIEMANN@63
- 1.2299 RIEMANN@64
- 0.8068 PARABOLA@67
  - 0.0146 Instance #1
  - 0.0156 Instance #2
  - 0.0156 Instance #3
  - 0.0163 Instance #4
  - 0.0163 Instance #5
  - 0.0174 Instance #6
  - 0.0167 Instance #7

X Traceback

PARABOLA@67  
PPMLR@51  
sweepx1\_LOOP.2.li.32@53  
sweepx1\_LOOP.1.li.31@32  
SWEEPX1@31  
VPHONE@232

Source - /home/users/heidi/reveal/parabola.f90

Up

66

f 67 do n = nmin, nmax

68 deltaa(n) = ar(n) - al(n)

69 a6(n) = 6. \* (a(n) - .5 \* (al(n) + ar(n)))

70 scrch1(n) = (ar(n) - a(n)) \* (a(n)-al(n))

71 scrch2(n) = deltaa(n) \* deltaa(n)

72 scrch3(n) = deltaa(n) \* a6(n)

73 enddo

74

Vr2 75 do n = nmin, nmax

76 if(scrch1(n) <= 0.0) then

77 ar(n) = a(n)

78 al(n) = a(n)

79 endif

Loopmark Legend

- A Pattern Matched
- C Collapsed
- E Cloned
- G Accelerated
- I Inlined
- II Not Inlined
- L Loop
- M Multithreaded
- R Region
- S Scoping Analysis
- V Vectorized
- a Atomic Memory Operation
- b Blocked
- c Conditional and/or Computed
- f Fused
- g Partitioned
- i Interchanged
- n Non-blocking Remote Transfer
- p Partial
- r Unrolled
- s Shortloop
- w Unwound

Info - Line 67

A loop starting at line 67 was fused with the loop starting at line 53.

Will not work with Intel Compiler

vphone.pl loaded. vphone\_loops.ap2 loaded.

COMPUTE

STORE

ANALYZE

# Access Cray Compiler Message Information



File Edit View Help  
vhone.pi X

**Navigation**

- riemann.f90
- remap.f90
- evolve.f90
- volume.f90
- forces.f90
- ppmlr.f90
- states.f90
- flatten.f90
- sweeptz.f90
- sweeppy.f90
- boundary.f90
- print.f90
- sweepx.f90
- 0.53% SWEEPX2
  - Loop@28
  - Loop@29
  - Loop@32
  - Loop@33
  - Loop@44
  - Loop@58
- sweepx1.f90

**Source - /lus/sonexion/heidi/reveal/sweepx2.f90**

```
L 32 do m = 1, npey
Lr8 33 do i = 1, isy
34   n = i + isy*(m-1) + 6
35   r(n) = recv2(1,k,i,j,m)
36   p(n) = recv2(2,k,i,j,m)
37   u(n) = recv2(3,k,i,j,m)
38   v(n) = recv2(4,k,i,j,m)
39   w(n) = recv2(5,k,i,j,m)
40   f(n) = recv2(6,k,i,j,m)
41   enddo
42   enddo
43
44 do i = 1,imax
45   n = i + 6
```

**Info - Line 33**

- A loop starting at line 33 was not vectored because it does not have enough iterations.
- A loop starting at line 33 was unrolled 6 times.

Double click on optimization message for more detailed information

**Reveal**

**OPT\_INFO: A loop starting at line %s was unrolled.**

The compiler unrolled the loop. Unrolling creates a number of copies of the loop body. When unrolling an outer loop, the compiler attempts to fuse replicated inner loops - a transformation known as unroll-and-jam. The compiler will always employ the unroll-and-jam mode when unrolling an outer loop; literal outer loop unrolling may occur when unrolling to satisfy a user directive (pragma).

This message indicates that unroll-and-jam was performed with respect to the identified loop. A different message is issued when literal outer loop unrolling is performed, as this transformation is far less likely to be beneficial.

For sake of illustration, the following contrasts unroll-and-jam with literal outer loop unrolling.

```
# 426 "/ptmp/uplib/buildslaves/pdgcs-81-edition-build/tbs/build/release/pdgcs/pdgc_ftn.msg.c"
DO J=1,10
  DO I=1,100
    A(I,J)=B(I,J) + 42.0
  ENDDO
ENDDO

DO J=1,10,2
  DO I=1,100
    A(I,J)=B(I,J) + 42.0 ! unroll-and-jam
    A(I,J+1)=B(I,J+1) + 42.0
  ENDDO
ENDDO

DO J=1,10,2
  DO I=1,100
    A(I,J)=B(I,J) + 42.0 ! literal outer unrol
  ENDDO
ENDDO
DO I=1,100
  A(I,J)=B(I,J+1) + 42.0
ENDDO
ENDDO
```

**Explain**

**Will not work with Intel Compiler**

The literal outer unroll code performs the same sequence of memory operations as the original nest, while the unroll-and-jam transformation interleaves operations from outer loop iterations. The compiler employs literal outerloop unrolling only when the data dependencies in the loop, or a control flow impediment, prevent fusion of the replicated inner loops. Literal outer loop unrolling is generally not desirable. It is provided to ensure expected behavior and for those rare instances where the user has determined that it is beneficial.

Explain other message... Close

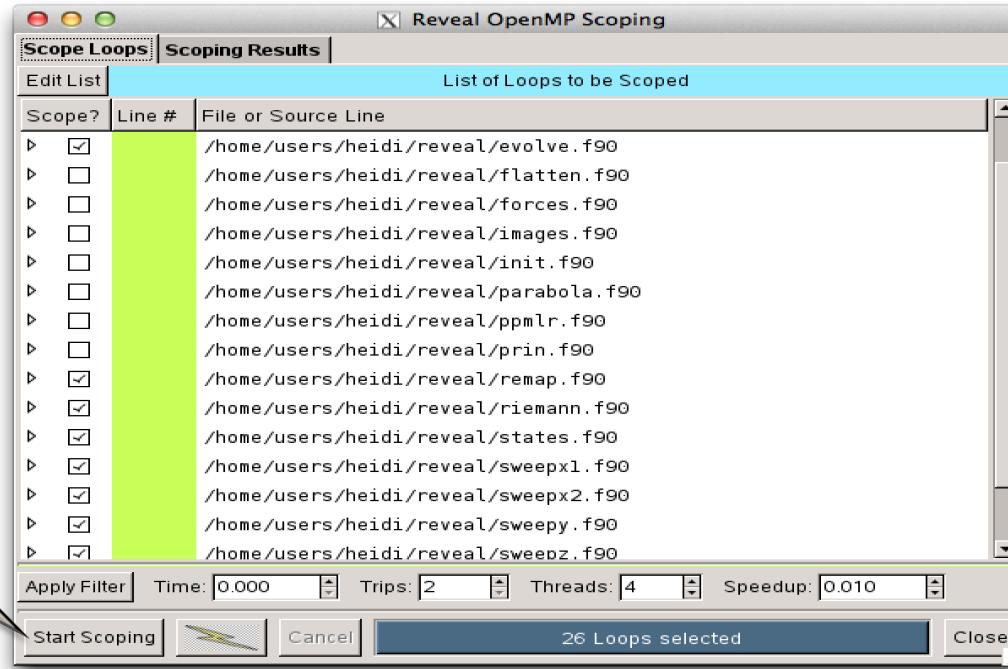
STORE

ANALYZE

# Scope Selected Loop(s)



Will not work  
with Intel Compiler



- Trigger dependence analysis
- scope loops above given threshold

# Review Scoping Results

Navigation

- Loop Performance

Source - /home/users/heidi/reveal/sweepz.f90

```
      49 call par
      50      do i = 1, isz
      51      radius = zxc(i+mpey*isz)
      52      theta = zyc(j+mpey*js)
      53      stheta = sin(theta)
      54      radius = radius * stheta
      55
      56      ! Put state variables into 1D arrays, padding with 6 ghost zones
      57      do m = 1, npez
      58          do k = 1, ks
      59              n = k + ks*(m-1) + 6
      60              r(n) = recv3(1,j,k,i,m)
      61              p(n) = recv3(2,j,k,i,m)
      62              u(n) = recv3(3,j,k,i,m)
      63              v(n) = recv3(4,j,k,i,m)
      64              w(n) = recv3(5,j,k,i,m)
      65
      66
```

Info - Line 51

- A loop starting at line 51 was scoped with errors. See Scoping Tool for more information.
  - "ppmlr" (called from "sweepz") was not inlined because I/O was detected in "volume".
  - "ppmlr" (called from "sweepz") was not inlined because the enclosing loop body did not completely flatten.
  - A loop starting at line 105 is flat (contains no external calls).
- A loop starting at line 105 was not vectorized because it does not map well onto the target architecture.
- A loop starting at line 105 was unrolled 8 times.
- A loop starting at line 51 was not vectorized because it contains a call to subroutine "ppmlr" on line 81.
- A loop starting at line 52 was not vectorized because it contains a call to subroutine "ppmlr" on line 81.
- A loop starting at line 59 is flat (contains no external calls).
- A loop starting at line 59 was not vectorized because a better candidate was found at line 60.
- A loop starting at line 60 is flat (contains no external calls).
- A loop starting at line 60 was not vectorized because it does not map well onto the target architecture.
- A loop starting at line 60 was unrolled 8 times.
- A loop starting at line 71 is flat (contains no external calls).
- A loop starting at line 71 was vectorized.

/home/users/heidi/reveal/vphone\_loops.ap2 loaded.

Will not work  
with Intel Compiler

Reveal OpenMP Scoping

Scope Loops Scoping Results

sweepz.f90: Loop#51

Name	Type	Scope	Info
wl@remap_	Scalar	Unresolved	<b>FAIL:</b> Possible recurrence involving this object. <b>FAIL:</b> Possible resolvable recurrence involving this object.
xa	Array	Unresolved	<b>FAIL:</b> Possible recurrence involving this object. <b>FAIL:</b> Possible resolvable recurrence involving this object. <b>WARN:</b> LastPrivate of array may be very expensive.
xa0	Array	Unresolved	<b>FAIL:</b> Possible recurrence involving this object. <b>FAIL:</b> Possible resolvable recurrence involving this object. <b>WARN:</b> LastPrivate of array may be very expensive.
i	Scalar	Private	
j	Scalar	Private	
k	Scalar	Private	
m	Scalar	Private	
n	Scalar	Private	
stheta	Scalar	Private	
theta	Scalar	Private	
gamm	Scalar	Shared	
isz	Scalar	Shared	
js	Scalar	Shared	
ks	Scalar	Shared	
mpey	Scalar	Shared	

First/Last Private

Enable FirstPrivate

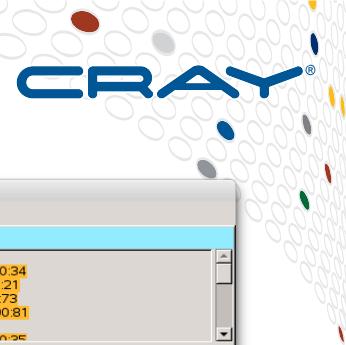
Enable LastPrivate

Find Name:

Insert Directive Show Directive

Close

Parallelization inhibitor messages are provided to assist user with analysis



# Review Scoping Results (continued)



Reveal OpenMP Scoping

Scope Loops Scoping Results

sweepy.f90: Loop@35

Call or I/O at line 62 of sweepy.f90  
4: /home/users/heidi/reveal/volume.f90:34  
3: /home/users/heidi/reveal/evolve.f90:21  
2: /home/users/heidi/reveal/volume.f90:35

Name	Type	Scope	Info
ks	Scalar	Shared	
mpey	Scalar	Shared	
ndim	Scalar	Shared	
npey	Scalar	Shared	
recv1	Array	Shared	
send2	Array	Shared	
svel RI	Scalar	Shared	WARN: atomic reduction operator required unless reduction fully
zdy	Array	Shared	
zxc	Array	Shared	
zya	Array	Shared	

First/Last Private  
 Enable FirstPrivate  
 Enable LastPrivate

Find Name:

Reduction  
None

Insert Directive Show Directive Close

Will not work with Intel Compiler

Reveal identifies calls that prevent parallelization

Reveal identifies shared reductions down the call chain

# Review Scoping Results (continued)



Reveal OpenMP Scoping

Scope Loops Scoping Results Footnote

m\_mat\_an.c: Loop@39

Name	Type	Scope	Info
a0i	Scalar	Private	
a0r	Scalar	Private	
a1i	Scalar	Private	
a1r	Scalar	Private	
a2i	Scalar	Private	
a2r	Scalar	Private	
b0i	Scalar	Private	
b0r	Scalar	Private	
b1i	Scalar	Private	
b1r	Scalar	Private	
b2i	Scalar	Private	
b2r	Scalar	Private	
j	Scalar	Private	
a	Scalar	Shared	<b>WARN:</b> Assuming no overlap with other objects. <b>INFO:</b> additional detail.
b	Scalar	Shared	<b>WARN:</b> Assuming no overlap with other objects. <b>INFO:</b> additional detail.
c	Scalar	Shared	<b>WARN:</b> Assuming no overlap with other objects. <b>INFO:</b> additional detail.

-First/Last Private  
 Enable FirstPrivate  
 Enable LastPrivate

-Reduction  
None

Find Name:

Insert Directive | Show Directive | Close

Will not work with Intel Compiler

Reveal OpenMP Scoping

Scope Loops Scoping Results Footnote

Scoping Footnote

Assume no overlap between lattice[\*].mom[\*] and tempmom[\*][\*]

Close

COMPUTE

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# Generate OpenMP Directives

```
! Directive inserted by Cray Reveal. May be incomplete.  
!$OMP parallel do default(None)  
!$OMP& unresolved (dvol,dx,dx0,e,f,flat,p,para,q,r,radius,svel,u,v,w,  
!$OMP& xa,xa0) &  
!$OMP& private (i,j,k,m,n,$$_n,delp2,delp1,shock,temp2,old_flat,  
!$OMP& onemfl,hdt,sinxf0,gamfac1,gamfac2,dtheta,deltx,fractn,  
!$OMP& ekin) &  
!$OMP& shared (gamm,isy,js,ks,mypey,ndim,ngeom,y,nlefty,npey,nrighty, &  
!$OMP& recv1,send2,zdy,zxc,zya) &  
do k = 1, ks  
do i = 1, isy  
    radius = zxc(i+mypey*isy)  
  
! Put state variables into 1D arrays, padding with 6 ghost zones  
do m = 1, npey  
do j = 1, js  
    n = j + js*(m-1) + 6  
    r(n) = recv1(1,k,j,i,m)  
    p(n) = recv1(2,k,j,i,m)  
    u(n) = recv1(4,k,j,i,m)  
    v(n) = recv1(5,k,j,i,m)  
    w(n) = recv1(3,k,j,i,m)  
    f(n) = recv1(6,k,j,i,m)  
enddo  
enddo  
  
do j = 1, jmax  
    n = j + 6
```

Will not work  
with Intel Compiler

Reveal generates OpenMP  
directive with illegal clause  
marking variables that need  
addressing

# Validate User Inserted Directives

Will not work  
with Intel Compiler



Reveal

File Edit View Help

vhone.pl

Navigation

Program View

Source - /ufs/home/users/heidi/reveal/riemann.f90

Up Down Save

Directive inserted by CRAY Reveal. May be incomplete.

R 64 !\$OMP parallel do default(none)

65 !\$OMP& private (l)

66 !\$OMP& shared (lmin,lmax,prgh,urgh,vrgh,plft,ulft,vlft,pmid,clft,

67 !\$OMP& crgh,gamfac1,gamfac2,plfti,pmold,prghi,umidl,umidr

68 !\$OMP& wlft,wrgh,zlft,zrgh,n)

LSg 69 do l = lmin, lmax

L 70 do n = 1, 12

71 pmold(l)

72 wlft (l)

73 wrgh (l)

74 wlft (l)

75 wrgh (l)

76 zlft (l)

User inserted directive with mis-scoped variable 'n'

Reveal Open Scope Loops Scoping Results

riemann.f90: Loop@69

Name	Type	Scope	Info
i	Scalar	Private	
n	Scalar	Private	WARN: Scope does not agree with user OMP directive.
clft	Array	Shared	
crgh	Array	Shared	
gamfac1	Scalar	Shared	
gamfac2	Scalar	Shared	

Info - Line 69

A loop starting at line 69 was not found

A loop starting at line 69 was parallel

First/Last Private

Enable FirstPrivate

Enable LastPrivate

Reduction

None

Find Name:

Insert Directive Show Directive Close

vhone.pl loaded

COMPUTE

STORE

ANALYZE

# Look For Vectorization Opportunities

Will not work **CRAY**  
with Intel Compiler

The screenshot shows the Reveal IDE interface. On the left, a large callout bubble points to the 'Compiler Messages' view in the navigation pane. The message list is titled 'Not Vectorized' and contains several entries, including:

- line 123
- images.f90 line 149
- init.f90 line 113 (0.000 sec)
- line 114 (0.000 sec)
- line 153 (0.000 sec)
- line 154 (0.000 sec)
- line 139
- prin.f90 line 125 (0.006 sec)
- line 42 (0.000 sec)
- line 43 (0.000 sec)
- line 127 (0.000 sec)
- line 128 (0.000 sec)
- line 129 (0.000 sec)
- line 104 (0.000 sec)
- riemann.f90 line 63 (0.387 sec)
- line 64 (0.224 sec)
- sweeplex1.f90 line 31 (4.053 sec)
- line 32 (4.053 sec)
- line 59 (0.093 sec)

The main window displays the source code for 'vhone.pl' with line numbers 62 through 79. The code is color-coded by column (FS, F, S, L) and contains loops and mathematical calculations. The 'Info - Line 64' panel at the bottom provides analysis for line 64:

- A loop starting at line 64 is flat (contains no external calls).
- A loop starting at line 64 was not vectorized because a recurrence was found on "pmid" at line 77.

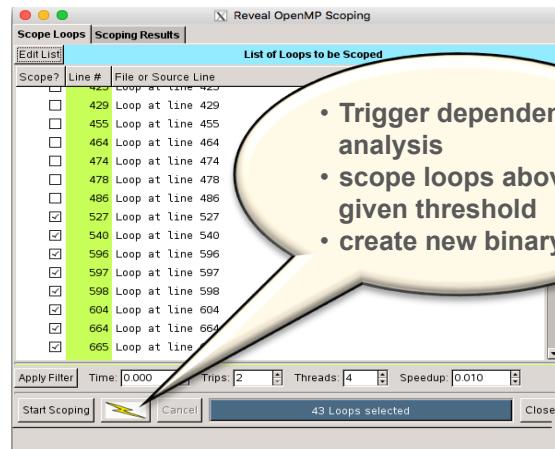
At the bottom of the interface, a status bar indicates: 'vhone.pl loaded. vhone\_loops.ap2 loaded.'

# Reveal Auto-Parallelization

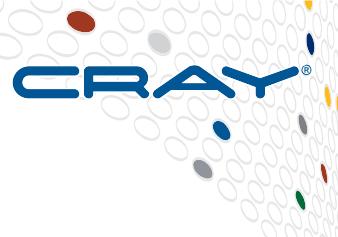
Will not work  
with Intel Compiler



- Build an **experimental binary** that includes automatic runtime-assisted parallelization
- **No source code changes required** to see if high level loops that contain calls can be automatically parallelized
- Result includes **parallelization of serial loops via traditional OpenMP** as well as **more extensive loop optimizations**
- **User Workflow:**
  1. Obtain **loop work estimates** using CCE and perftools-lite-loops
  2. Use Reveal and CCE's program library to **parallelize loops and create experimental binary**
  3. Run experimental binary and compare performance against baseline
  4. If auto-parallelization is successful, **use Reveal to insert parallel directives into source**



# Reveal Analysis Feedback



Will not work  
with Intel Compiler

Reveal

File Edit View Help

mg.f

Navigation

- Loop Performance
- 12.5484 MG\_MPI@245 ★
- 6.5747 MG\_MPI@664 ★
- 0.1878 Instance #1
- 0.2593 Instance #2
- 0.0130 Instance #3
- 2.8343 Instance #4
- 2.8454 Instance #5
- 0.1455 Instance #6
- 0.1447 Instance #7
- 0.1448 Instance #8
- 6.5736 MG\_MPI@665 ★
- 3.8584 MG\_MPI@666 ★
- 2.8689 MG\_MPI@596 ★
- 2.8683 MG\_MPI@597 ★
- 2.5819 MG\_MPI@672 ★
- 1.8738 MG\_MPI@1069 ★
- 1.7607 MG\_MPI@598 ★
- 1.5160 MG\_MPI@540 ★
- 1.4182 MG\_MPI@527 ★
- 1.0593 MG\_MPI@834 ★
- 1.0590 MG\_MPI@835 ★
- 1.0583 MG\_MPI@750 ★
- 1.0580 MG\_MPI@753 ★
- 1.0363 MG\_MPI@604 ★
- 0.9935 MG\_MPI@1154 ★

Source - /home/heidi/demos/NPB/NPB3.2-MPI-mg/MG/mg.f

663 do i3=2,n3-1

664 Fg

665 do i2=n2-1

666 F

667 do il=1,nl

668 ul(il) = u(il,i2-1,i3) + u(il,i2+1,i3)

669 + u(il,i2,i3-1) + u(il,i2,i3+1)

670 u2(il) = u(il,i2-1,i3-1) + u(il,i2+1,i3-1)

671 > + u(il,i2-1,i3+1) + u(il,i2+1,i3+1)

672 enddo

673 do il=2,nl-1

674 r(il,i2,i3) = v(il,i2,i3)

Info - Line 664

- A loop starting at line 664 was scoped without errors.
- A loop starting at line 664 is flat (contains no external calls).
- A loop starting at line 664 was not vectorized because a recurrence was found on "u1" between lines 667 and 672.
- A loop starting at line 664 was partitioned.
- A loop starting at line 665 is flat (contains no external calls).
- A loop starting at line 665 was not vectorized because a recurrence was found on "u1" between lines 667 and 672.
- A loop starting at line 666 is flat (contains no external calls).
- A loop starting at line 666 was unrolled 4 times.
- A loop starting at line 666 was vectorized.
- A loop starting at line 672 is flat (contains no external calls).
- A loop starting at line 672 was unrolled 4 times.
- A loop starting at line 672 was vectorized.

/home/heidi/demos/NPB/NPB3.2-MPI-mg/bin/mg.C.16+17976-76t.ap2 loaded.

Reveal OpenMP Scoping

Scope Loops Scoping Results Build Results

Your binary was rebuilt with the following changes.

/home/heidi/demos/NPB/NPB3.2-MPI-mg/MG/mg.f

- OMP loop at line 596
- OMP loop at line 664
- OMP loop at line 750
- OMP loop at line 834
- OMP loop at line 995
- OMP loop at line 1154
- Autothreaded loop at line 1199
- Autothreaded loop at line 1213
- Autothreaded loop at line 1262
- Autothreaded loop at line 1276
- Autothreaded loop at line 1326
- Autothreaded loop at line 1335
- Autothreaded loop at line 1370
- Autothreaded loop at line 1379
- OMP loop at line 2173
- OMP loop at line 2435

Close

COMPUTE

STORE

ANALYZE

# Reveal Auto-Parallelization Recap



- Minimal user time investment includes time to set up and run optimization experiment
  - Collect loop work estimates
  - Build program library
  - Click button in Reveal
  - Run experimental binary and compare against original program
- Even if experiment does not yield a performance improvement, Reveal will provide insight into parallelization issues
- Target: when a pure MPI solution does not scale
- Use on X86 hardware (AMD Interlagos, Intel Broadwell, etc.)

# Optimal Programming for the KNL



- **First and foremost, make sure you are taking advantage of available memory bandwidth**
  - If you are memory bandwidth limited, vectorization and threading will not help
  - If on the other hand, if you take steps to improve memory bandwidth utilization then vectorization and threading will give some improvement
- **If you do not vectorize or cannot not vectorize your application, KNL will be of little value**
- **Bottom line – you must take advantage of MCDRAM and then vectorize**
- **Utilization of available cores then the final important step for taking advantage of the KNL**

# Taking advantage of available Bandwidth



- Need less than 16 Gbytes?      You got it made
- Need more than 16 Gbytes
  - MUST look at cache utilization.
    - KNL has small amount of lower level cache, hard to block for Lower level cache; however, blocking for MCDRAM is important and doable
  - Striding and indirect address needs to be examined
    - Can it be eliminated?
    - Can the working set be blocked to fit into MCDRAM used as Cache?
  - Eliminate unnecessary stores

# Vectorization



- Vectorize
- Check Alignment

---

COMPUTE

STORE

ANALYZE

# Explanation of tests

- These are small tests which is not really good for timing on KNL since the cache is always flushed prior to the test. This means that some are latency dominated, so DDR may be faster than MCDRAM.
- Trying to illustrate what the compilers, CCE and Intel do with different types of loops. Performance variations are largely dependent upon cache reuse as much as vectorization

# Impact of size of stride



## ORIGINAL

```
48.      1          C      DIMENSION A(128,N)
49.      1
50.      1 Vr2--<      DO 41080  I = 1,N
51.      1 Vr2          A( 1,I) = C1*A(13,I) + C2* A(12,I) + C3*A(11,I) +
52.      1 Vr2          *          C4*A(10,I) + C5* A( 9,I) + C6*A( 8,I) +
53.      1 Vr2          *          C7*A( 7,I) + C0*(A( 5,I) + A( 6,I) ) + A( 3,I)
54.      1 Vr2--> 41080 CONTINUE
```

## RESTRUCTURED

```
75.      1          C      DIMENSION B(13,N)
76.      1
77.      1 Vr2--<      DO 41081  I = 1,N
78.      1 Vr2          B( 1,I) = C1*B(13,I) + C2* B(12,I) + C3*B(11,I) +
79.      1 Vr2          *          C4*B(10,I) + C5* B( 9,I) + C6*B( 8,I) +
80.      1 Vr2          *          C7*B( 7,I) + C0*(B( 5,I) + B( 6,I) ) + B( 3,I)
81.      1 Vr2--> 41081 CONTINUE
```



# Impact of size of stride Intel -qopt-report-phase=all

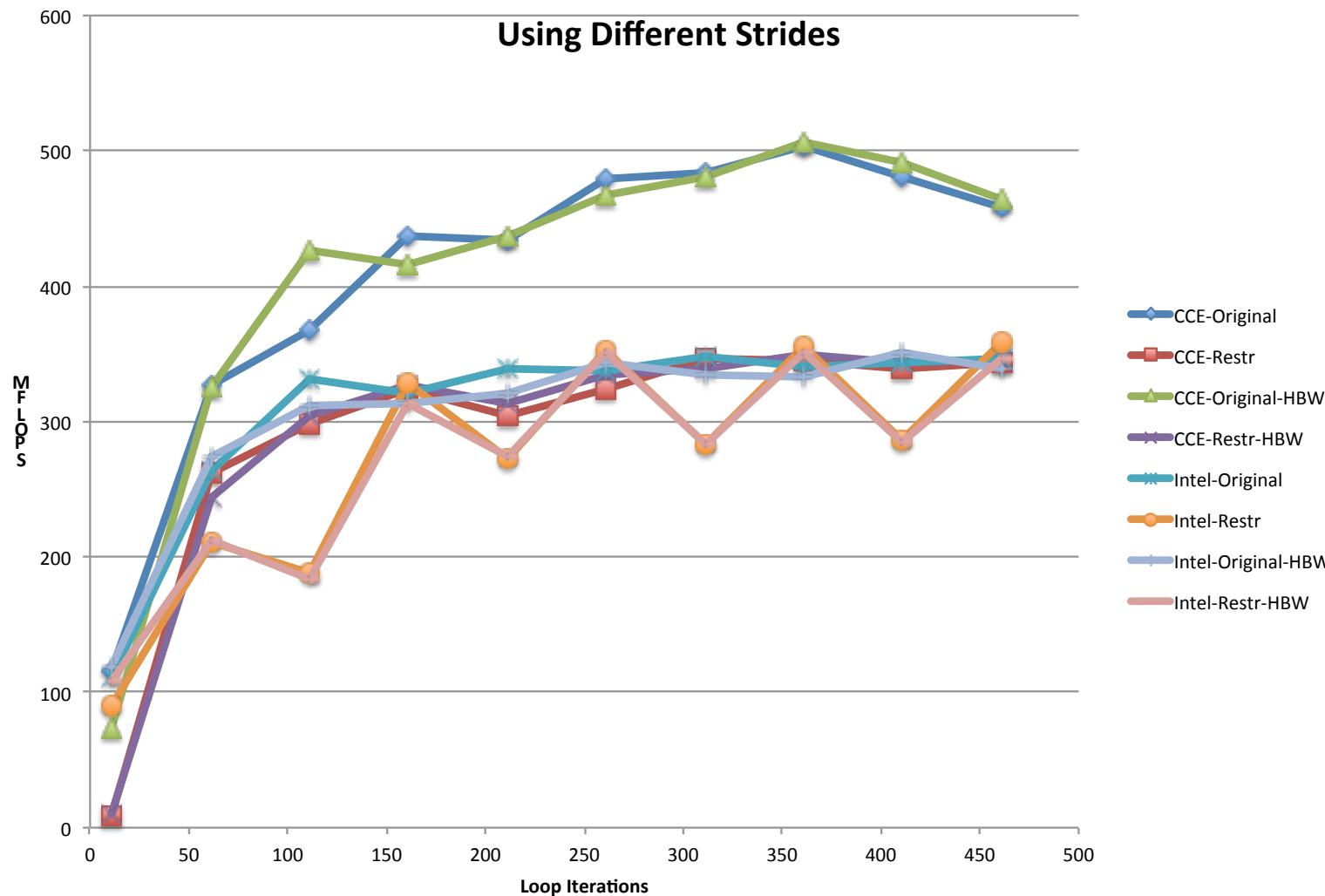
LOOP BEGIN at lp41080.f(50,10)  
  remark #15300: LOOP WAS VECTORIZED  
LOOP END

LOOP BEGIN at lp41080.f(50,10)  
<Remainder loop for vectorization>  
  remark #15301: REMAINDER LOOP WAS VECTORIZED  
LOOP END

LOOP BEGIN at lp41080.f(77,10)  
  remark #15300: LOOP WAS VECTORIZED  
LOOP END

LOOP BEGIN at lp41080.f(77,10)  
<Remainder loop for vectorization>  
  remark #15301: REMAINDER LOOP WAS VECTORIZED  
LOOP END

## Using Different Strides



# Discussion – Impact of Stride

- **My Guess**
  - With the larger stride, CCE used scalar instructions to fetch elements
  - With the smaller stride, CCE and Intel used gather/scatter
- **Need to confirm with Compiler people**

# Impact of size of stride

```
59.      1           C      THE ORIGINAL
60.      1
61. + 1 2-----<      DO 41090 K = KA, KE, -1
62.   1 2 iVp-----<      DO 41090 J = JA, JE
63. + 1 2 iVp i-----<      DO 41090 I = IA, IE
64.   1 2 iVp i          A(K,L,I,J) = A(K,L,I,J) - B(J,1,i,k)*A(K+1,L,I,1)
65.   1 2 iVp i          * - B(J,2,i,k)*A(K+1,L,I,2) - B(J,3,i,k)*A(K+1,L,I,3)
66.   1 2 iVp i          * - B(J,4,i,k)*A(K+1,L,I,4) - B(J,5,i,k)*A(K+1,L,I,5)
67.   1 2 iVp i----->>> 41090 CONTINUE

96.      1           C      THE RESTRUCTURED
97.      1
98. + 1 2-----<      DO 41091 K = KA, KE, -1
99. + 1 2 3-----<      DO 41091 J = JA, JE
100.    1 2 3 Vr2-----<      DO 41091 I = IA, IE
101.    1 2 3 Vr2          AA(I,K,L,J) = AA(I,K,L,J) - BB(I,J,1,K)*AA(I,K+1,L,1)
102.    1 2 3 Vr2          * - BB(I,J,2,K)*AA(I,K+1,L,2) - BB(I,J,3,K)*AA(I,K+1,L,3)
103.    1 2 3 Vr2          * - BB(I,J,4,K)*AA(I,K+1,L,4) - BB(I,J,5,K)*AA(I,K+1,L,5)
104.    1 2 3 Vr2----->>> 41091 CONTINUE
```

# Stride versus no Stride



```
LOOP BEGIN at lp41090.f(48,10)
  remark #15542: loop was not vectorized: inner loop was already vectorized

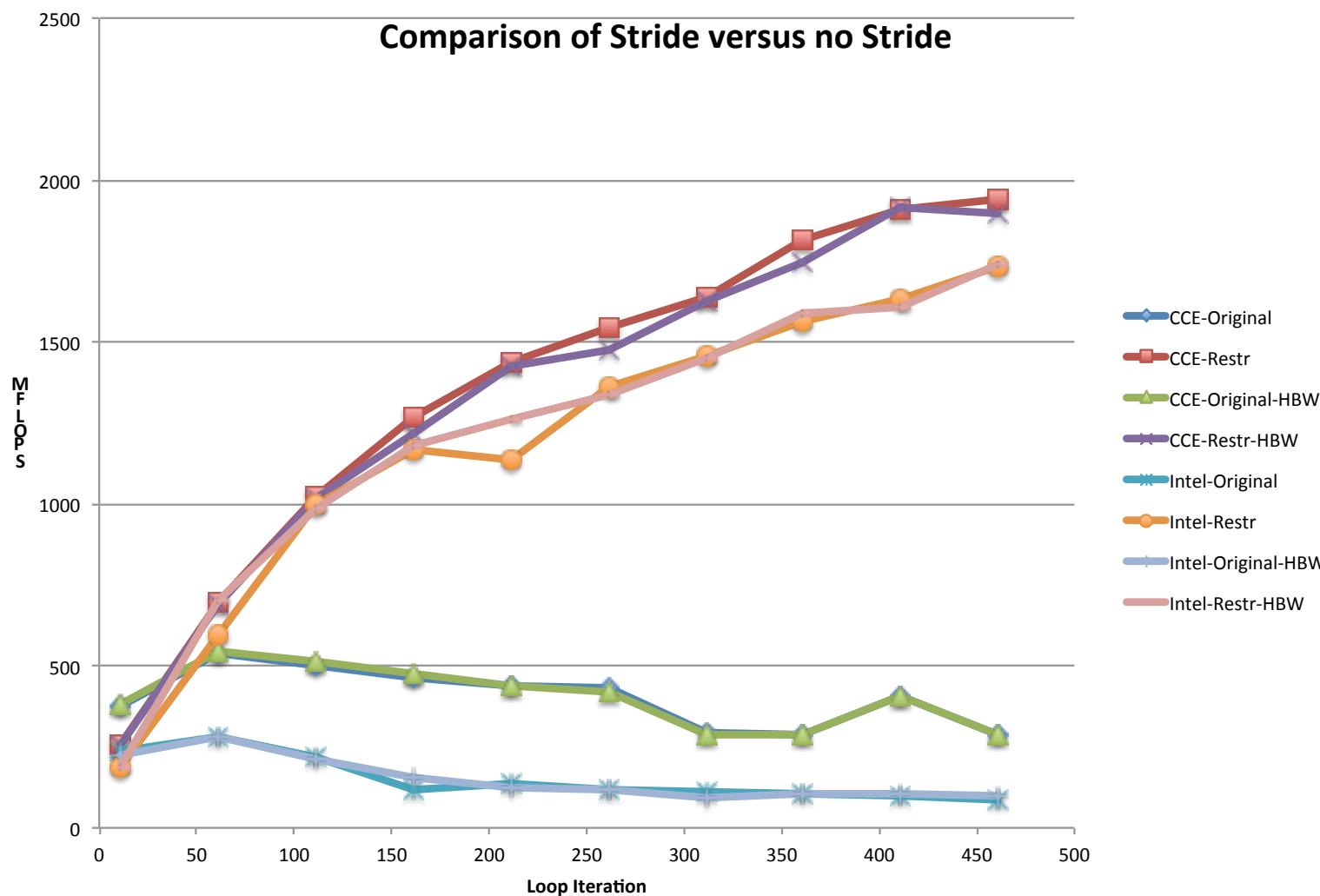
LOOP BEGIN at lp41090.f(61,10)
  remark #15542: loop was not vectorized: inner loop was already vectorized

LOOP BEGIN at lp41090.f(62,12)
  remark #15542: loop was not vectorized: inner loop was already vectorized

LOOP BEGIN at lp41090.f(63,14)
  remark #15300: LOOP WAS VECTORIZED
  LOOP END

  LOOP BEGIN at lp41090.f(63,14)
  <Remainder loop for vectorization>
    remark #15335: remainder loop was not vectorized: vectorization possible but seems inefficient.
                  Use vector always directive or -vec-threshold0 to override
    LOOP END
  LOOP END
LOOP END
```

## Comparison of Stride versus no Stride





# Impact of size of stride

ORIGINAL

```
43.    1           C  GAUSS ELIMINATION
44.    1
45. + 1 2-----<      DO 43020 I = 1, MATDIM
46.    1 2           A(I,I) = 1. / A(I,I)
47.    1 2 r2-----<    DO 43020 J = I+1, MATDIM
48.    1 2 r2           A(J,I) = A(J,I) * A(I,I)
49.    1 2 r2 Vr2----<  DO 43020 K = I+1, MATDIM
50.    1 2 r2 Vr2           A(J,K) = A(J,K) - A(J,I) * A(I,K)
51.    1 2 r2 Vr2-->>> 43020 CONTINUE
```

RESTRUCTURED

```
83. + 1 2-----<      DO 43021 I = 1, MATDIM
84.    1 2           A(I,I) = 1. / A(I,I)
85.    1 2 r2-----<    DO 43021 J = I+1, MATDIM
86.    1 2 r2           A(J,I) = A(J,I) * A(I,I)
88.    1 2 r2           CDIR$ IVDEP
90.    1 2 r2 Vr2----<  DO 43021 K = I+1, MATDIM
91.    1 2 r2 Vr2           A(J,K) = A(J,K) - A(J,I) * A(I,K)
92.    1 2 r2 Vr2-->>> 43021 CONTINUE
```

# Gaussian Elimination



```
LOOP BEGIN at lp43020.f(45,10)
  remark #15344: loop was not vectorized: vector dependence prevents vectorization.
  First dependence is shown below. Use level 5 report for details
  remark #15346: vector dependence: assumed OUTPUT dependence between line 46 and line 50

  LOOP BEGIN at lp43020.f(47,11)
  <Distributed chunk1>
    remark #25426: Loop Distributed (2 way)
    remark #15335: loop was not vectorized: vectorization possible but seems inefficient.
    Use vector always directive or -vec-threshold0 to override
    remark #25439: unrolled with remainder by 8
  LOOP END

  LOOP BEGIN at lp43020.f(47,11)
  <Distributed chunk2>
    remark #15335: loop was not vectorized: vectorization possible but seems inefficient.
    Use vector always directive or -vec-threshold0 to override

  LOOP BEGIN at lp43020.f(49,12)
    remark #15335: loop was not vectorized: vectorization possible but seems inefficient.
    Use vector always directive or -vec-threshold0 to override
    remark #25439: unrolled with remainder by 2
  LOOP END

  LOOP BEGIN at lp43020.f(49,12)
  <Remainder>
  LOOP END
LOOP END

  LOOP BEGIN at lp43020.f(47,11)
  <Remainder, Distributed chunk1>
    remark #25436: completely unrolled by 7
  LOOP END
LOOP END
```



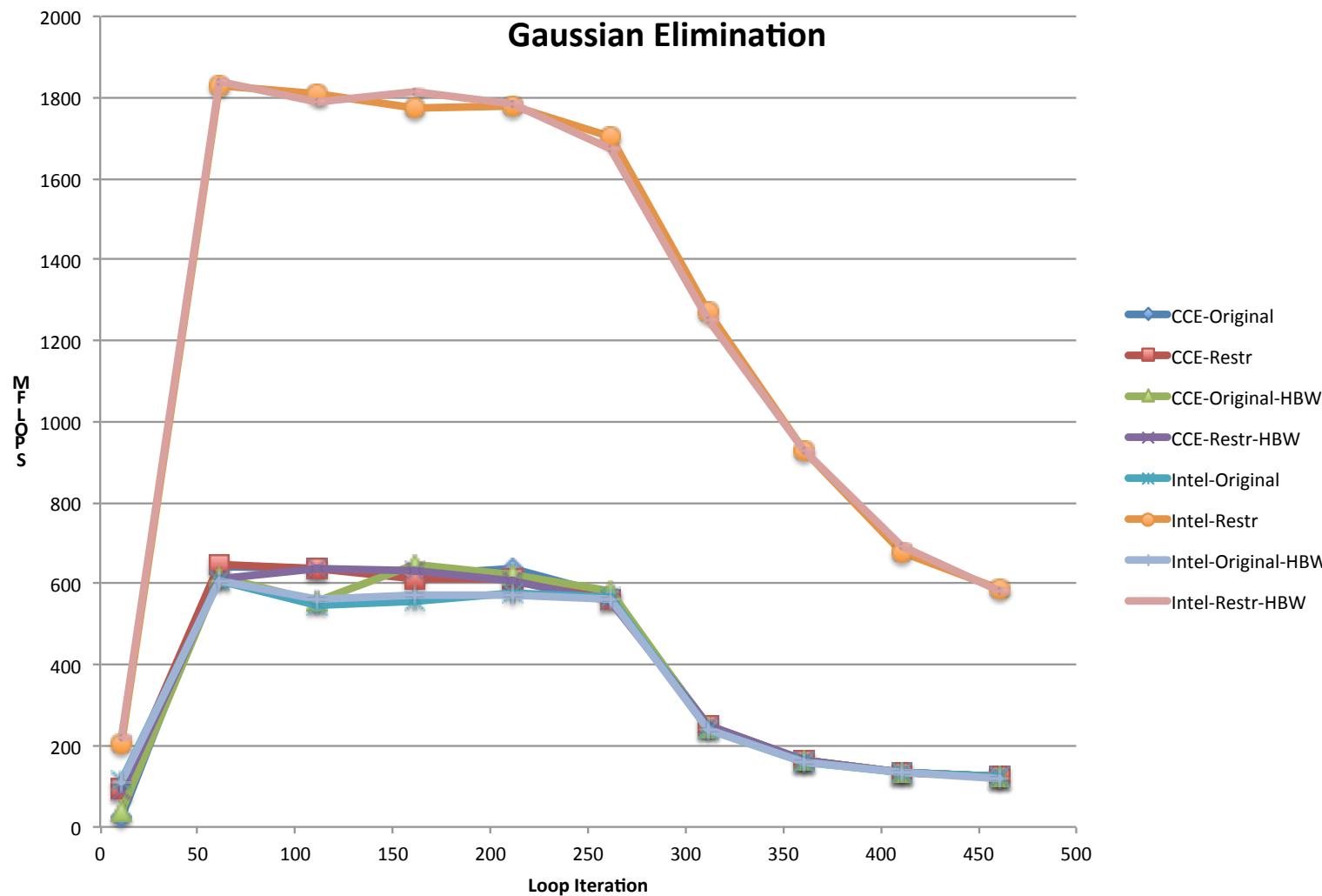
# Impact of size of stride

```
LOOP BEGIN at lp43020.f(83,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

LOOP BEGIN at lp43020.f(85,11)
    remark #15301: OUTER LOOP WAS VECTORIZED

LOOP BEGIN at lp43020.f(90,12)
    remark #15335: loop was not vectorized: vectorization possible but seems inefficient.
                  Use vector always directive or -vec-threshold0 to override
    LOOP END
    LOOP END
LOOP END
```

# Gaussian Elimination



# Recursion on One of Two dimensions



```
47.    1          C      THE ORIGINAL
48.    1
49.    1 V-----<      DO 43100 J = 1, N
50.    1 V          AH = B(J) - B(J-1)
51.    1 V r4-----<      DO 43100 I = 2, N
52.    1 V r4          A(I,J) = AH * A(I-1,J) + C(I,J)
53.    1 V r4-----> 43100 CONTINUE
```

```
76.    1          C      THE RESTRUCTURED
77.    1
78.    1 Vr2-----<      DO 43101 J = 1, N
79.    1 Vr2          VAH(J) = B(J) - B(J-1)
80.    1 Vr2-----> 43101 CONTINUE
81.    1
82.    1 ir4-----<      DO 43102 I = 2, N
83.    1 ir4 iV-----<      DO 43102 J = 1, N
84.    1 ir4 iV          A(I,J) = VAH(J) * A(I-1,J) + C(I,J)
85.    1 ir4 iV-----> 43102 CONTINUE
```

# Recursion on One of Two dimensions



```
LOOP BEGIN at lp43100.f(49,10)
  remark #15301: OUTER LOOP WAS VECTORIZED

  LOOP BEGIN at lp43100.f(51,11)
    remark #15344: loop was not vectorized: vector dependence prevents vectorization.
      First dependence is shown below. Use level 5 report for details
    remark #15346: vector dependence: assumed FLOW dependence between line 52 and line 52
    LOOP END
  LOOP END

  LOOP BEGIN at lp43100.f(49,10)
<Remainder loop for vectorization>
  remark #15301: REMAINDER LOOP WAS VECTORIZED

  LOOP BEGIN at lp43100.f(51,11)
    remark #25460: No loop optimizations reported
    LOOP END
  LOOP END
```

# Recursion on One of Two dimensions



```
LOOP BEGIN at lp43100.f(78,10)
<Peeled loop for vectorization>
    remark #15301: PEEL LOOP WAS VECTORIZED
LOOP END

LOOP BEGIN at lp43100.f(78,10)
    remark #15300: LOOP WAS VECTORIZED
LOOP END

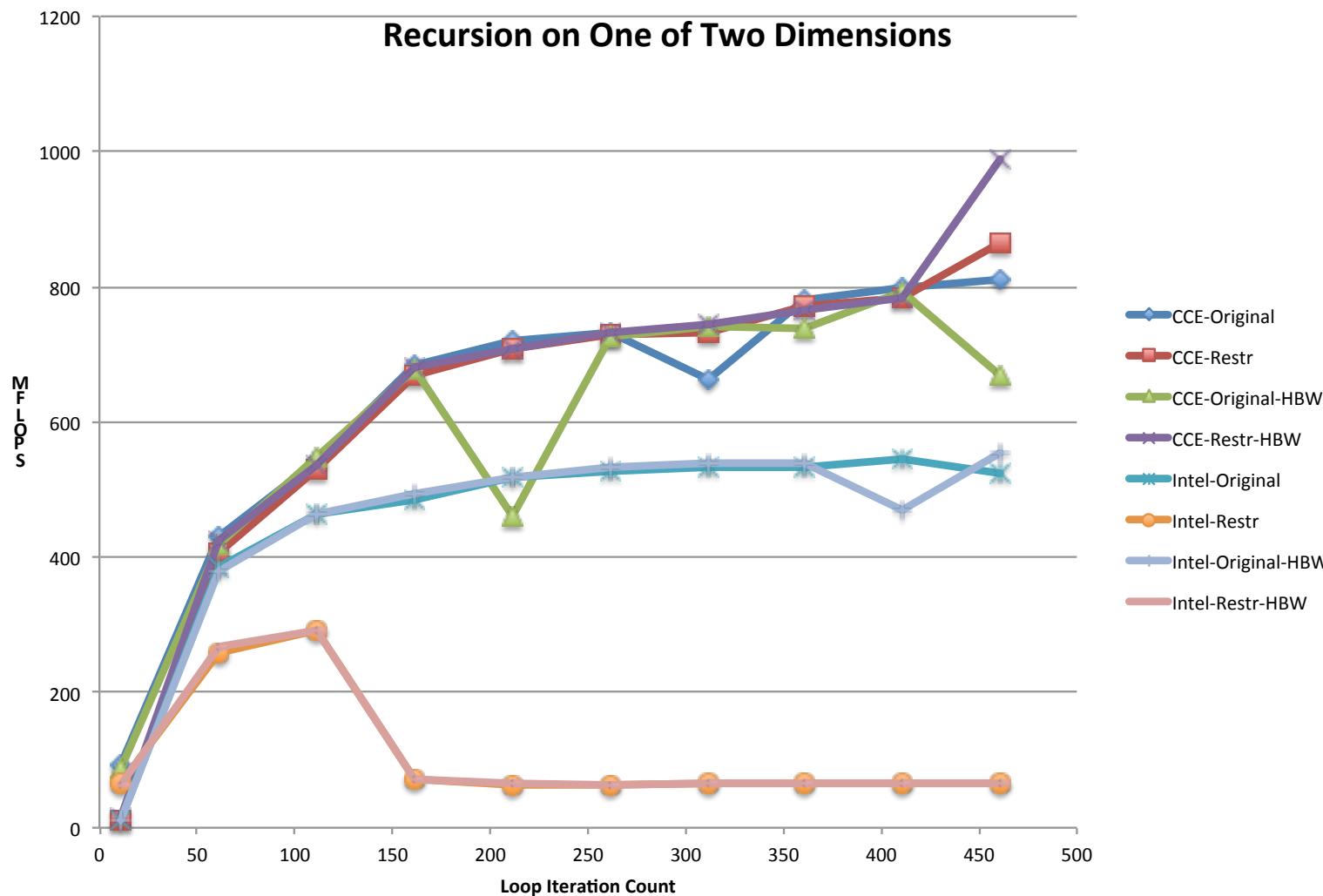
LOOP BEGIN at lp43100.f(78,10)
<Remainder loop for vectorization>
    remark #15301: REMAINDER LOOP WAS VECTORIZED
LOOP END

LOOP BEGIN at lp43100.f(82,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

LOOP BEGIN at lp43100.f(83,11)
    remark #15300: LOOP WAS VECTORIZED
LOOP END

LOOP BEGIN at lp43100.f(83,11)
<Remainder loop for vectorization>
    remark #15335: remainder loop was not vectorized: vectorization possible but seems inefficient.
                    Use vector always directive or -vec-threshold0 to override
LOOP END
```

## Recursion on One of Two Dimensions



# Alternating Direction Implicit



```
48.      1           C      THE ORIGINAL
49.      1
50.      1 V-----<      DO 43111 J = 2, N
51.      1 V           AH = B(J) - B(J-1)
52.      1 V r4-----<      DO 43110 I = 2, N
53.      1 V r4           A(I,J) = AH * A(I-1,J) + C(I,J)
54.      1 V r4-----> 43110 CONTINUE
55.      1 V
56.      1 V           BH = D(J) - D(J-1)
57.      1 V r4-----<      DO 43112 I = N, 2, - 1
58.      1 V r4           A(I,J) = BH * A(I+1,J) + C(I,J)
59.      1 V r4-----> 43112 CONTINUE
60.      1 V
61.      1 V-----> 43111 CONTINUE
```

# Alternating Direction Implicit



```
85.      1           C      THE RESTRUCTURED
86.      1
87.      1 fV-----<      DO 43113 J = 2, N
88.      1 fV           VAH(J) = B(J) - B(J-1)
89.      1 fV-----> 43113 CONTINUE
90.      1
91.      1 ir4-----<      DO 43114 I = 2, N
92.      1 ir4 fi----<      DO 43114 J = 2, N
93.      1 ir4 fi           A(I,J) = VAH(J) * A(I-1,J) + C(I,J)
94.      1 ir4 fi---->> 43114 CONTINUE
95.      1
96.      1 f-----<      DO 43115 J = 2, N
97.      1 f           VBH(J) = D(J) - D(J-1)
98.      1 f-----> 43115 CONTINUE
99.      1
100.     1 ir4-----<      DO 43116 I = N, 2, - 1
101.     1 ir4 iV----<      DO 43116 J = 2, N
102.     1 ir4 iV           A(I,J) = VBH(J) * A(I+1,J) + C(I,J)
103.     1 ir4 iV---->> 43116 CONTINUE
```

# Alternating Direction Implicit



```
LOOP BEGIN at lp43110.f(50,10)
    remark #15301: OUTER LOOP WAS VECTORIZED

    LOOP BEGIN at lp43110.f(52,11)
        remark #15344: loop was not vectorized: vector dependence prevents vectorization.
            First dependence is shown below. Use level 5 report for details
        remark #15346: vector dependence: assumed FLOW dependence between line 53 and line 53
    LOOP END

    LOOP BEGIN at lp43110.f(57,11)
        remark #15344: loop was not vectorized: vector dependence prevents vectorization.
            First dependence is shown below. Use level 5 report for details
        remark #15346: vector dependence: assumed FLOW dependence between line 58 and line 58
    LOOP END
    LOOP END

    LOOP BEGIN at lp43110.f(50,10)
    <Remainder loop for vectorization>
        remark #15301: REMAINDER LOOP WAS VECTORIZED

        LOOP BEGIN at lp43110.f(52,11)
            remark #25460: No loop optimizations reported
        LOOP END

        LOOP BEGIN at lp43110.f(57,11)
            remark #25460: No loop optimizations reported
    
```

---

C O M P U T E | S T O R E | A N A L Y Z E

LOOP END      LOOP END

# Alternating Direction Implicit



```
LOOP BEGIN at lp43110.f(87,10)
  <Peeled loop for vectorization>
    remark #15301: PEEL LOOP WAS VECTORIZED
  LOOP END

  LOOP BEGIN at lp43110.f(87,10)
    remark #15300: LOOP WAS VECTORIZED
  LOOP END

  LOOP BEGIN at lp43110.f(87,10)
  <Remainder loop for vectorization>
    remark #15301: REMAINDER LOOP WAS VECTORIZED
  LOOP END
  LOOP BEGIN at lp43110.f(91,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

  LOOP BEGIN at lp43110.f(92,11)
    remark #15300: LOOP WAS VECTORIZED
  LOOP END

  LOOP BEGIN at lp43110.f(92,11)
  <Remainder loop for vectorization>
    remark #15335: remainder loop was not vectorized: vectorization possible but seems inefficient.
                  Use vector always directive or -vec-threshold0 to override
  LOOP END
LOOP END
```

# Alternating Direction Implicit



```
LOOP BEGIN at lp43110.f(96,10)
<Peeled loop for vectorization>
    remark #15301: PEEL LOOP WAS VECTORIZED
LOOP END

LOOP BEGIN at lp43110.f(96,10)
    remark #15300: LOOP WAS VECTORIZED
LOOP END

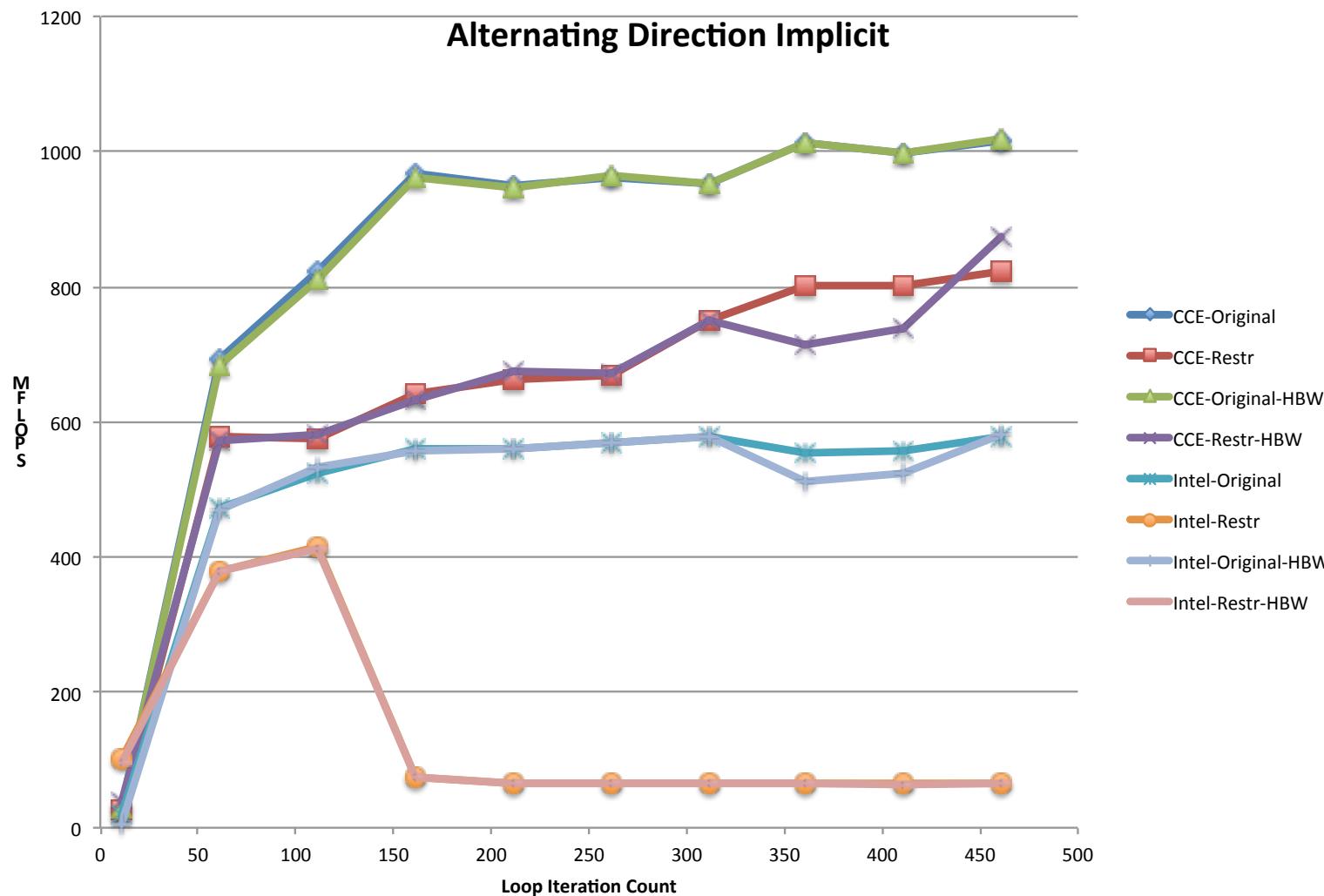
LOOP BEGIN at lp43110.f(96,10)
<Remainder loop for vectorization>
    remark #15301: REMAINDER LOOP WAS VECTORIZED
LOOP END

LOOP BEGIN at lp43110.f(100,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

LOOP BEGIN at lp43110.f(101,11)
    remark #15300: LOOP WAS VECTORIZED
LOOP END

LOOP BEGIN at lp43110.f(101,11)
<Remainder loop for vectorization>
    remark #15335: remainder loop was not vectorized: vectorization possible but seems inefficient. Use v
LOOP END
LOOP END
```

# Alternating Direction Implicit



# Strongly Implicit Procedure



```
51.      1           C      THE ORIGINAL
52.      1
53. + 1 2-----<      DO 43140 J = 2, N
54. + 1 2 r2-----<      DO 43140 I = 2, N
55.      1 2 r2           A(I,J,1) = A(I,J,1) - B(I,J) * A(I-1,J,1)
56.      1 2 r2           *                               - C(I,J) * A(I,J-1,1)
57.      1 2 r2           A(I,J,2) = A(I,J,2) - B(I,J) * A(I-1,J,2)
58.      1 2 r2           *                               - C(I,J) * A(I,J-1,2)
59.      1 2 r2           A(I,J,3) = A(I,J,3) - B(I,J) * A(I-1,J,3)
60.      1 2 r2           *                               - C(I,J) * A(I,J-1,3)
61.      1 2 r2----->> 43140 CONTINUE
```

# Strongly Implicit Procedure



```
89.      1          C      THE RESTRUCTURED
90.      1
91.      1          NDIAGS = 2 * N - 3
92.      1          ISTART = 1
93.      1          JSTART = 2
94.      1          LDIAG  = 0
95. + 1 2-----< DO 43141 IDIAGS = 1, NDIAGS
96.    1 2          IF(IDIAGS .LE. N-1 ) THEN
97.      1 2          ISTART = ISTART + 1
98.      1 2          LDIAG  = LDIAG  + 1
99.      1 2          ELSE
100.     1 2          JSTART = JSTART + 1
101.     1 2          LDIAG  = LDIAG  - 1
102.     1 2          ENDIF
103.     1 2          I = ISTART + 1
104.     1 2          J = JSTART - 1
105.     1 2          CDIR$ IVDEP
106.     1 2          CVD$ NODEPCHK
107.     1 2          *VDIR NODEP
108.    1 2 Vr2-----< DO 43142 IPOINT = 1, LDIAG
109.    1 2 Vr2          I = I - 1
110.    1 2 Vr2          J = J + 1
111.    1 2 Vr2          A(I,J,1) = A(I,J,1) - B(I,J) * A(I-1,J,1)
112.    1 2 Vr2          *                                - C(I,J) * A(I,J-1,1)
113.    1 2 Vr2          A(I,J,2) = A(I,J,2) - B(I,J) * A(I-1,J,2)
114.    1 2 Vr2          *                                - C(I,J) * A(I,J-1,2)
115.    1 2 Vr2          A(I,J,3) = A(I,J,3) - B(I,J) * A(I-1,J,3)
116.    1 2 Vr2          *                                - C(I,J) * A(I,J-1,3)
117.    1 2 Vr2-----> 43142 CONTINUE
118.    1 2-----> 43141 CONTINUE
```

COMPUTE

STORE

ANALYZE

# Strongly Implicit Procedure



```
89.    1          C      THE RESTRUCTURED
90.    1
91.    1          NDIAGS = 2 * N - 3
92.    1          ISTART = 1
93.    1          JSTART = 2
94.    1          LDIAG  = 0
95. + 1 2-----< DO 43141 IDIAGS = 1, NDIAGS
96.    1 2          IF(IDIAGS .LE. N-1 ) THEN
97.    1 2          ISTART = ISTART + 1
98.    1 2          LDIAG  = LDIAG  + 1
99.    1 2          ELSE
100.   1 2          JSTART = JSTART + 1
101.   1 2          LDIAG  = LDIAG  - 1
102.   1 2          ENDIF
103.   1 2          I = ISTART + 1
104.   1 2          J = JSTART - 1
105.   1 2          CDIR$ IVDEP
106.   1 2          CVD$ NODEPCHK
107.   1 2          *VDIR NODEP
108.   1 2 Vr2-----< DO 43142 IPOINT = 1, LDIAG
109.   1 2 Vr2          I = I - 1
110.   1 2 Vr2          J = J + 1
111.   1 2 Vr2          A(I,J,1) = A(I,J,1) - B(I,J) * A(I-1,J,1)
112.   1 2 Vr2          *                                - C(I,J) * A(I,J-1,1)
113.   1 2 Vr2          A(I,J,2) = A(I,J,2) - B(I,J) * A(I-1,J,2)
114.   1 2 Vr2          *                                - C(I,J) * A(I,J-1,2)
115.   1 2 Vr2          A(I,J,3) = A(I,J,3) - B(I,J) * A(I-1,J,3)
116.   1 2 Vr2          *                                - C(I,J) * A(I,J-1,3)
117.   1 2 Vr2-----> 43142 CONTINUE
118.   1 2-----> 43141 CONTINUE
```

COMPUTE

STORE

ANALYZE

# Strongly Implicit Procedure



```
LOOP BEGIN at lp43140.f(53,10)
    remark #15344: loop was not vectorized: vector dependence prevents vectorization.
        First dependence is shown below. Use level 5 report for details
    remark #15346: vector dependence: assumed FLOW dependence between line 55 and line 55

    LOOP BEGIN at lp43140.f(54,11)
        remark #15344: loop was not vectorized: vector dependence prevents vectorization.
            First dependence is shown below. Use level 5 report for details
        remark #15346: vector dependence: assumed FLOW dependence between line 55 and line 55
        remark #25456: Number of Array Refs Scalar Replaced In Loop: 7
    LOOP END
LOOP END
```

# Strongly Implicit Procedure

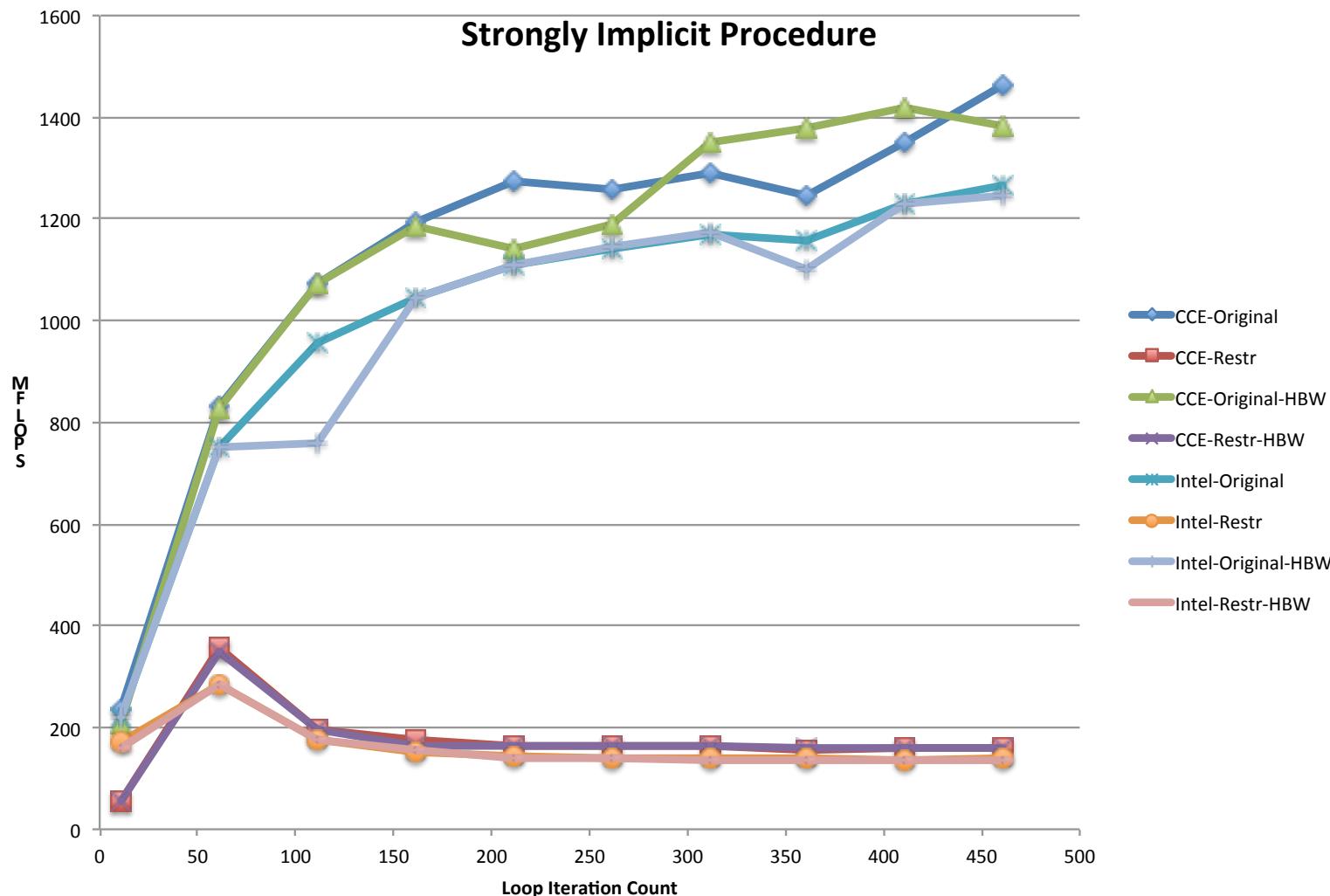


```
LOOP BEGIN at lp43140.f(95,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

    LOOP BEGIN at lp43140.f(108,11)
        remark #15300: LOOP WAS VECTORIZED
    LOOP END

    LOOP BEGIN at lp43140.f(108,11)
    <Remainder loop for vectorization>
        remark #15335: remainder loop was not vectorized: vectorization possible but seems inefficient.
            Use vector always directive or -vec-threshold0 to override
        remark #25456: Number of Array Refs Scalar Replaced In Loop: 4
    LOOP END
LOOP END
```

# Strongly Implicit Procedure



# Loop Carried Dependent Scalar



```
62.    1          C      THE ORIGINAL
63.    1
64.    1          PF = 0.0
65. + 1 2-----< DO 44030 I = 2, N
66.    1 2          AV = B(I) * RV
67.    1 2          PB = PF
68.    1 2          PF = C(I)
69.    1 2          IF ((D(I) + D(I+1)) .LT. 0.) PF = -C(I+1)
70.    1 2          AA = E(I) - E(I-1) + F(I) - F(I-1)
71.    1 2          1   + G(I) + G(I-1) - H(I) - H(I-1)
72.    1 2          BB = R(I) + S(I-1) + T(I) + T(I-1)
73.    1 2          1   - U(I) - U(I-1) + V(I) + V(I-1)
74.    1 2          2   - W(I) + W(I-1) - X(I) + X(I-1)
75.    1 2          A(I) = AV * (AA + BB + PF - PB + Y(I) - Z(I)) + A(I)
76.    1 2-----> 44030 CONTINUE
77.    1
```

# Loop Carried Dependent Scalar



```
98.    1      C      THE RESTRUCTURED
99.    1
100.   1      VPF(1) = 0.0
101.   1 Vr2---< DO 44031 I = 2, N
102.   1 Vr2      AV = B(I) * RV
103.   1 Vr2      VPF(I) = C(I)
104.   1 Vr2      IF ((D(I) + D(I+1)) .LT. 0.) VPF(I) = -C(I+1)
105.   1 Vr2      AA = E(I) - E(I-1) + F(I) - F(I-1)
106.   1 Vr2      1   + G(I) + G(I-1) - H(I) - H(I-1)
107.   1 Vr2      BB = R(I) + S(I-1) + T(I) + T(I-1)
108.   1 Vr2      1   - U(I) - U(I-1) + V(I) + V(I-1)
109.   1 Vr2      2   - W(I) + W(I-1) - X(I) + X(I-1)
110.   1 Vr2      A(I) = AV * (AA + BB + VPF(I) - VPF(I-1) + Y(I) - Z(I)) + A(I)
111.   1 Vr2---> 44031 CONTINUE 77.    1
```

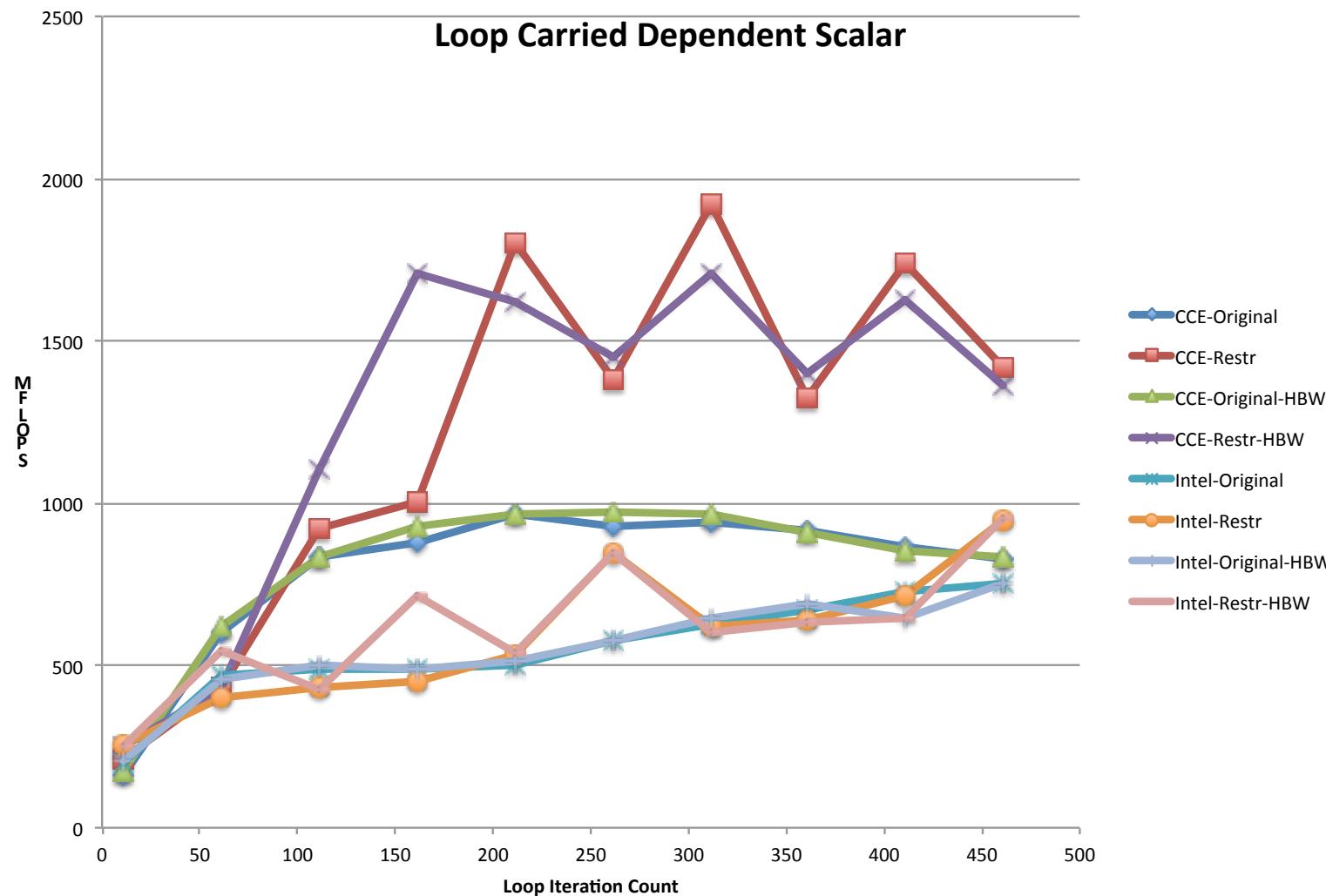
# Loop Carried Dependent Scalar



```
LOOP BEGIN at lp44030.f(65,10)
    remark #15344: loop was not vectorized: vector dependence prevents vectorization.
        First dependence is shown below. Use level 5 report for details
    remark #15346: vector dependence: assumed FLOW dependence between line 68 and line 67
    remark #25456: Number of Array Refs Scalar Replaced In Loop: 10
LOOP END
```

```
LOOP BEGIN at lp44030.f(101,10)
    remark #15344: loop was not vectorized: vector dependence prevents vectorization.
        First dependence is shown below. Use level 5 report for details
    remark #15346: vector dependence: assumed FLOW dependence between line 104 and line 110
    remark #25456: Number of Array Refs Scalar Replaced In Loop: 12
LOOP END
```

# Loop Carried Dependent Scalar



# Complex Reduction



```
61.    1      C      THE ORIGINAL
62.    1
63.    1 V---<    DO 44040 I = 2, N
64.    1 V          RR      = 1. / A(I,1)
65.    1 V          U       = A(I,2) * RR
66.    1 V          V       = A(I,3) * RR
67.    1 V          W       = A(I,4) * RR
68.    1 V          SNDSP   = SQRT (GD * (A(I,5) * RR + .5* (U*U + V*V + W*W)))
69.    1 V          SIGA   = ABS (XT + U*B(I) + V*C(I) + W*D(I))
70.    1 V          *      + SNDSP * SQRT (B(I)**2 + C(I)**2 + D(I)**2)
71.    1 V          SIGB   = ABS (YT + U*E(I) + V*F(I) + W*G(I))
72.    1 V          *      + SNDSP * SQRT (E(I)**2 + F(I)**2 + G(I)**2)
73.    1 V          SIGC   = ABS (ZT + U*H(I) + V*R(I) + W*S(I))
74.    1 V          *      + SNDSP * SQRT (H(I)**2 + R(I)**2 + S(I)**2)
75.    1 V          SIGABC = AMAX1 (SIGA, SIGB, SIGC)
76.    1 V          IF (SIGABC.GT.SIGMAX) THEN
77.    1 V          IMAX   = I
78.    1 V          SIGMAX = SIGABC
79.    1 V          ENDIF
80.    1 V---> 44040 CONTINUE
```

# Complex Reduction



```
99.    1      C      THE RESTRUCTURED
100.   1
101.   1 fV--<    DO 44041 I = 2, N
102.   1 fV          RR      = 1. / A(I,1)
103.   1 fV          U       = A(I,2) * RR
104.   1 fV          V       = A(I,3) * RR
105.   1 fV          W       = A(I,4) * RR
106.   1 fV          SNDSP   = SQRT (GD * (A(I,5) * RR + .5* (U*U + V*V + W*W)))
107.   1 fV          SIGA   = ABS (XT + U*B(I) + V*C(I) + W*D(I))
108.   1 fV          *      + SNDSP * SQRT (B(I)**2 + C(I)**2 + D(I)**2)
109.   1 fV          SIGB   = ABS (YT + U*E(I) + V*F(I) + W*G(I))
110.   1 fV          *      + SNDSP * SQRT (E(I)**2 + F(I)**2 + G(I)**2)
111.   1 fV          SIGC   = ABS (ZT + U*H(I) + V*R(I) + W*S(I))
112.   1 fV          *      + SNDSP * SQRT (H(I)**2 + R(I)**2 + S(I)**2)
113.   1 fV          VSIGABC(I) = AMAX1 (SIGA, SIGB, SIGC)
114.   1 fV--> 44041 CONTINUE
115.   1
116.   1 f---<    DO 44042 I = 2, N
117.   1 f          IF (VSIGABC(I) .GT. SIGMAX) THEN
118.   1 f          IMAX     = I
119.   1 f          SIGMAX   = VSIGABC(I)
120.   1 f          ENDIF
121.   1 f---> 44042 CONTINUE
```

# Complex Reduction



```
LOOP BEGIN at lp44040.f(63,10)
    remark #15344: loop was not vectorized: vector dependence prevents vectorization.
        First dependence is shown below. Use level 5 report for details
    remark #15346: vector dependence: assumed ANTI dependence between line 76 and line 78
    remark #25456: Number of Array Refs Scalar Replaced In Loop: 9
LOOP END

LOOP BEGIN at lp44040.f(101,10)
    remark #25045: Fused Loops: ( 101 116 )

    remark #15344: loop was not vectorized: vector dependence prevents vectorization.
        First dependence is shown below. Use level 5 report for details
    remark #15346: vector dependence: assumed ANTI dependence between line 117 and line 119
    remark #25456: Number of Array Refs Scalar Replaced In Loop: 9
LOOP END
```



# Matrix Multiply



```
44.      1
45.      1           C      THE ORIGINAL
46.      1
47.      + 1 2-----<      DO 44050 I = 1, N
48.      + 1 2 3-----<      DO 44050 J = 1, N
49.      1 2 3           A(I,J) = 0.0
50.      + 1 2 3 4-----<      DO 44050 K = 1, N
51.      1 2 3 4 A---<>      A(I,J) = A(I,J) + B(I,K) * C(K,J)
52.      1 2 3 4----->>> 44050 CONTINUE

77.      1
78.      1           C      THE RESTRUCTURED
79.      1
80.      + 1 2-----<      DO 44051 J = 1, N
81.      1 2 A-----<      DO 44051 I = 1, N
82.      1 2 A           A(I,J) = 0.0
83.      1 2 A----->> 44051 CONTINUE
84.      1
85.      + 1 2-----<      DO 44052 K = 1, N
86.      + 1 2 3-----<      DO 44052 J = 1, N
87.      + 1 2 3 4-----<      DO 44052 I = 1, N
88.      1 2 3 4 A---<>      A(I,J) = A(I,J) + B(I,K) * C(K,J)
89.      1 2 3 4----->>> 44052 CONTINUE
90.      1           C
```

# Matrix Multiply



```
LOOP BEGIN at lp44050.f(47,10)
  <Remainder, Distributed chunk2>
    remark #15542: loop was not vectorized: inner loop was already vectorized

  LOOP BEGIN at lp44050.f(48,11)
  <Distributed chunk2>
    remark #15542: loop was not vectorized: inner loop was already vectorized

  LOOP BEGIN at lp44050.f(50,12)
  <Peeled loop for vectorization>
    remark #15335: peel loop was not vectorized: vectorization possible but seems inefficient.
                  Use vector always directive or -vec-threshold0 to override
  LOOP END

  LOOP BEGIN at lp44050.f(50,12)
    remark #25085: Preprocess Loopnests: Moving Out Load and Store      [ lp44050.f(51,10) ]
    remark #15300: LOOP WAS VECTORIZED
  LOOP END

  LOOP BEGIN at lp44050.f(50,12)
  <Remainder loop for vectorization>
    remark #15301: REMAINDER LOOP WAS VECTORIZED
  LOOP END
  LOOP END
  LOOP END
```

# Matrix Multiply



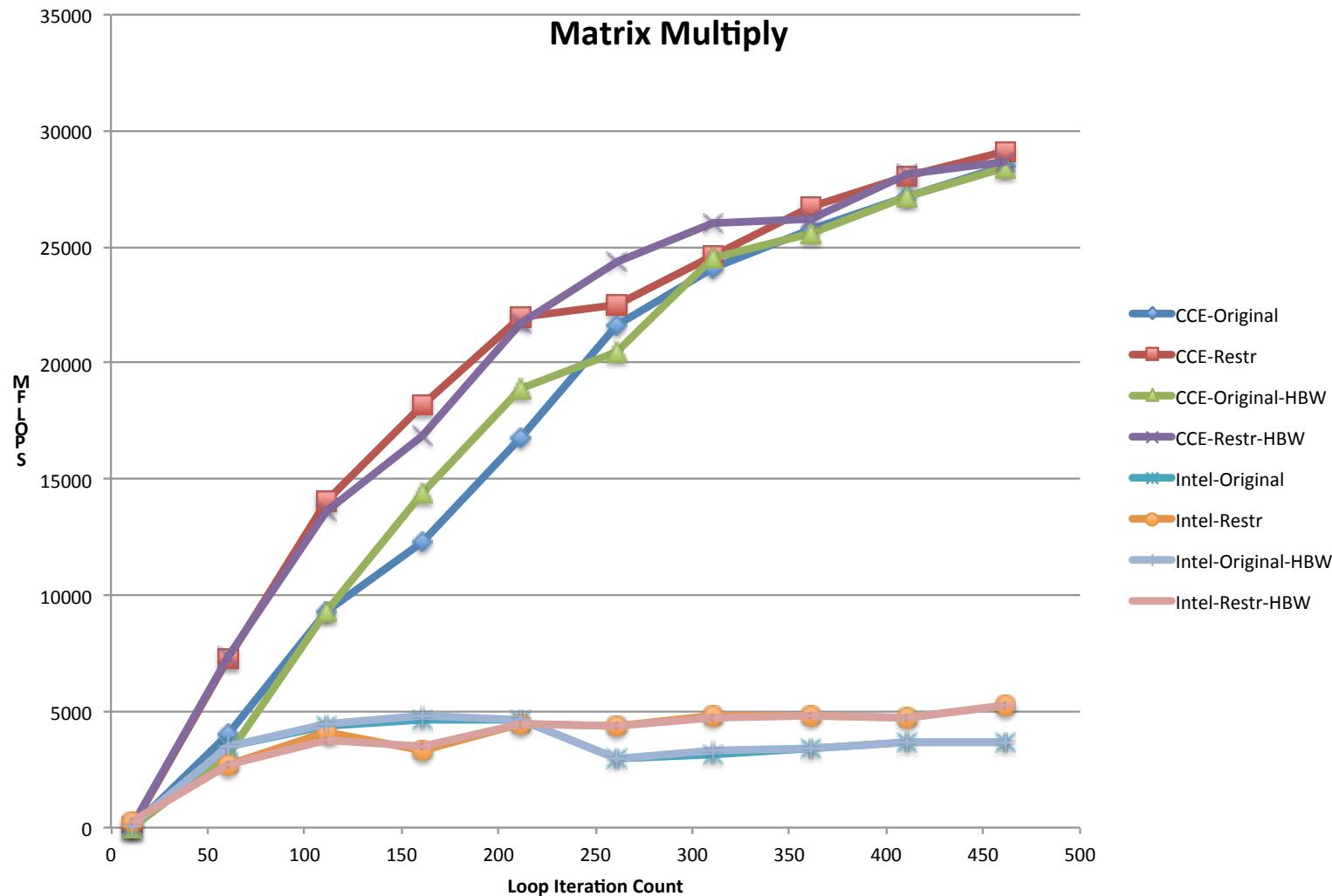
```
LOOP BEGIN at lp44050.f(80,10)
  remark #15542: loop was not vectorized: inner loop was already vectorized

  LOOP BEGIN at lp44050.f(81,11)
  <Peeled loop for vectorization>
    remark #15335: peel loop was not vectorized: vectorization possible but seems inefficient.
                  Use vector always directive or -vec-threshold0 to override
  LOOP END

  LOOP BEGIN at lp44050.f(81,11)
    remark #15300: LOOP WAS VECTORIZED
  LOOP END

  LOOP BEGIN at lp44050.f(81,11)
  <Remainder loop for vectorization>
    remark #15301: REMAINDER LOOP WAS VECTORIZED
  LOOP END
  LOOP END
```

# Matrix Multiply



# Vector Versus Reduction



```
62.    1 2          C      THE ORIGINAL
63.    1 2
64.    1 2
65. + 1 2 3-----<      J = 1
66.    1 2 3      DO 45011 I = 1, IM
67.    1 2 3      SUM1 = 0.0
68.    1 2 3      SUM2 = 0.0
69.    1 2 3      SUM4 = 0.0
70.    1 2 3 Vr2-----<      SUM5 = 0.0
71.    1 2 3 Vr2      DO 45010 K = 2, KM
72.    1 2 3 Vr2      KK = KM - K + 1
73.    1 2 3 Vr2      SUM1 = SUM1 + 4.0 * ( A(J+1,K,I,1) * A(J+1,K,I,6)
74.    1 2 3 Vr2      *           +     A(J+1,KK,I,1) * A(J+1,KK,I,6) )
75.    1 2 3 Vr2      *           -     A(J+2,K,I,1) * A(J+2,K,I,6)
76.    1 2 3 Vr2      *           -     A(J+2,KK,I,1) * A(J+2,KK,I,6)
77.    1 2 3 Vr2      SUM2 = SUM2 + 4.0 * ( A(J+1,K,I,2) * A(J+1,K,I,6)
78.    1 2 3 Vr2      *           +     A(J+1,KK,I,2) * A(J+1,KK,I,6) )
79.    1 2 3 Vr2      *           -     A(J+2,K,I,2) * A(J+2,K,I,6)
80.    1 2 3 Vr2      *           -     A(J+2,KK,I,2) * A(J+2,KK,I,6)
81.    1 2 3 Vr2      SUM4 = SUM4 + 4.0 * ( A(J+1,K,I,4) * A(J+1,K,I,6)
82.    1 2 3 Vr2      *           +     A(J+1,KK,I,4) * A(J+1,KK,I,6) )
83.    1 2 3 Vr2      *           -     A(J+2,K,I,4) * A(J+2,K,I,6)
84.    1 2 3 Vr2      *           -     A(J+2,KK,I,4) * A(J+2,KK,I,6)
85.    1 2 3 Vr2      SUM5 = SUM5 + 4.0 * ( A(J+1,K,I,5) * A(J+1,K,I,6)
86.    1 2 3 Vr2      *           +     A(J+1,KK,I,5) * A(J+1,KK,I,6) )
87.    1 2 3 Vr2      *           -     A(J+2,K,I,5) * A(J+2,K,I,6)
88.    1 2 3 Vr2      *           -     A(J+2,KK,I,5) * A(J+2,KK,I,6)
89.    1 2 3 Vr2      92.    1 2 3 Vr2-----> 45010      CONTINUE
90.    1 2 3 Vr2
91.    1 2 3 Vr2      93.    1 2 3
92.    1 2 3 Vr2-----<      DO 45011 K = 2, KM
93.    1 2 3      A(J,K,I,1) = SUM1 / (6.0 * (KM-2) * A(J,K,I,6))
94.    1 2 3 Vr2      A(J,K,I,2) = SUM2 / (6.0 * (KM-2) * A(J,K,I,6))
95.    1 2 3 Vr2      A(J,K,I,3) = 0.0
96.    1 2 3 Vr2      A(J,K,I,4) = SUM4 / (6.0 * (KM-2) * A(J,K,I,6))
97.    1 2 3 Vr2      A(J,K,I,5) = SUM5 / (6.0 * (KM-2) * A(J,K,I,6))
```

# Vector Versus Reduction



```

139.    1 2 iVr2-----<
140.    1 2 iVr2
141. + 1 2 iVr2 fi-----<
142.    1 2 iVr2 fi
143.    1 2 iVr2 fi
144.    1 2 iVr2 fi
145.    1 2 iVr2 fi
146.    1 2 iVr2 fi
147.    1 2 iVr2 fi
148.    1 2 iVr2 fi
149.    1 2 iVr2 fi
150.    1 2 iVr2 fi
151.    1 2 iVr2 fi
152.    1 2 iVr2 fi
153.    1 2 iVr2 fi
154.    1 2 iVr2 fi
155.    1 2 iVr2 fi
156.    1 2 iVr2 fi
157.    1 2 iVr2 fi
158.    1 2 iVr2 fi
159.    1 2 iVr2 fi
160.    1 2 iVr2 fi
161.    1 2 iVr2 fi
162.    1 2 iVr2 fi----->> 45013 CONTINUE
163.    1 2
164.    1 2 iVr2-----<
165. + 1 2 iVr2 fi-----<
166.    1 2 iVr2 fi
167.    1 2 iVr2 fi
168.    1 2 iVr2 fi
169.    1 2 iVr2 fi
170.    1 2 iVr2 fi
171.    1 2 iVr2 fi----->> 45014 CONTINUE

```

DO 45013 K = 2, KM  
 KK = KM - K + 1  
 DO 45013 I = 1, IM  
 VSUM1(I) = VSUM1(I) + 4.0 \* ( A(J+1,K,I,1) \* A(J+1,K,I,6)  
 \* + A(J+1,KK,I,1) \* A(J+1,KK,I,6) )  
 \* - A(J+2,K,I,1) \* A(J+2,K,I,6)  
 \* - A(J+2,KK,I,1) \* A(J+2,KK,I,6)  
 VSUM2(I) = VSUM2(I) + 4.0 \* ( A(J+1,K,I,2) \* A(J+1,K,I,6)  
 \* + A(J+1,KK,I,2) \* A(J+1,KK,I,6) )  
 \* - A(J+2,K,I,2) \* A(J+2,K,I,6)  
 \* - A(J+2,KK,I,2) \* A(J+2,KK,I,6)  
 VSUM4(I) = VSUM4(I) + 4.0 \* ( A(J+1,K,I,4) \* A(J+1,K,I,6)  
 \* + A(J+1,KK,I,4) \* A(J+1,KK,I,6) )  
 \* - A(J+2,K,I,4) \* A(J+2,K,I,6)  
 \* - A(J+2,KK,I,4) \* A(J+2,KK,I,6)  
 VSUM5(I) = VSUM5(I) + 4.0 \* ( A(J+1,K,I,5) \* A(J+1,K,I,6)  
 \* + A(J+1,KK,I,5) \* A(J+1,KK,I,6) )  
 \* - A(J+2,K,I,5) \* A(J+2,K,I,6)  
 \* - A(J+2,KK,I,5) \* A(J+2,KK,I,6)  
 DO 45014 K = 2, KM  
 DO 45014 I = 1, IM  
 A(J,K,I,1) = VSUM1(I) / (6.0 \* (KM-2) \* A(J,K,I,6))  
 A(J,K,I,2) = VSUM2(I) / (6.0 \* (KM-2) \* A(J,K,I,6))  
 A(J,K,I,3) = 0.0  
 A(J,K,I,4) = VSUM4(I) / (6.0 \* (KM-2) \* A(J,K,I,6))  
 A(J,K,I,5) = VSUM5(I) / (6.0 \* (KM-2) \* A(J,K,I,6))

# Vector Versus Reduction



```
LOOP BEGIN at lp45010.f(65,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

    LOOP BEGIN at lp45010.f(70,11)
        remark #15300: LOOP WAS VECTORIZED
    LOOP END

    LOOP BEGIN at lp45010.f(70,11)
    <Remainder loop for vectorization>
        remark #15301: REMAINDER LOOP WAS VECTORIZED
    LOOP END

    LOOP BEGIN at lp45010.f(94,11)
        remark #15300: LOOP WAS VECTORIZED
    LOOP END

    LOOP BEGIN at lp45010.f(94,11)
    <Remainder loop for vectorization>
        remark #15301: REMAINDER LOOP WAS VECTORIZED
    LOOP END
LOOP END
```

# Vector Versus Reduction



```
LOOP BEGIN at lp45010.f(139,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

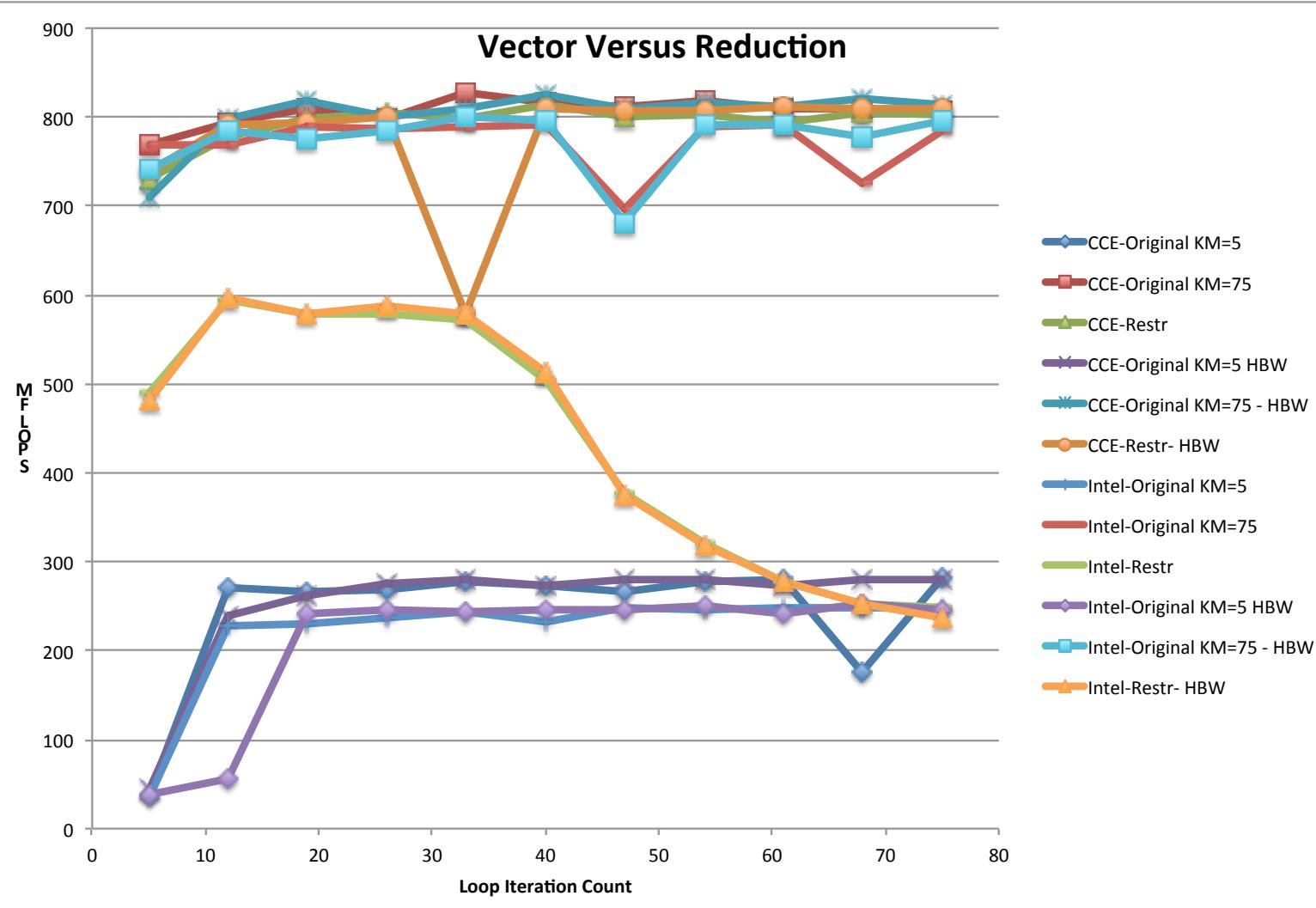
    LOOP BEGIN at lp45010.f(141,11)
        remark #15300: LOOP WAS VECTORIZED
    LOOP END

    LOOP BEGIN at lp45010.f(141,11)
    <Remainder loop for vectorization>
        remark #15301: REMAINDER LOOP WAS VECTORIZED
    LOOP END
LOOP END

LOOP BEGIN at lp45010.f(164,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

    LOOP BEGIN at lp45010.f(165,11)
        remark #15300: LOOP WAS VECTORIZED
    LOOP END

    LOOP BEGIN at lp45010.f(165,11)
    <Remainder loop for vectorization>
        remark #15301: REMAINDER LOOP WAS VECTORIZED
    LOOP END
LOOP END
```



# Rectangular Matrix Multiply



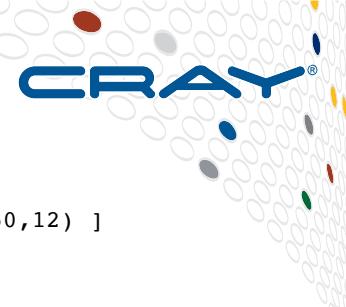
```
44.      1
45.      1          C      THE ORIGINAL
46.      1
47.      + 1 2-----<      DO 46020 I = 1,N
48.      + 1 2 3-----<      DO 46020 J = 1,4
49.      1 2 3          A(I,J) = 0.
50.      + 1 2 3 4-----<      DO 46020 K = 1,4
51.      1 2 3 4 A-->>      A(I,J) = A(I,J) + B(I,K) * C(K,J)
52.      1 2 3 4--->>> 46020 CONTINUE
```

# Rectangular Matrix Multiply



```
69.      1           C      THE RESTRUCTURED
70.      1
71.      1 V-----<    DO 46021 I = 1, N
72.      1 V           A(I,1) = B(I,1) * C(1,1) + B(I,2) * C(2,1)
73.      1 V           *      + B(I,3) * C(3,1) + B(I,4) * C(4,1)
74.      1 V           A(I,2) = B(I,1) * C(1,2) + B(I,2) * C(2,2)
75.      1 V           *      + B(I,3) * C(3,2) + B(I,4) * C(4,2)
76.      1 V           A(I,3) = B(I,1) * C(1,3) + B(I,2) * C(2,3)
77.      1 V           *      + B(I,3) * C(3,3) + B(I,4) * C(4,3)
78.      1 V           A(I,4) = B(I,1) * C(1,4) + B(I,2) * C(2,4)
79.      1 V           *      + B(I,3) * C(3,4) + B(I,4) * C(4,4)
80.      1 V-----> 46021 CONTINUE
```

# Rectangular Matrix Multiply



```
LOOP BEGIN at lp46020.f(47,10)
  remark #15541: outer loop was not auto-vectorized: consider using SIMD directive [ lp46020.f(50,12) ]
  remark #25456: Number of Array Refs Scalar Replaced In Loop: 16

LOOP BEGIN at lp46020.f(48,11)
  remark #15541: outer loop was not auto-vectorized: consider using SIMD directive [ lp46020.f(50,12) ]
  remark #25436: completely unrolled by 4

LOOP BEGIN at lp46020.f(50,12)
  remark #25085: Preprocess Loopnests: Moving Out Load and Store [ lp46020.f(51,10) ]
  remark #15335: loop was not vectorized: vectorization possible but seems inefficient. Use vector
always directive or -vec-threshold0 to override
  remark #25436: completely unrolled by 4
LOOP END

LOOP BEGIN at lp46020.f(50,12)
LOOP END

LOOP BEGIN at lp46020.f(50,12)
LOOP END

LOOP BEGIN at lp46020.f(50,12)
LOOP END
LOOP END
LOOP END
```

# Rectangular Matrix Multiply

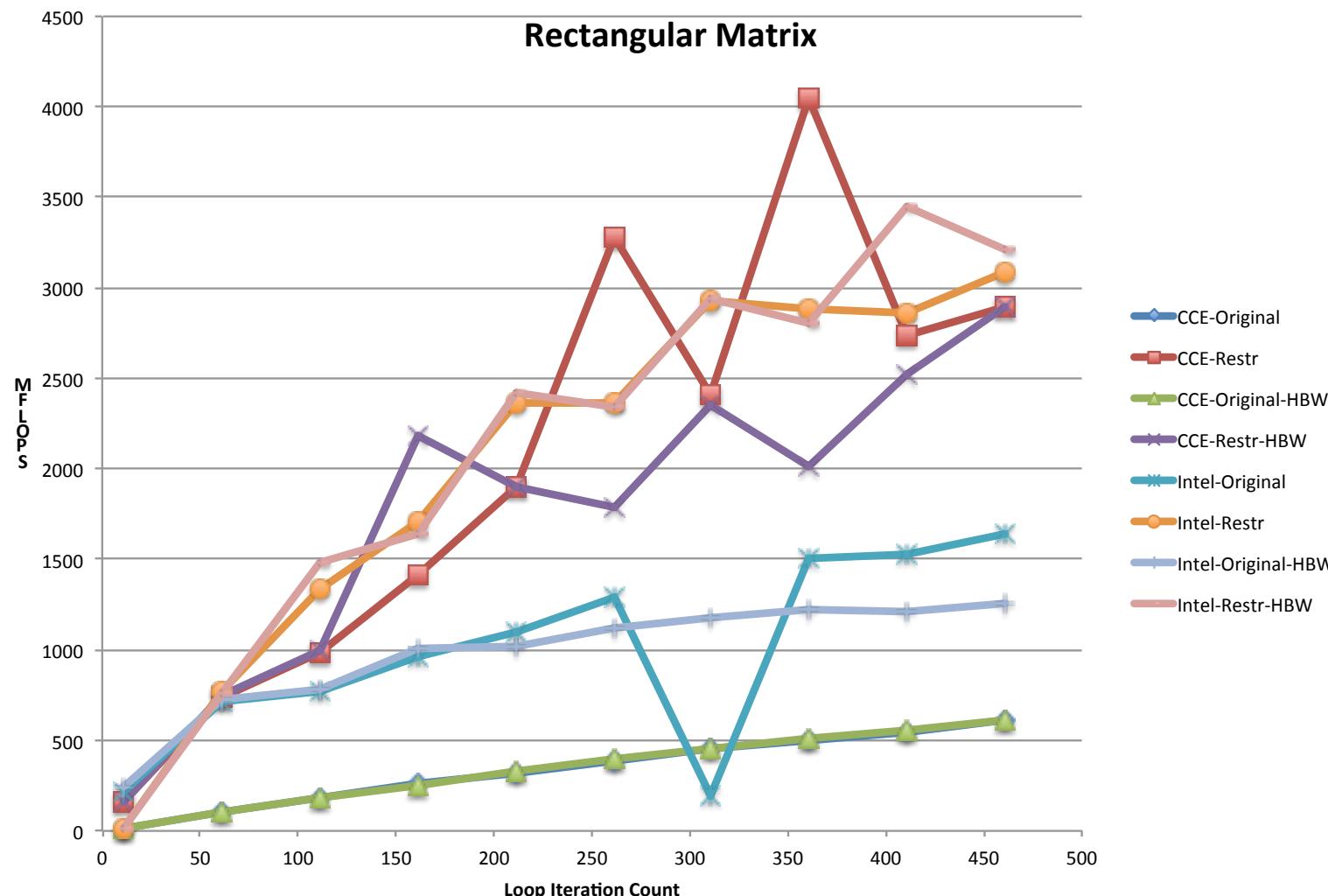


```
LOOP BEGIN at lp46020.f(71,10)
  <Peeled loop for vectorization>
    remark #15301: PEEL LOOP WAS VECTORIZED
  LOOP END

  LOOP BEGIN at lp46020.f(71,10)
    remark #15300: LOOP WAS VECTORIZED
    remark #25456: Number of Array Refs Scalar Replaced In Loop: 12
  LOOP END

  LOOP BEGIN at lp46020.f(71,10)
  <Remainder loop for vectorization>
    remark #15301: REMAINDER LOOP WAS VECTORIZED
  LOOP END
```

# Rectangular Matrix



# Optimized Matrix Multiply



```
42.      1           C      THE ORIGINAL
43.      1
44. + 1 2-----<      DO 46030 J = 1, N
45.   1 2 A-----<      DO 46030 I = 1, N
46.   1 2 A           A(I,J) = 0.
47.   1 2 A----->> 46030 CONTINUE
48.
49.   1
50. + 1 2 3-----<      DO 46031   K = 1, N
51. + 1 2 3 4-----<      DO 46031   J = 1, N
52.   1 2 3 4 A---->>      A(I,J) = A(I,J) + B(I,K) * C(K,J)
53.   1 2 3 4----->>> 46031 CONTINUE
```

# Optimized Matrix Multiply



```
70.      1           C      THE RESTRUCTURED
71.      1
72. + 1 2-----<      DO 46032 J = 1, N
73. 1 2 A-----<      DO 46032 I = 1, N
74. 1 2 A          A(I,J)=0.
75. 1 2 A----->> 46032 CONTINUE
76.      1           C
77. + 1 2-----<      DO 46033 K = 1, N-5, 6
78. + 1 2 r4-----<      DO 46033 J = 1, N
79. 1 2 r4 Vr2---<      DO 46033 I = 1, N
80. 1 2 r4 Vr2          A(I,J) = A(I,J) + B(I,K ) * C(K ,J)
81. 1 2 r4 Vr2          *          + B(I,K+1) * C(K+1,J)
82. 1 2 r4 Vr2          *          + B(I,K+2) * C(K+2,J)
83. 1 2 r4 Vr2          *          + B(I,K+3) * C(K+3,J)
84. 1 2 r4 Vr2          *          + B(I,K+4) * C(K+4,J)
85. 1 2 r4 Vr2          *          + B(I,K+5) * C(K+5,J)
86. 1 2 r4 Vr2-->>> 46033 CONTINUE
87.      1           C
88. + 1 2-----<      DO 46034 KK = K, N
89. + 1 2 3-----<      DO 46034 J = 1, N
90. + 1 2 3 4-----<      DO 46034 I = 1, N
91. 1 2 3 4 A---<>      A(I,J) = A(I,J) + B(I,KK) * C(KK ,J)
92. 1 2 3 4----->>> 46034 CONTINUE
```

# Optimized Matrix Multiply



```
LOOP BEGIN at lp46030.f(30,10)
  remark #15542: loop was not vectorized: inner loop was already vectorized

LOOP BEGIN at lp46030.f(44,10)
  remark #15542: loop was not vectorized: inner loop was already vectorized

LOOP BEGIN at lp46030.f(45,11)
<Peeled loop for vectorization>
  remark #15335: peel loop was not vectorized: vectorization possible but seems inefficient. Use vector
always directive or -vec-threshold0 to override
  LOOP END

LOOP BEGIN at lp46030.f(45,11)
  remark #15300: LOOP WAS VECTORIZED
  LOOP END

LOOP BEGIN at lp46030.f(45,11)
<Remainder loop for vectorization>
  remark #15301: REMAINDER LOOP WAS VECTORIZED
  LOOP END
LOOP END
```



# Optimized Matrix Multiply

```
LOOP BEGIN at lp46030.f(72,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

    LOOP BEGIN at lp46030.f(73,11)
    <Peeled loop for vectorization>
        remark #15335: peel loop was not vectorized: vectorization possible but seems inefficient.
                        Use vector always directive or -vec-threshold0 to override
    LOOP END

    LOOP BEGIN at lp46030.f(73,11)
        remark #15300: LOOP WAS VECTORIZED
    LOOP END

    LOOP BEGIN at lp46030.f(73,11)
    <Remainder loop for vectorization>
        remark #15301: REMAINDER LOOP WAS VECTORIZED
    LOOP END
LOOP END

LOOP BEGIN at lp46030.f(77,10)
    remark #15542: loop was not vectorized: inner loop was already vectorized

    LOOP BEGIN at lp46030.f(78,11)
        remark #15542: loop was not vectorized: inner loop was already vectorized

    LOOP BEGIN at lp46030.f(79,12)
        remark #15300: LOOP WAS VECTORIZED
    LOOP END

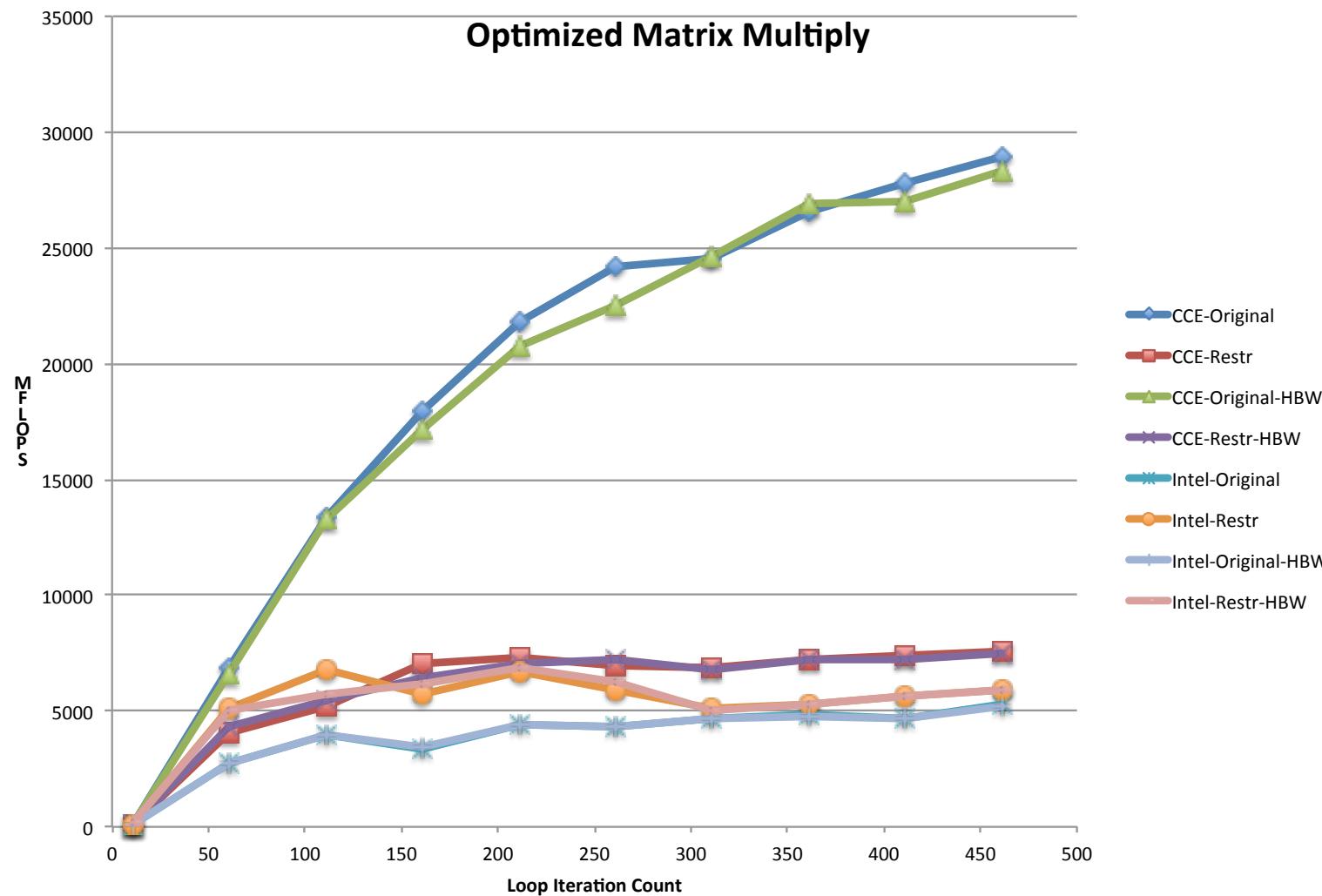
    LOOP BEGIN at lp46030.f(79,12)
    <Remainder loop for vectorization>
        remark #15301: REMAINDER LOOP WAS VECTORIZED
    LOOP END
LOOP END
LOOP END
LOOP END
```

C O M P U T E

S T O R E

A N A L Y Z E

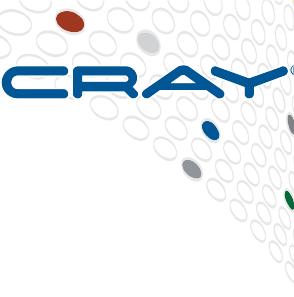
# Optimized Matrix Multiply



# IF tests on Loop Indices



```
53.    1           C      THE ORIGINAL
54.    1
55. + 1 i-----<      DO 47020   J = 1, JMAX
56. + 1 i i-----<      DO 47020   K = 1, KMAX
57. + 1 i i V-----<    DO 47020   I = 1, IMAX
58.    1 i i V          JP        = J + 1
59.    1 i i V          JR        = J - 1
60.    1 i i V          KP        = K + 1
61.    1 i i V          KR        = K - 1
62.    1 i i V          IP        = I + 1
63.    1 i i V          IR        = I - 1
64.    1 i i V          IF (J .EQ. 1) GO TO 50
65.    1 i i V          IF( J .EQ. JMAX) GO TO 51
66.    1 i i V          XJ = ( A(I,JP,K) - A(I,JR,K) ) * DA2
67.    1 i i V          YJ = ( B(I,JP,K) - B(I,JR,K) ) * DA2
68.    1 i i V          ZJ = ( C(I,JP,K) - C(I,JR,K) ) * DA2
69.    1 i i V          GO TO 70
70.    1 i i V          50     J1 = J + 1
71.    1 i i V          J2 = J + 2
72.    1 i i V          XJ = (-3. * A(I,J,K) + 4. * A(I,J1,K) - A(I,J2,K) ) * DA2
73.    1 i i V          YJ = (-3. * B(I,J,K) + 4. * B(I,J1,K) - B(I,J2,K) ) * DA2
74.    1 i i V          ZJ = (-3. * C(I,J,K) + 4. * C(I,J1,K) - C(I,J2,K) ) * DA2
75.    1 i i V          GO TO 70
76.    1 i i V          51     J1 = J - 1
77.    1 i i V          J2 = J - 2
78.    1 i i V          XJ = ( 3. * A(I,J,K) - 4. * A(I,J1,K) + A(I,J2,K) ) * DA2
79.    1 i i V          YJ = ( 3. * B(I,J,K) - 4. * B(I,J1,K) + B(I,J2,K) ) * DA2
80.    1 i i V          ZJ = ( 3. * C(I,J,K) - 4. * C(I,J1,K) + C(I,J2,K) ) * DA2
81.    1 i i V          70     CONTINUE
82.    1 i i V          IF (K .EQ. 1) GO TO 52
```



# IF tests on Loop Indices

```
83. 1 i i v
84. 1 i i v
85. 1 i i v
86. 1 i i v
87. 1 i i v
88. 1 i i v
89. 1 i i v
90. 1 i i v
91. 1 i i v
92. 1 i i v
93. 1 i i v
94. 1 i i v
95. 1 i i v
96. 1 i i v
97. 1 i i v
98. 1 i i v
99. 1 i i v
100. 1 i i v
101. 1 i i v
102. 1 i i v
103. 1 i i v
104. 1 i i v
105. 1 i i v
106. 1 i i v
107. 1 i i v
108. 1 i i v
109. 1 i i v
110. 1 i i v
111. 1 i i v
112. 1 i i v
113. 1 i i v
114. 1 i i v
115. 1 i i v
116. 1 i i v
117. 1 i i v
118. 1 i i v
119. 1 i i v
120. 1 i i v
121. 1 i i v----->>> 47020 CONTINUE
```

IF (K .EQ. KMAX) GO TO 53  
XX = ( A(I,J,KP) - A(I,J,KR) ) \* DB2  
YK = ( B(I,J,KP) - B(I,J,KR) ) \* DB2  
ZK = ( C(I,J,KP) - C(I,J,KR) ) \* DB2  
GO TO 71

52      K1 = K + 1  
K2 = K + 2  
XX = (-3. \* A(I,J,K) + 4. \* A(I,J,K1) - A(I,J,K2) ) \* DB2  
YK = (-3. \* B(I,J,K) + 4. \* B(I,J,K1) - B(I,J,K2) ) \* DB2  
ZK = (-3. \* C(I,J,K) + 4. \* C(I,J,K1) - C(I,J,K2) ) \* DB2  
GO TO 71

53      K1 = K - 1  
K2 = K - 2  
XX = ( 3. \* A(I,J,K) - 4. \* A(I,J,K1) + A(I,J,K2) ) \* DB2  
YK = ( 3. \* B(I,J,K) - 4. \* B(I,J,K1) + B(I,J,K2) ) \* DB2  
ZK = ( 3. \* C(I,J,K) - 4. \* C(I,J,K1) + C(I,J,K2) ) \* DB2

71      CONTINUE  
IF (I .EQ. 1)      GO TO 54  
IF (I .EQ. IMAX) GO TO 55  
  XI = ( A(IP,J,K) - A(IR,J,K) ) \* DC2  
  YI = ( B(IP,J,K) - B(IR,J,K) ) \* DC2  
  ZI = ( C(IP,J,K) - C(IR,J,K) ) \* DC2  
  GO TO 60

54      I1 = I + 1  
I2 = I + 2  
  XI = (-3. \* A(I,J,K) + 4. \* A(I1,J,K) - A(I2,J,K) ) \* DC2  
  YI = (-3. \* B(I,J,K) + 4. \* B(I1,J,K) - B(I2,J,K) ) \* DC2  
  ZI = (-3. \* C(I,J,K) + 4. \* C(I1,J,K) - C(I2,J,K) ) \* DC2  
  GO TO 60

55      I1 = I - 1  
I2 = I - 2  
  XI = ( 3. \* A(I,J,K) - 4. \* A(I1,J,K) + A(I2,J,K) ) \* DC2  
  YI = ( 3. \* B(I,J,K) - 4. \* B(I1,J,K) + B(I2,J,K) ) \* DC2  
  ZI = ( 3. \* C(I,J,K) - 4. \* C(I1,J,K) + C(I2,J,K) ) \* DC2  
  GO TO 60

60      CONTINUE

---

DINV = XI \* YK \* ZI + YJ \* ZK \* XI + ZJ \* XK \* YI  
      - XJ \* ZK \* YI - YJ \* XK \* ZI - ZJ \* YK \* XI  
D(I,J,K) = 1. / (DINV + 1.E-20)

COMPUTE | STORE | ANALYZE



# IF tests on Loop Indices

```
142.    1          C      THE RESTRUCTURED
143.    1
144. + 1 2-----< DO 47029 J = 1, JMAX
145. + 1 2 3-----< DO 47029 K = 1, KMAX
146.    1 2 3
147.    1 2 3          IF(J.EQ.1)THEN
148.    1 2 3
149.    1 2 3          J1      = 2
150.    1 2 3          J2      = 3
151.    1 2 3 Vr2-----< DO 47021 I = 1, IMAX
152.    1 2 3 Vr2          VAJ(I) = (-3. * A(I,J,K) + 4. * A(I,J1,K) - A(I,J2,K) ) * DA2
153.    1 2 3 Vr2          VBJ(I) = (-3. * B(I,J,K) + 4. * B(I,J1,K) - B(I,J2,K) ) * DA2
154.    1 2 3 Vr2          VCJ(I) = (-3. * C(I,J,K) + 4. * C(I,J1,K) - C(I,J2,K) ) * DA2
155.    1 2 3 Vr2-----> 47021 CONTINUE
156.    1 2 3
157.    1 2 3          ELSE IF(J.NE.JMAX) THEN
158.    1 2 3
159.    1 2 3          JP      = J+1
160.    1 2 3          JR      = J-1
161.    1 2 3 Vr2-----< DO 47022 I = 1, IMAX
162.    1 2 3 Vr2          VAJ(I) = ( A(I,JP,K) - A(I,JR,K) ) * DA2
163.    1 2 3 Vr2          VBJ(I) = ( B(I,JP,K) - B(I,JR,K) ) * DA2
164.    1 2 3 Vr2          VCJ(I) = ( C(I,JP,K) - C(I,JR,K) ) * DA2
165.    1 2 3 Vr2-----> 47022 CONTINUE
166.    1 2 3
167.    1 2 3          ELSE
168.    1 2 3
169.    1 2 3          J1      = JMAX-1
170.    1 2 3          J2      = JMAX-2
171.    1 2 3 Vr2-----< DO 47023 I = 1, IMAX
172.    1 2 3 Vr2          VAJ(I) = ( 3. * A(I,J,K) - 4. * A(I,J1,K) + A(I,J2,K) ) * DA2
173.    1 2 3 Vr2          VBJ(I) = ( 3. * B(I,J,K) - 4. * B(I,J1,K) + B(I,J2,K) ) * DA2
174.    1 2 3 Vr2          VCJ(I) = ( 3. * C(I,J,K) - 4. * C(I,J1,K) + C(I,J2,K) ) * DA2
175.    1 2 3 Vr2-----> 47023 CONTINUE
176.    1 2 3
177.    1 2 3          ENDIF
178.    1 2 3
179.    1 2 3          IF(K.EQ.1) THEN
180.    1 2 3
181.    1 2 3          K1      = 2
182.    1 2 3          K2      = 3
183.    1 2 3 Vr2-----< DO 47024 I = 1, IMAX
184.    1 2 3 Vr2          VAK(I) = (-3. * A(I,J,K) + 4. * A(I,J,K1) - A(I,J,K2) ) * DB2
185.    1 2 3 Vr2          VBK(I) = (-3. * B(I,J,K) + 4. * B(I,J,K1) - B(I,J,K2) ) * DB2
186.    1 2 3 Vr2          VCK(I) = (-3. * C(I,J,K) + 4. * C(I,J,K1) - C(I,J,K2) ) * DB2
187.    1 2 3 Vr2-----> 47024 CONTINUE
```

C O M P U T E

S T O R E

A N A L Y Z E

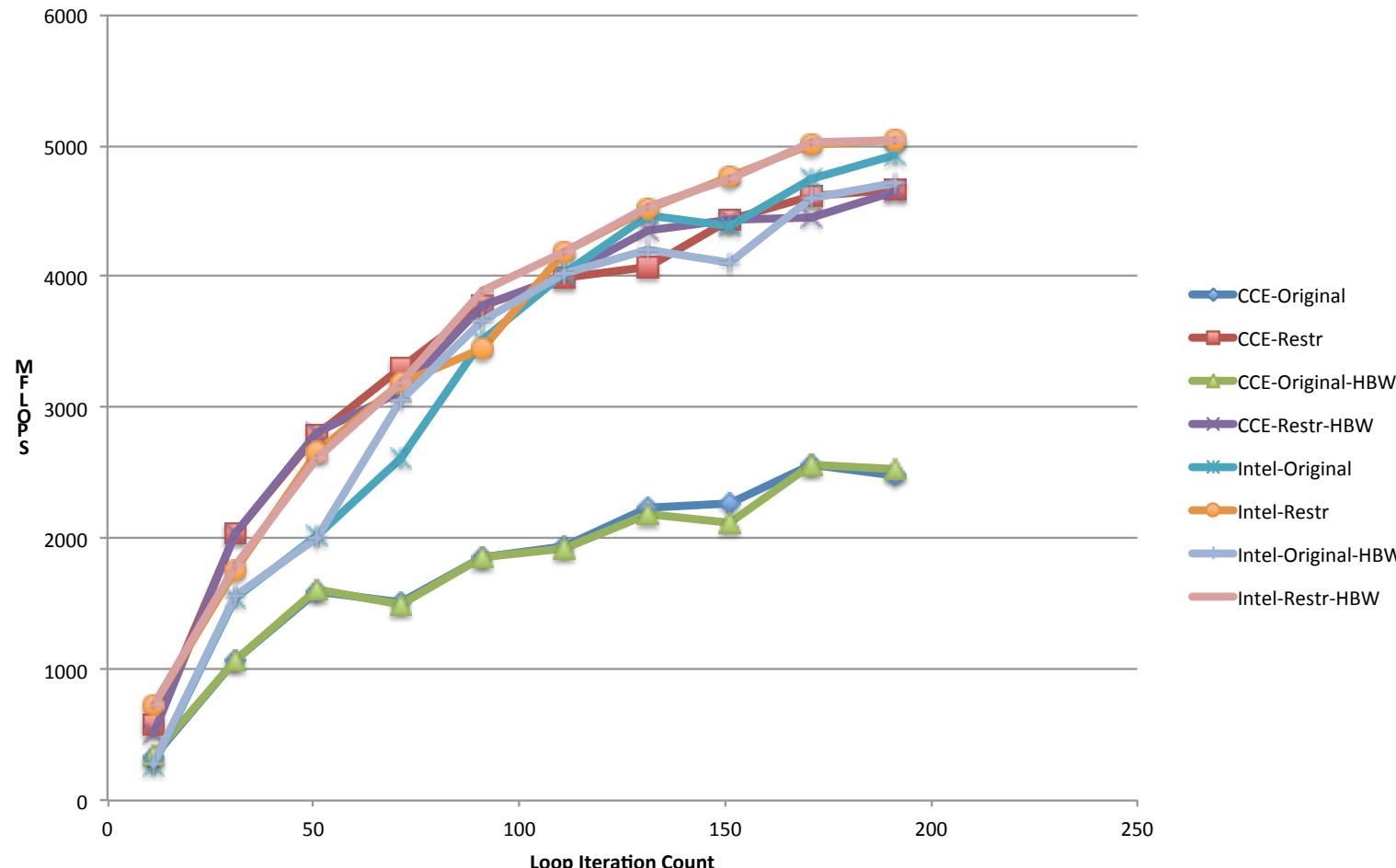
# IF tests on Loop Indices

```

186. 1 2 3
187. 1 2 3 ELSE IF (K.NE.KMAX) THEN
190. 1 2 3
191. 1 2 3 KP = K + 1
192. 1 2 3 KR = K - 1
193. 1 2 3 Vr2----< DO 47025 I = 1, IMAX
194. 1 2 3 Vr2 VAK(I) = ( A(I,J,KP) - A(I,J,KR) ) * DB2
195. 1 2 3 Vr2 VBK(I) = ( B(I,J,KP) - B(I,J,KR) ) * DB2
196. 1 2 3 Vr2 VCK(I) = ( C(I,J,KP) - C(I,J,KR) ) * DB2
197. 1 2 3 Vr2----> 47025 CONTINUE
198. 1 2 3
199. 1 2 3 ELSE
200. 1 2 3
201. 1 2 3 K1 = KMAX - 1
202. 1 2 3 K2 = KMAX - 2
203. 1 2 3 Vr2----< DO 47026 I = 1, IMAX
204. 1 2 3 Vr2 VAK(I) = ( 3. * A(I,J,K) - 4. * A(I,J,K1) + A(I,J,K2) ) * DB2
205. 1 2 3 Vr2 VBK(I) = ( 3. * B(I,J,K) - 4. * B(I,J,K1) + B(I,J,K2) ) * DB2
206. 1 2 3 Vr2 VCK(I) = ( 3. * C(I,J,K) - 4. * C(I,J,K1) + C(I,J,K2) ) * DB2
207. 1 2 3 Vr2----> 47026 CONTINUE
208. 1 2 3 ENDIF
209. 1 2 3
210. 1 2 3 I = 1
211. 1 2 3 I1 = 2
212. 1 2 3 I2 = 3
213. 1 2 3 VAI(I) = ( -3. * A(I,J,K) + 4. * A(I1,J,K) - A(I2,J,K) ) * DC2
214. 1 2 3 VBI(I) = ( -3. * B(I,J,K) + 4. * B(I1,J,K) - B(I2,J,K) ) * DC2
215. 1 2 3 VCI(I) = ( -3. * C(I,J,K) + 4. * C(I1,J,K) - C(I2,J,K) ) * DC2
216. 1 2 3
217. 1 2 3 Vr2----< DO 47027 I = 2, IMAX-1
218. 1 2 3 Vr2 IP = I + 1
219. 1 2 3 IR = I - 1
220. 1 2 3 Vr2 VAI(I) = ( A(IP,J,K) - A(IR,J,K) ) * DC2
221. 1 2 3 Vr2 VBI(I) = ( B(IP,J,K) - B(IR,J,K) ) * DC2
222. 1 2 3 Vr2 VCI(I) = ( C(IP,J,K) - C(IR,J,K) ) * DC2
223. 1 2 3 Vr2----> 47027 CONTINUE
224. 1 2 3
225. 1 2 3 I = IMAX
226. 1 2 3 I1 = IMAX - 1
227. 1 2 3 I2 = IMAX - 2
228. 1 2 3 VAI(I) = ( 3. * A(I,J,K) - 4. * A(I1,J,K) + A(I2,J,K) ) * DC2
229. 1 2 3 VBI(I) = ( 3. * B(I,J,K) - 4. * B(I1,J,K) + B(I2,J,K) ) * DC2
230. 1 2 3 VCI(I) = ( 3. * C(I,J,K) - 4. * C(I1,J,K) + C(I2,J,K) ) * DC2
231. 1 2 3
232. 1 2 3 Vr2----< DO 47028 I = 1, IMAX
233. 1 2 3 Vr2 DINV = VAJ(I) * VBK(I) * VCI(I) + VBJ(I) * VCK(I) * VAI(I)
234. 1 2 3 Vr2 1 + VCJ(I) * VAK(I) * VBI(I) - VAJ(I) * VCK(I) * VBI(I)
235. 1 2 3 Vr2 2 - VBJ(I) * VAK(I) * VCI(I) - VCJ(I) * VBK(I) * VAI(I)
236. 1 2 3 Vr2 D(I,J,K) = 1. / (DINV + 1.E-20)
237. 1 2 3 Vr2----> 47028 CONTINUE
238. 1 2 3----->> 47029 CONTINUE

```

## Tests on IF Tests



# Complex Decision



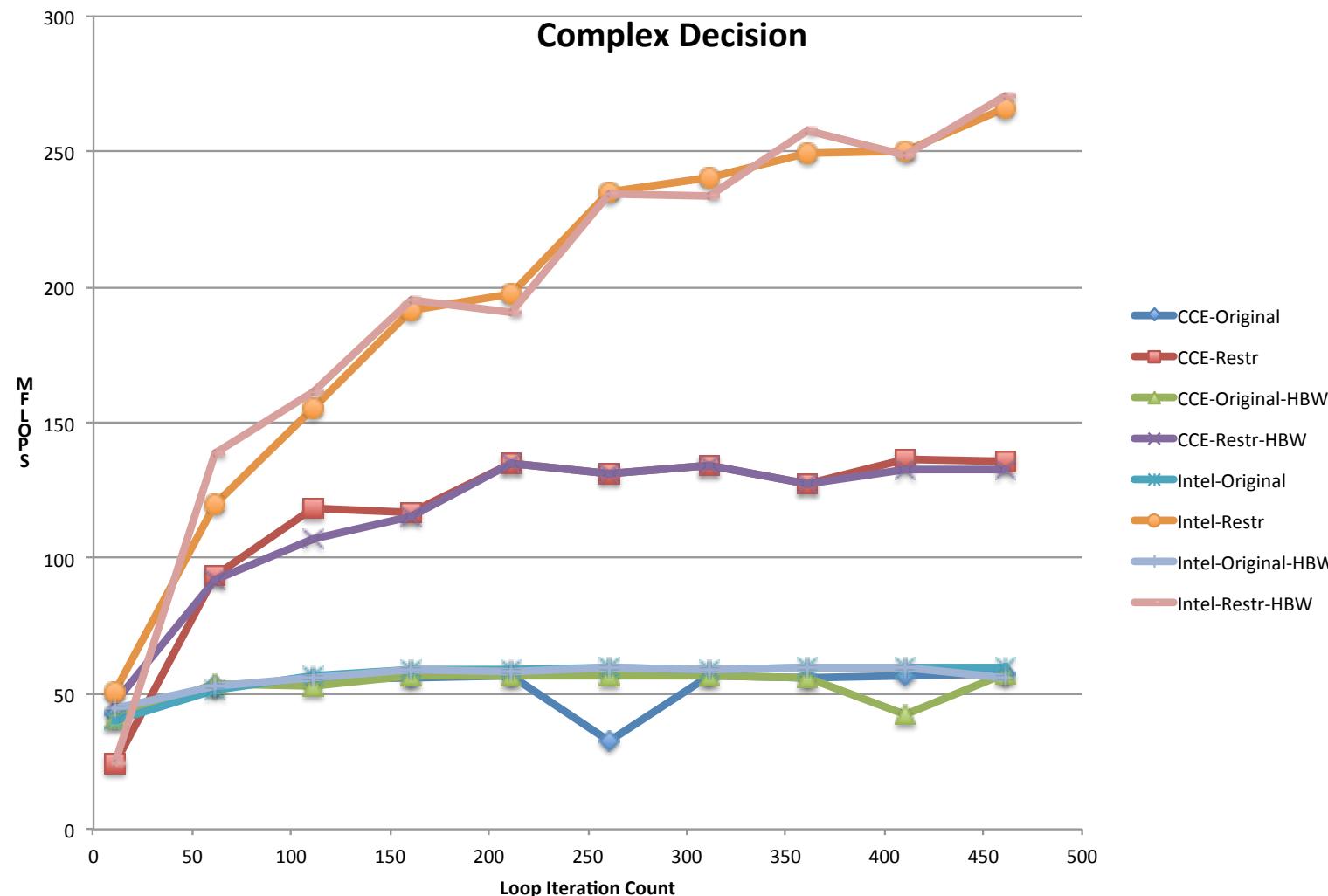
```
80.      1           C      THE ORIGINAL
81.      1
82.      1           SUM = 0.0
83. + 1 i-----< DO 47080 J = 2, JMAX
84. + 1 i i-----< DO 47080 I = 2, N
85.      1 i i       IF (I .EQ. N) GO TO 47080
86.      1 i i       IF (A(1,J) .LT. B(1,I)) GO TO 47080
87.      1 i i       IF (A(1,1) .GT. B(1,I)) GO TO 47080
88.      1 i i       IF (A(1,J) .GE. B(1,I+1) .AND. I .NE. N) GO TO 500
89.      1 i i       IF (J.EQ.1) GO TO 47080
90.      1 i i       IF (A(1,J-1) .LT. B(1,I-1) .AND. I*J .NE. 1) GO TO 500
91.      1 i i       IF (A(1,J-1) .LT. B(1,I)) GO TO 47080
92.      1 i i           500   CONTINUE
93.      1 i i       P1     = C(1,I-1)
94.      1 i i       P2     = D(I-1)
95.      1 i i       DD     = B(1,I) - B(1,I-1)
96.      1 i i       P3     = (3.0 * E(I) - 2.0 * P2 - D(I)) / DD
97.      1 i i       P4     = ( P2 + D(I) - 2.0 * E(I) ) / DD**2
98.      1 i i       SUMND = DD * (P1           + DD * (P2 / 2.
99.      1 i i           *           + DD * (P3 / 3. + DD * P4 / 4.) ) )
100.     1 i i           SUM   = SUM + SUMND
101.    1 i i--->> 47080 CONTINUE
```

# Complex Decision



```
121.    1          C      THE RESTRUCTURED
122.    1
123.    1          SUM = 0.0
124.    1 i-----<    DO 47081 J = 2, JMAX
125.    1 i iV---<    DO 47081 I = 2, N-1
126.    1 i iV
127.    1 i iV          LOG1 = A(1,J) .GE. B(1,I)
128.    1 i iV          LOG2 = A(1,1) .LE. B(1,I)
129.    1 i iV          LOG3 = A(1,J) .GE. B(1,I+1)
130.    1 i iV          LOG4 = J .NE. 1
131.    1 i iV          LOG5 = A(1,J-1) .LT. B(1,I-1)
132.    1 i iV          LOG6 = A(1,J-1) .GE. B(1,I)
133.    1 i iV          LOG7 = LOG1 .AND. LOG2 .AND. LOG3 .OR.
134.    1 i iV          *      LOG1 .AND. LOG2 .AND. LOG4 .AND. LOG5 .OR.
135.    1 i iV          *      LOG1 .AND. LOG2 .AND. LOG4 .AND. LOG6
136.    1 i iV          P1   = C(1,I-1)
137.    1 i iV          P2   = D(I-1)
138.    1 i iV          DD   = B(1,I) - B(1,I-1)
139.    1 i iV          IF(.NOT. LOG7) DD=1.0
140.    1 i iV          P3   = (3.0 * E(I) - 2.0 * P2 - D(I)) / DD
141.    1 i iV          P4   = ( P2 + D(I) - 2.0 * E(I) ) / DD**2
142.    1 i iV          SUMND = 0.0
143.    1 i iV          IF(LOG7) SUMND = DD * (P1      + DD * (P2 / 2.
144.    1 i iV          *           + DD * (P3 / 3. + DD * P4 / 4.) ) )
145.    1 i iV          SUM   = SUM + SUMND
146.    1 i iV-->> 47081 CONTINUE
```

# Complex Decision



# VH1 Profile



69.7%	275.784732	--	--	9,830,702.0	USER
<hr/>					
19.1%	75.655916	2.409537	3.1%	409,600.0	riemann_
10.1%	39.933326	3.687220	8.5%	3,686,400.0	parabola_
7.8%	30.814069	1.117453	3.5%	100.0	sweeypy_
6.9%	27.441332	1.908996	6.5%	409,600.0	remap_
5.8%	22.891882	0.595616	2.5%	50.0	sweepz_
3.9%	15.461957	1.012099	6.1%	50.0	sweepx1_
3.5%	14.017892	2.581348	15.6%	819,200.0	paraset_
2.9%	11.384113	1.949496	14.6%	409,600.0	evolve_
2.2%	8.649484	0.464873	5.1%	409,600.0	ppmlr_
2.1%	8.169585	0.321632	3.8%	409,600.0	states_
1.5%	6.071514	0.543860	8.2%	1,228,800.0	volume_
1.3%	5.015920	0.372648	6.9%	409,600.0	flatten_
1.2%	4.799249	0.146185	3.0%	50.0	sweepx2_
<hr/>					

C O M P U T E

S T O R E

A N A L Y Z E

# Important Loop in Riemann – VH1



```
63. + 1----< do l = lmin, lmax
64. + 1 2--<   do n = 1, 12
65.    1 2          pmold(l) = pmid(l)
66.    1 2          wlft (l) = 1.0 + gamfac1*(pmid(l) - plft(l)) * plfti(l)
67.    1 2          wrgh (l) = 1.0 + gamfac1*(pmid(l) - prgh(l)) * prghi(l)
68.    1 2          wlft (l) = clft(l) * sqrt(wlft(l))
69.    1 2          wrgh (l) = crgh(l) * sqrt(wrgh(l))
70.    1 2          zlft (l) = 4.0 * vlft(l) * wlft(l) * wlft(l)
71.    1 2          zrgh (l) = 4.0 * vrgh(l) * wrgh(l) * wrgh(l)
72.    1 2          zlft (l) = -zlft(l) * wlft(l)/(zlft(l) - gamfac2*(pmid(l) - plft(l)))
73.    1 2          zrgh (l) = zrgh(l) * wrgh(l)/(zrgh(l) - gamfac2*(pmid(l) - prgh(l)))
74.    1 2          umidl(l) = ulft(l) - (pmid(l) - plft(l)) / wlft(l)
75.    1 2          umidr(l) = urgh(l) + (pmid(l) - prgh(l)) / wrgh(l)
76.    1 2          pmid (l) = pmid(l) + (umidr(l) - umidl(l))*(zlft(l) * zrgh(l)) / (zrgh(l)
77.    1 2          pmid (l) = max(smallp,pmid(l))
78.    1 2          if (abs(pmid(l)-pmold(l))/pmid(l) < tol ) exit
79.    1 2-->  enddo
80. 1----> enddo
```

# Important Loop in Riemann – VH1



```
70. + 1-----< do n = 1, 12
71. + 1           if(any(not_converge(lmin:lmax)))then
72.   1 Vr2--<  do l = lmin, lmax
73.   1 Vr2           if (not_converge(l))then
74.   1 Vr2           pmold(l) = pmid(l)
75.   1 Vr2           wlft (l) = 1.0 + gamfac1*(pmid(l) - plft(l)) * plfti(l)
76.   1 Vr2           wrgh (l) = 1.0 + gamfac1*(pmid(l) - prgh(l)) * prghi(l)
77.   1 Vr2           wlft (l) = clft(l) * sqrt(wlft(l))
78.   1 Vr2           wrgh (l) = crgh(l) * sqrt(wrgh(l))
79.   1 Vr2           zlft (l) = 4.0 * vlft(l) * wlft(l) * wlft(l)
80.   1 Vr2           zrgh (l) = 4.0 * vrgh(l) * wrgh(l) * wrgh(l)
81.   1 Vr2           zlft (l) = -zlft(l) * wlft(l)/(zlft(l) - gamfac2*(pmid(l) - plft(l)))
82.   1 Vr2           zrgh (l) = zrgh(l) * wrgh(l)/(zrgh(l) - gamfac2*(pmid(l) - prgh(l)))
83.   1 Vr2           umidl(l) = ulft(l) - (pmid(l) - plft(l)) / wlft(l)
84.   1 Vr2           umidr(l) = urgh(l) + (pmid(l) - prgh(l)) / wrgh(l)
85.   1 Vr2           pmid (l) = pmid(l) + (umidr(l) - umidl(l))*(zlft(l) * zrgh(l)) / (zrgh(l) - zlft(l))
86.   1 Vr2           pmid (l) = max(smallp,pmid(l))
87.   1 Vr2           if (abs(pmld(l)-pmold(l))/pmid(l) < tol ) then
88.   1 Vr2               not_converge(l) = .false.
89.   1 Vr2           endif
90.   1 Vr2           endif
91.   1 Vr2-->  enddo
92.   1           endif
93.   1           enddo
94. 1-----> enddo
```

# Important Loop in Parabola- VH1



```
35.          !      zero out da(n) if a(n) is a local max/min
36. f----< do n = nmin-1, nmax+1
37. f      if(diffa(n-1)*diffa(n) < 0.0) da(n) = 0.0
38. f----> enddo
39. V----< do n = nmin-1, nmax
40. V      ar(n) = a(n) + para(n,1)*diffa(n) + para(n,2)*da(n+1) + para(n,3)*da(n)
41. V      al(n+1) = ar(n)
42. V----> enddo
43. fVr2--< do n = nmin, nmax
44. fVr2      onemfl= 1.0 - flat(n)
45. fVr2      ar(n) = flat(n) * a(n) + onemfl * ar(n)
46. fVr2      al(n) = flat(n) * a(n) + onemfl * al(n)
47. fVr2--> enddo
48. f----< do n = nmin, nmax
49. f      deltaa(n) = ar(n) - al(n)
50. f      a6(n)      = 6. * (a(n) - .5 * (al(n) + ar(n)))
51. f      scrch1(n) = (ar(n) - a(n)) * (a(n)-al(n))
52. f      scrch2(n) = deltaa(n) * deltaa(n)
53. f      scrch3(n) = deltaa(n) * a6(n)
54. f----> enddo
55. f----< do n = nmin, nmax
56. f      if(scrch1(n) <= 0.0) then
57. f          ar(n) = a(n)
58. f          al(n) = a(n)
59. f      endif
60. f      if(scrch2(n) < +scrch3(n)) al(n) = 3. * a(n) - 2. * ar(n)
61. f      if(scrch2(n) < -scrch3(n)) ar(n) = 3. * a(n) - 2. * al(n)
62. f----> enddo
63. f----< do n = nmin, nmax
64. f      deltaa(n)= ar(n) - al(n)
65. f      a6(n) = 6. * (a(n) - .5 * (al(n) + ar(n)))
66. f----> enddo
```

# Important Loop in Parabola- VH1



```
35.  
36.           !      zero out da(n) if a(n) is a local max/min  
37.   f----< do n = nmin-1, nmax+1  
38.   f          if(diffa(n-1)*diffa(n) < 0.0) da(n) = 0.0  
39.   f----> enddo  
40.   V----< do n = nmin-1, nmax  
41.   V          ar(n) = a(n) + para(n,1)*diffa(n) + para(n,2)*da(n+1) + para(n,3)*da(n)  
42.   V          al(n+1) = ar(n)  
43.   V----> enddo  
44.   Vr2---< do n = nmin, nmax  
45.   Vr2          onemfl= 1.0 - flat(n)  
46.   Vr2          ar(n) = flat(n) * a(n) + onemfl * ar(n)  
47.   Vr2          al(n) = flat(n) * a(n) + onemfl * al(n)  
48.   Vr2          deltaan = ar(n) - al(n)  
49.   Vr2          a6n     = 6. * (a(n) - .5 * (al(n) + ar(n)))  
50.   Vr2          scrch1n = (ar(n) - a(n)) * (a(n)-al(n))  
51.   Vr2          scrch2n = deltaan * deltaan  
52.   Vr2          scrch3n = deltaan * a6n  
53.   Vr2          if(scrch1n <= 0.0) then  
54.   Vr2              ar(n) = a(n)  
55.   Vr2              al(n) = a(n)  
56.   Vr2          endif  
57.   Vr2          if(scrch2n < +scrch3n) al(n) = 3. * a(n) - 2. * ar(n)  
58.   Vr2          if(scrch2n < -scrch3n) ar(n) = 3. * a(n) - 2. * al(n)  
59.   Vr2          deltaa(n)= ar(n) - al(n)  
60.   Vr2          a6(n) = 6. * (a(n) - .5 * (al(n) + ar(n)))  
61.   Vr2---> enddo
```

---

C O M P U T E | S T O R E | A N A L Y Z E

# VH1 Profile after optimization



64.5%	233.814721	--	--	9,830,702.0	USER
<hr/>					
11.3%	40.996789	3.114000	7.1%	3,686,400.0	parabola_
8.7%	31.436593	1.140046	3.5%	100.0	sweeypy_
8.1%	29.199941	23.521407	44.6%	409,600.0	riemann_
7.7%	27.759222	1.217341	4.2%	409,600.0	remap_
6.3%	22.928404	0.550848	2.3%	50.0	sweepz_
4.2%	15.280661	1.449079	8.7%	50.0	sweepx1_
4.1%	14.690055	0.674970	4.4%	819,200.0	paraset_
3.3%	11.803449	0.951406	7.5%	409,600.0	evolve_
2.5%	8.876112	0.571191	6.0%	409,600.0	ppmlr_
2.4%	8.679586	0.386993	4.3%	409,600.0	states_
1.8%	6.487615	0.360199	5.3%	1,228,800.0	volume_
1.4%	5.148014	0.201405	3.8%	409,600.0	flatten_
1.3%	4.748816	0.153916	3.1%	50.0	sweepx2_
1.0%	3.595215	0.336065	8.6%	1,228,800.0	forces_
<hr/>					

# Where are we



- **Understand that Memory Bandwidth Utilization is first and foremost**
- **Then we vectorize major computational kernels**
- **Next we need to investigate how best to utilize available cores**
  - First we will look at the best OpenMP examples
  - Then we will look at smaller OpenMP examples
  - Conclude what is important in the OpenMP implementation
  - Lastly look at examples of identifying how many MPI tasks we should run on the node versus how many shared memory threads we should use



# SPMD OpenMP

---

C O M P U T E

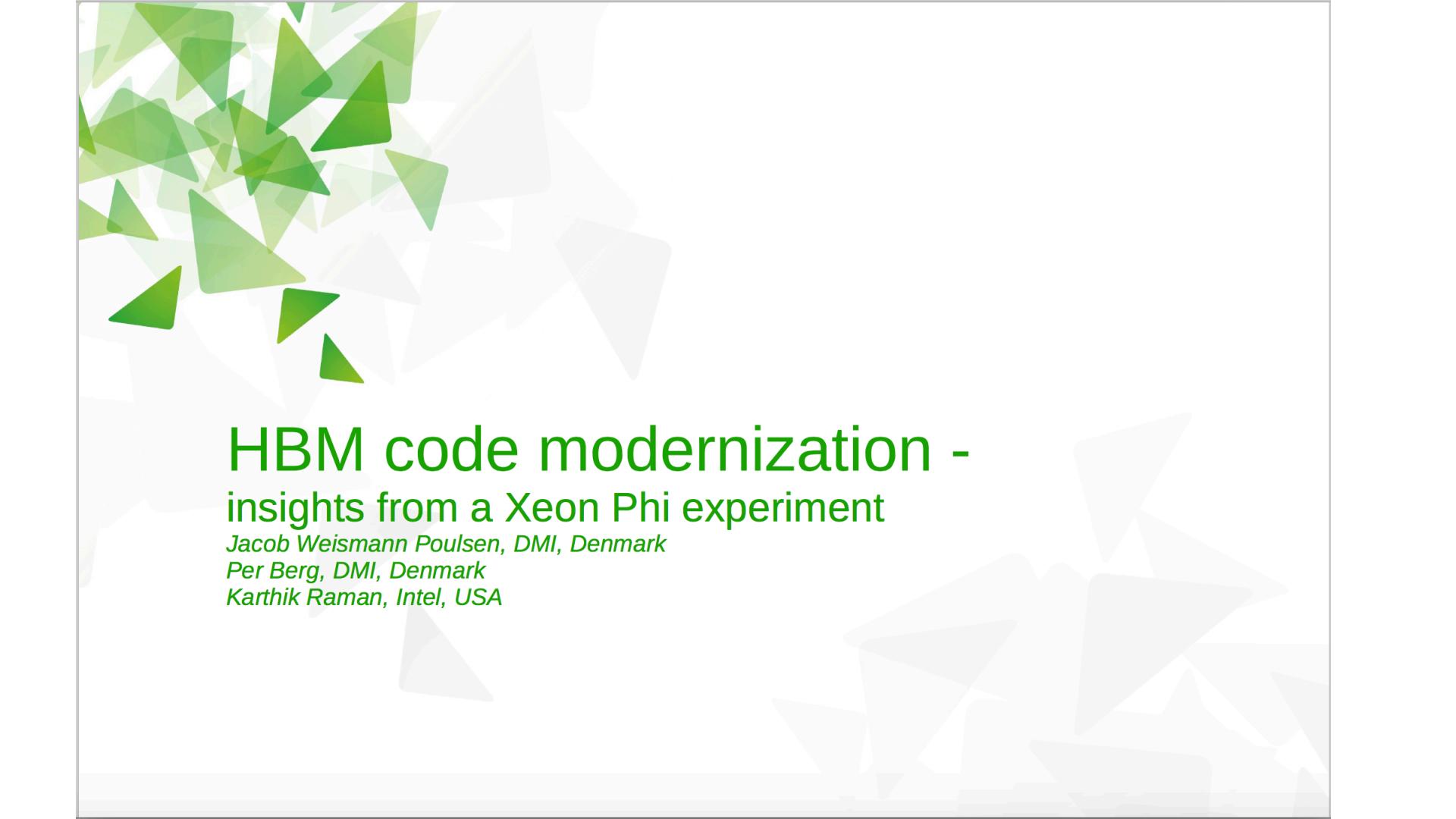
| S T O R E

| A N A L Y Z E

# Can we take clues from all-MPI advantages and disadvantages?



- **Can we force locality like MPI does?**
  - This requires that the thread that uses the data, allocates the data within its NUMA region
  - Tradition OpenMP has no notion of locality
- **Can we allow threads to work asynchronously?**
  - Tradition OpenMP implies barriers after a parallel region
  - Loop level parallelism forces too much synchronization
  - With SPMD OpenMP, user has complete control over synchronization
- **Can we somehow control scheduling of the threads to enable more dynamic re-distribution of work**
  - With SPMD OpenMP, user can write sophistication thread scheduling routine



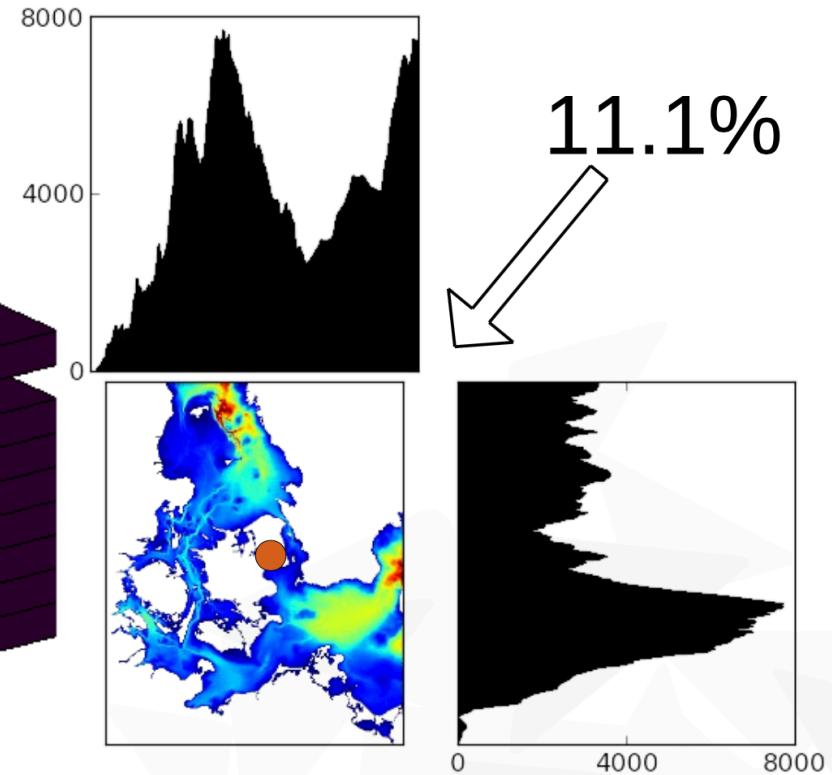
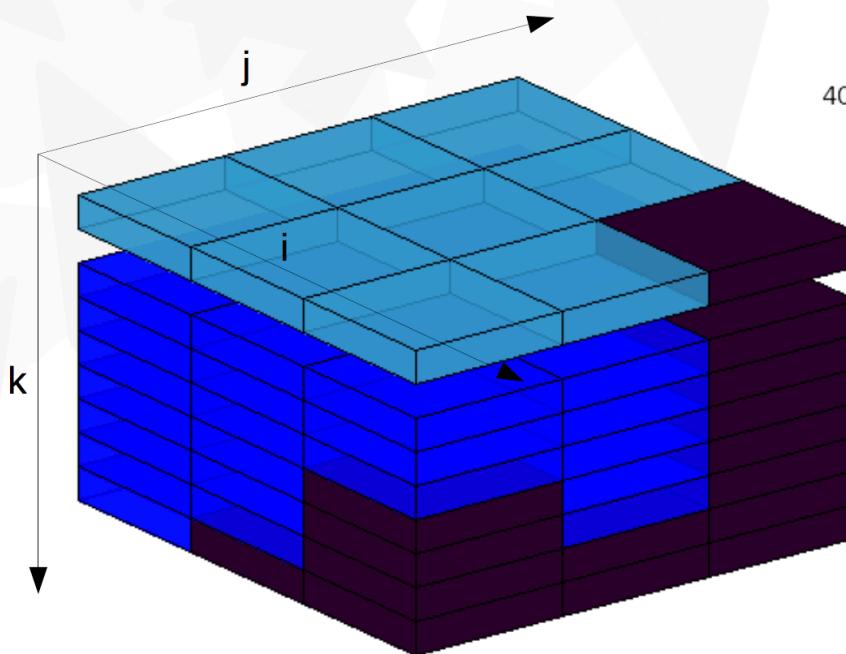
# HBM code modernization - insights from a Xeon Phi experiment

*Jacob Weismann Poulsen, DMI, Denmark*

*Per Berg, DMI, Denmark*

*Karthik Raman, Intel, USA*

# The data is sparse and highly irregular



# Data layout for threads (or tasks + explicit halo)

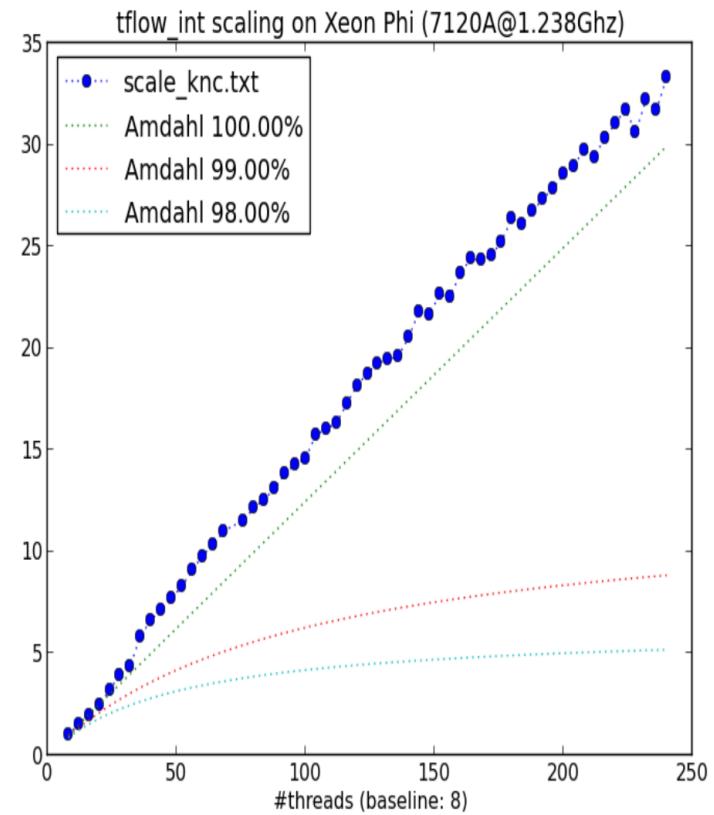
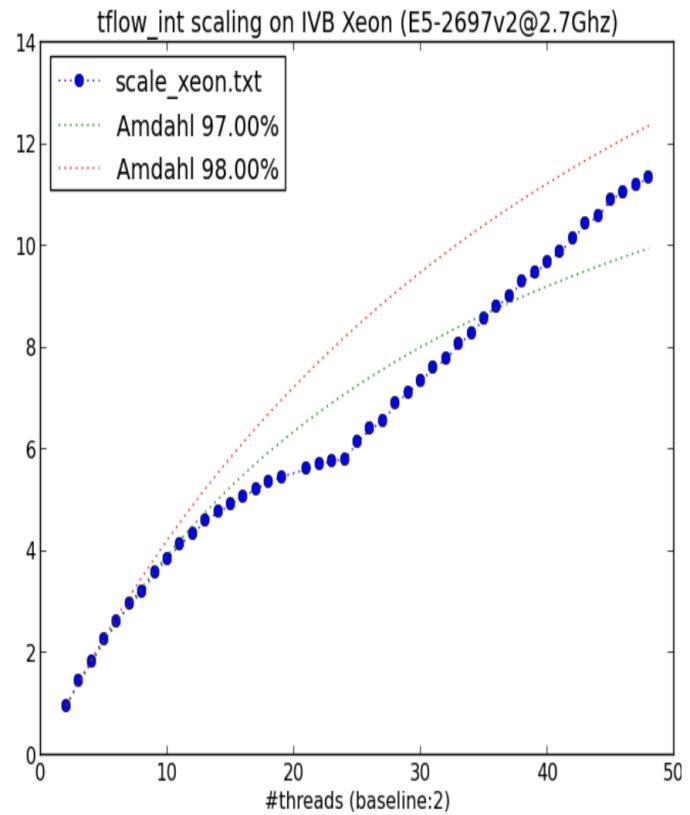
- Each thread will handle a subinterval of columns:



- Another layout of the columns will impose another decomposition for the threads (and the tasks).

```
...
!$OMP PARALLEL DEFAULT(SHARED)
call foo( ... );call bar(...); ...
!$OMP BARRIER
call halo_update(...)
!$OMP BARRIER
call baz( ... );call quux(...); ...
!$OMP END PARALLEL
...
subroutine foo(...)
  ...
  call domp_get_domain(kh, 1, iw2, nl, nu, idx)
  do iw=nl,nu
    i = ind(1,iw)
    j = ind(2,iw)
    ! all threadlocal wet-points (:,:,:,:) are reached here
  ...
enddo
!
```







# High-level OpenMP and Thread Scalable MPI-RMA:

**Application Study with the Wombat Astrophysical MHD Code**

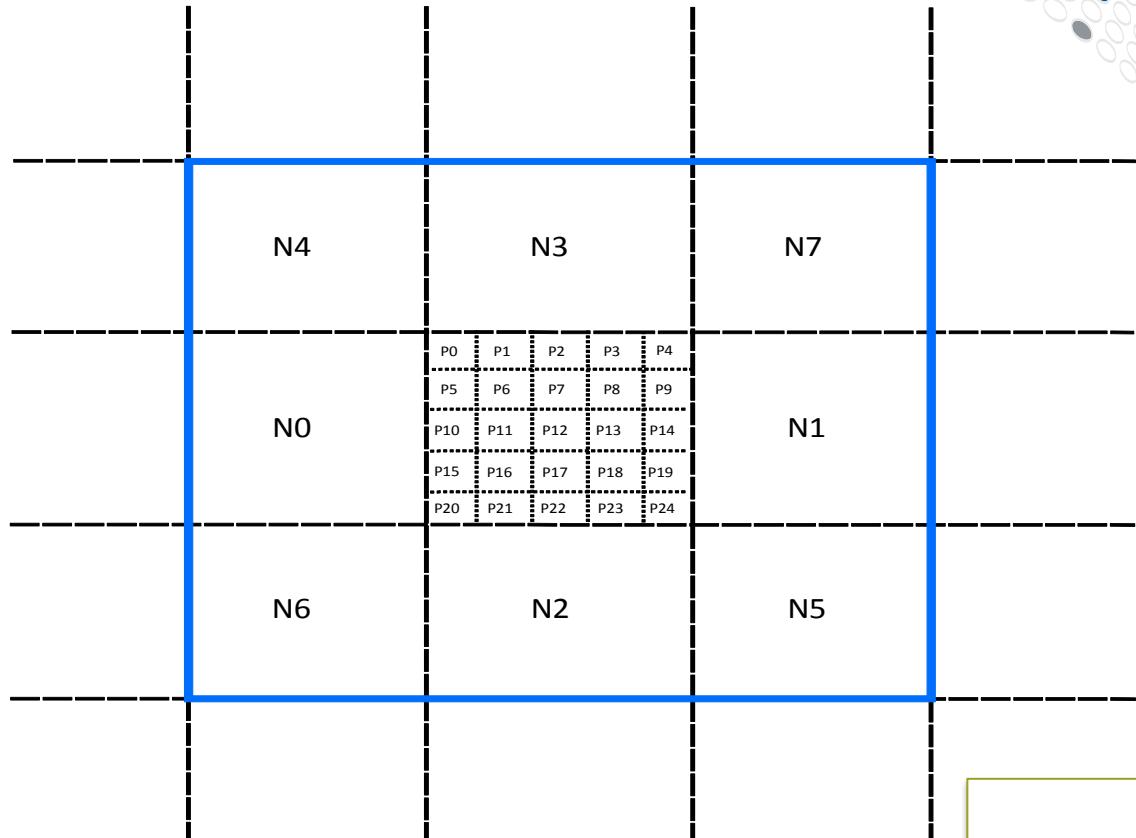
Dr. Peter Mendygral  
Cray Inc.

# Primary Development Concerns

- Two main issues drove the design of Wombat and explain why other codes have not been sufficient for the science
  - Scaling to extreme core count required for resolution requirements
  - Load balancing for SMR/AMR and dark matter particles
- The approach to these problems in Wombat was
  - Make communication matter as little as possible
  - Wide OpenMP on a node to soften impacts of load imbalances (and hardware imbalances) as much as possible before communicating work between ranks
  - Data structures that reduce AMR/SMR complexity and avoid significant global communication for refined patch tracking
  - Do it all in Fortran as that's what works best for me
    - Wombat uses object oriented features from Fortran 2003 and 2008

# Domain Decomposition

- Domain decomposition is represented in the data classes and structures
- “Domain” is an array of “Patches” managed by a MPI rank
- “Patch” is a self-contained, self-describing piece of the world grid at some fixed logical location
- Tunable neighborhood (blue box) sets the horizon of ranks that can be communicated with



# Domains and Patch

- Domains manage bookkeeping for an array of Patches
  - Track a Patch's neighbors
  - Manage allocating/deallocating a Patch's internal grid arrays as needed
  - Multiple Domains are used on a rank for accepting Patches from neighbors
- Patches are
  - Of some uniform fixed size (for a given Domain refinement level)
  - Fixed in a location that is known for all times by all ranks
- Patches provide
  - An atomic unit of work
  - Units to thread across
  - Unit to transfer for load balancing
  - A level of cache blocking

# High-level OpenMP

- Two choices (could be used at the same time) for threading the Domain/Patch design

Option A

! Move OpenMP near the top of the call stack

```
!#OMP PARALLEL  
DO WHILE (t .LT. tend)
```

```
!#OMP DO SCHEDULE(GUIDED)  
DO patch = 1, npatches
```

```
CALL update_patch()
```

! All threads drive MPI

```
END DO
```

```
END DO
```

Option B

! Keep OpenMP within a “compute” loop

```
DO WHILE (t .LT. tend)
```

```
DO patch = 1, npatches
```

```
CALL update_patch()
```

! MPI driven by single thread

```
END DO
```

```
END DO
```

```
SUBROUTINE update_patch()
```

```
!$OMP PARALLEL DO SCHEDULE(STATIC)  
DO i = 1, nx  
...do work...  
END DO
```

```
END SUBROUTINE
```



# High-level OpenMP

- Benefits of OpenMP near the top of the stack
  - Application more closely mimics completely independent processes
    - Less likely to be in the same portion of code at the same time
    - Bandwidth competition may decrease
    - Amdahl's law
  - Threads are less coupled => infrequent thread synchronization
  - Much less likely to have issues with memory conflicts between threads
  - Simpler to implement when done right
    - Large reduction in the amount of OpenMP directives
    - Very little variable scoping needed as most everything is shared => reduced memory footprint

# Wombat Driver and Parallel Region



Setup and object constructors

Thread parallel region

Array allocation and initialization

Time step loop

```
DomainSolver%solve(mhd)  
DomainSolver%solve(ct)  
DomainSolver%solve(passive)  
DomainSolver%solve(particle kick+drift)  
FMGSolver%solve(all levels)  
DomainSolver%solve(particle kick)
```

I/O data dump(s)

Update time step

Array cleanup

Object destructors

Simulation complete

COMPUTE

STORE

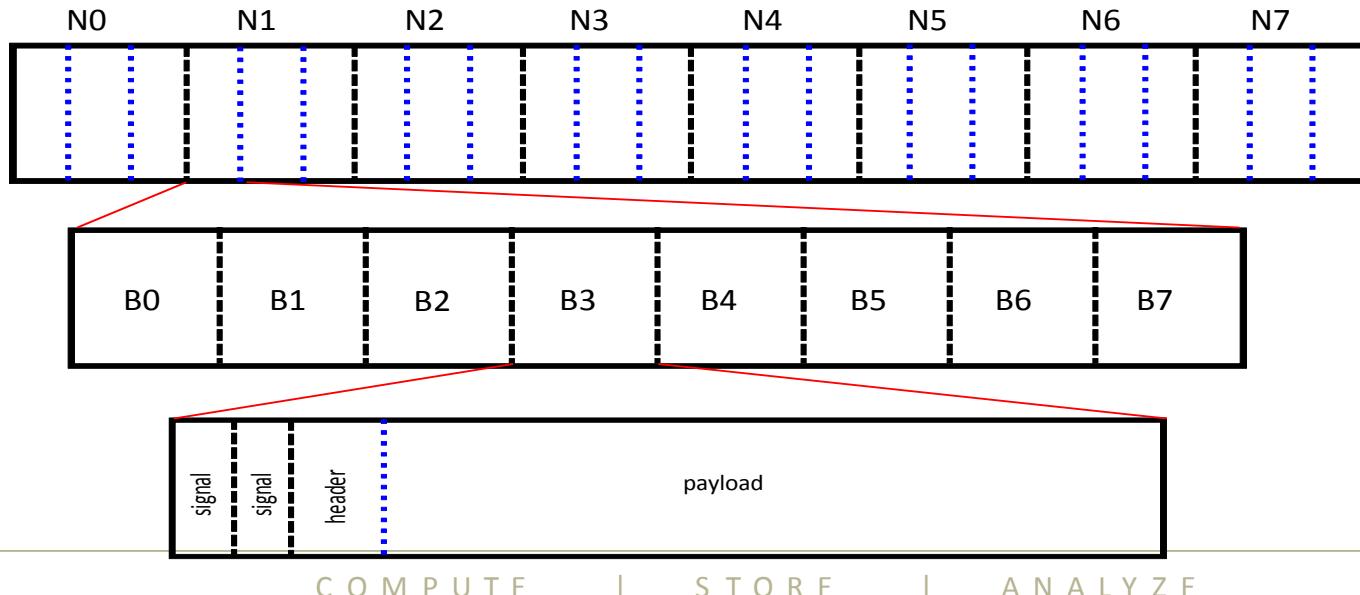
ANALYZE

# Communication Concerns

- If a rank is made much wider with threads, serialization around MPI will limit thread scaling and overall performance
  - Nearly all MPI libraries implement thread safety with a global lock
  - Cray is addressing this issue
    - Released per-object lock library
    - Threading enhancements under design now for two-sided (released per-object lock library a first step)
- Wide OpenMP also means more communication to process
  - Every Patch now has its own smaller boundaries to communicate
  - Starts tipping the behavior towards the message rate limit
  - Two-sided tag matching cannot be done in parallel and will limit thread scaling
    - May start hitting tag limit
- Slower serial performance of KNL => maybe look for the lightest weight MPI layer available
  - MPI-RMA over DMAPP on Cray systems is a thin software layer that achieves similar performance to SHMEM

# Single RMA Window Buffer

- Single buffer used for (almost) all communication
- Messages can be processed concurrently if MPI allows it
- Design is similar to mailboxes within MPI
  - Can process an arbitrary amount of communication



# Generalized Update+Comm Engine



- DomainSolver class
- Single class method capable of driving any of the solvers' data exchange and update process
- Used by every solver in Wombat

```
while # completed Patches < total # Patches
    Decomposition%reset_signals()

    Domain%mark_all_patch_bounds_unresolved()

    while # progressed Patches < total # Patches
        Domain%pack_some_patch_bounds() [all local bounds too if not done]
        Domain%unpack_all_local_patch_bounds() [if not already done]

        ⚡ poll: Decomposition%issue_gets() + update a Patch
        update any ready Patches

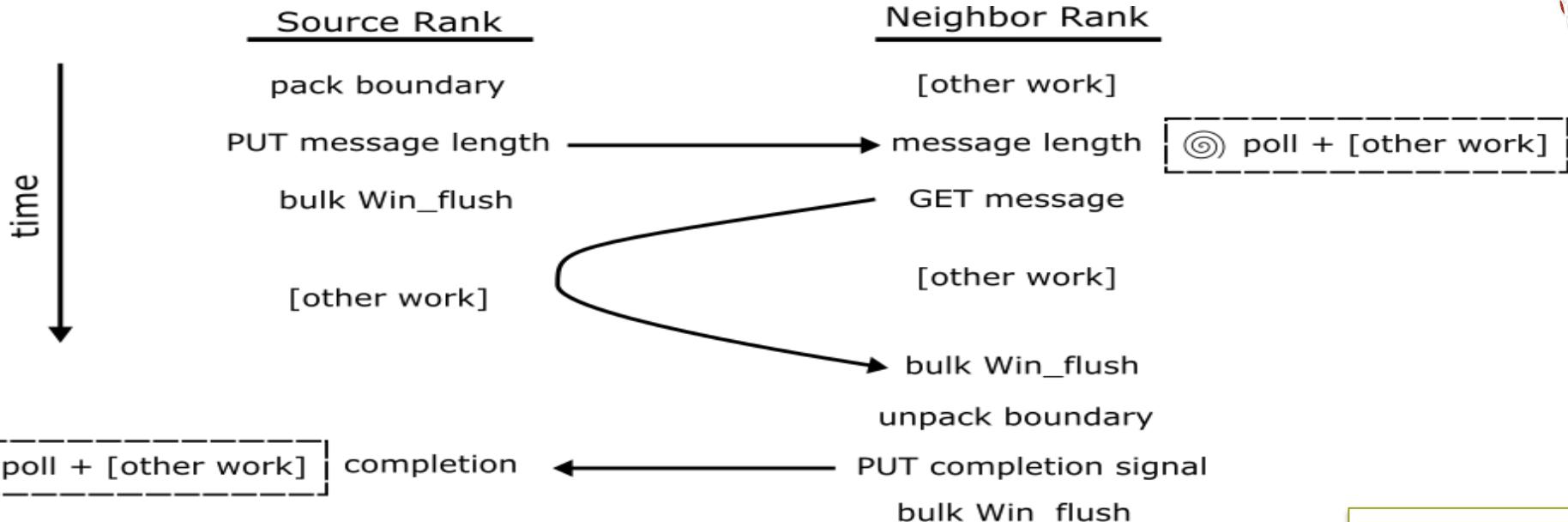
        Decomposition%unpack_mailboxes() [messages flushed]

        ⚡ poll: Decomposition%check_complete_signals + update a Patch

    update Patch progress counters
```

# RMA Boundary Communication Cycle

- Single passive RMA exposure epoch used for the duration of the application
  - No explicit synchronization between ranks
  - RMA semantics make computation/communication overlap simpler to achieve



# Thread Hot MPI-RMA

- **DMAPP library was enhanced to be “thread hot” for SHMEM**
  - “thread hot” is more than “thread safe”
  - “thread hot” implies concurrency and performance across threads was central to the design
- **MPI-RMA over DMAPP leverages this feature as of MPT 7.3.2**
  - No locks used in DMAPP layer
  - Very light weight locking in MPI layer
  - Design makes it very likely that locks are uncontended
  - Network resources efficiently managed among threads
  - Performance approaches that of N independent processes when using N threads
- **Example on HSW with 16 threads each on 2 nodes**
  - OSU passive MPI\_Put bandwidth for 8 B message
  - MPT 7.3.1 = 5.27 MB/s
  - MPT 7.3.2 = 399.9 MB/s
  - 75X improvement

# Wombat Performance

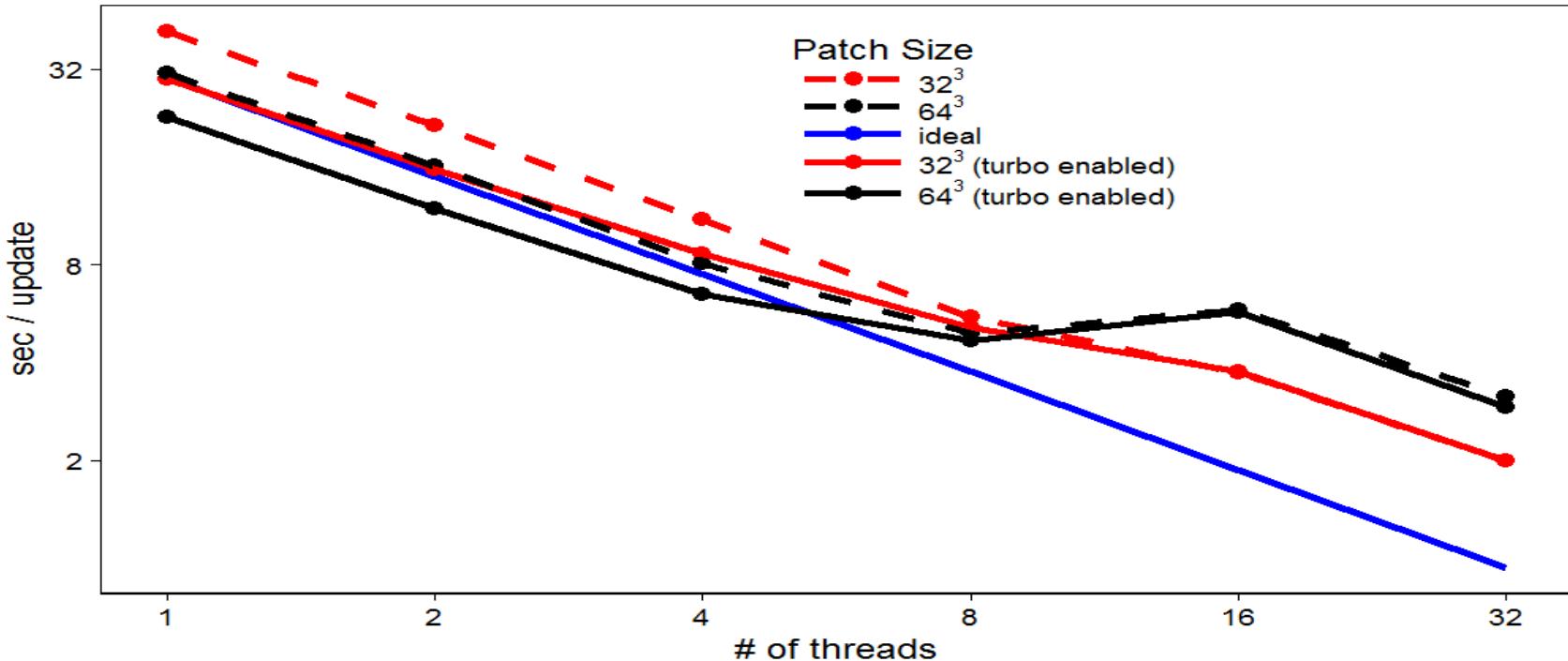


# Benchmarking Notes

- Only MHD was enabled on fixed grid running an MHD shock tube test
- In all cases 2 MB hugepages were used
  - module load craype-hugepages2MB
  - Loaded at link and run time
- The following environment variables were set
  - export HUGETLB\_NO\_RESERVE=yes
  - export MPICH\_MAX\_THREAD\_SAFETY=multiple
  - export MPICH\_RMA\_OVER\_DMAPP=1
- CCE was always used
  - Intel 15 fails to vectorize key routines (see later slide on vector length sensitivity)
  - Intel 16 hits double free corruption in OpenMP run-time. Investigation on-going.
- KNL was configured with MCDRAM as cache

# Haswell Thread Strong Scaling

32 Core - 2.3 GHz - 8,388,608 Zones



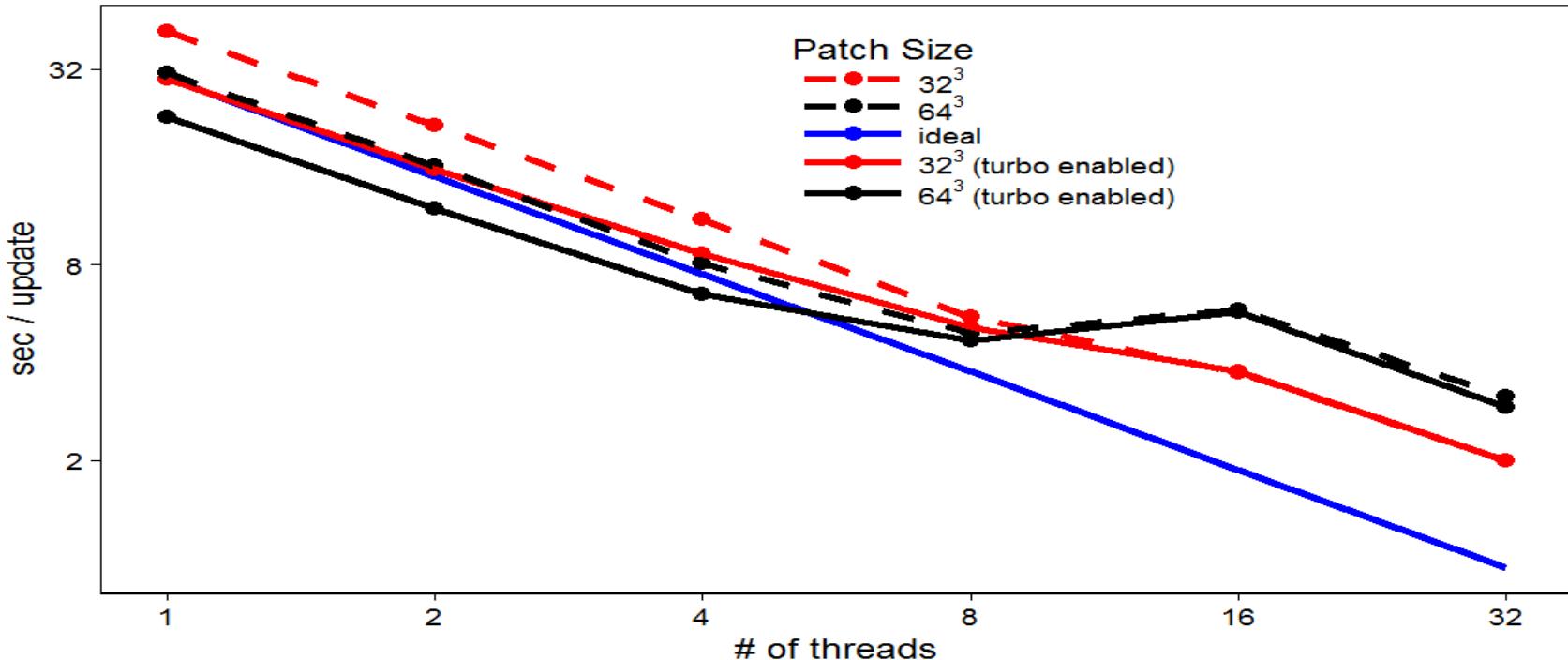
C O M P U T E

| S T O R E

| A N A L Y Z E

## Haswell Thread Strong Scaling

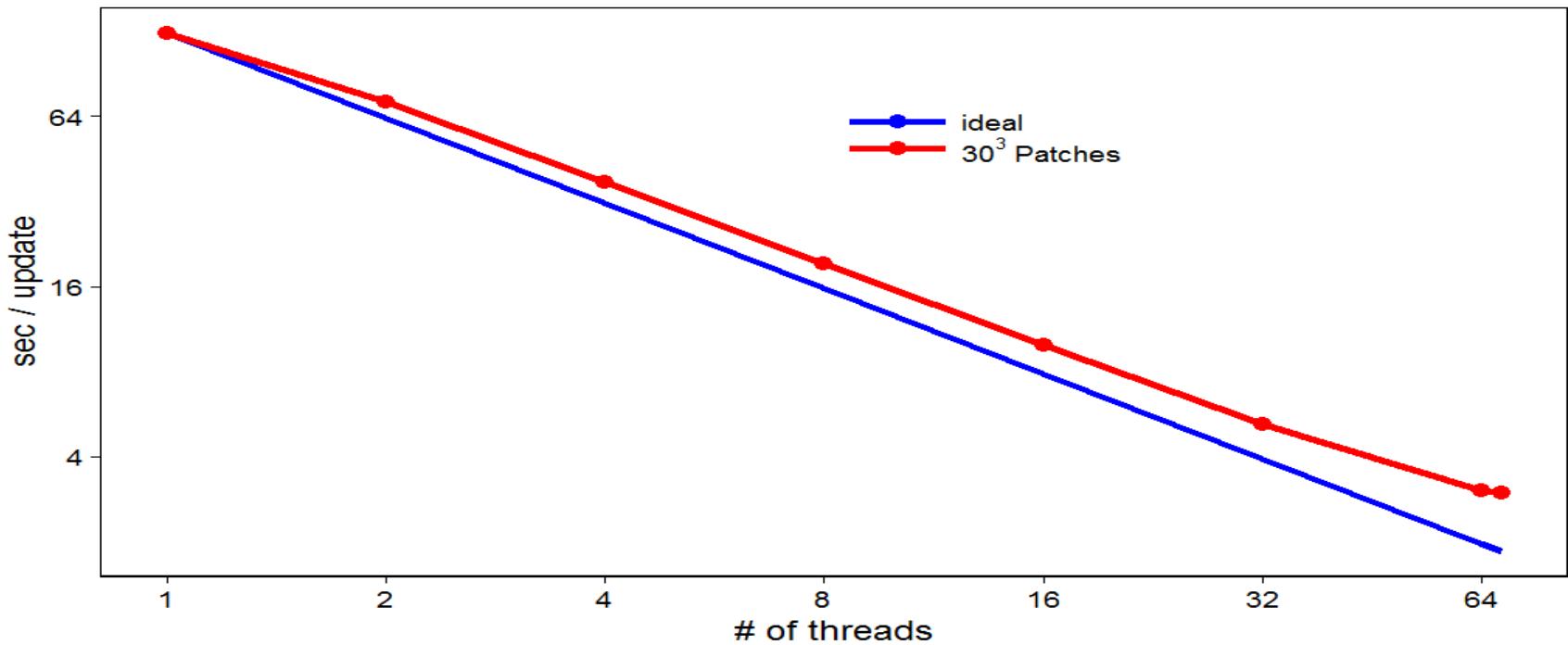
32 Core - 2.3 GHz - 8,388,608 Zones



- Tunable Patch size very important to performance

## KNL Thread Strong Scaling

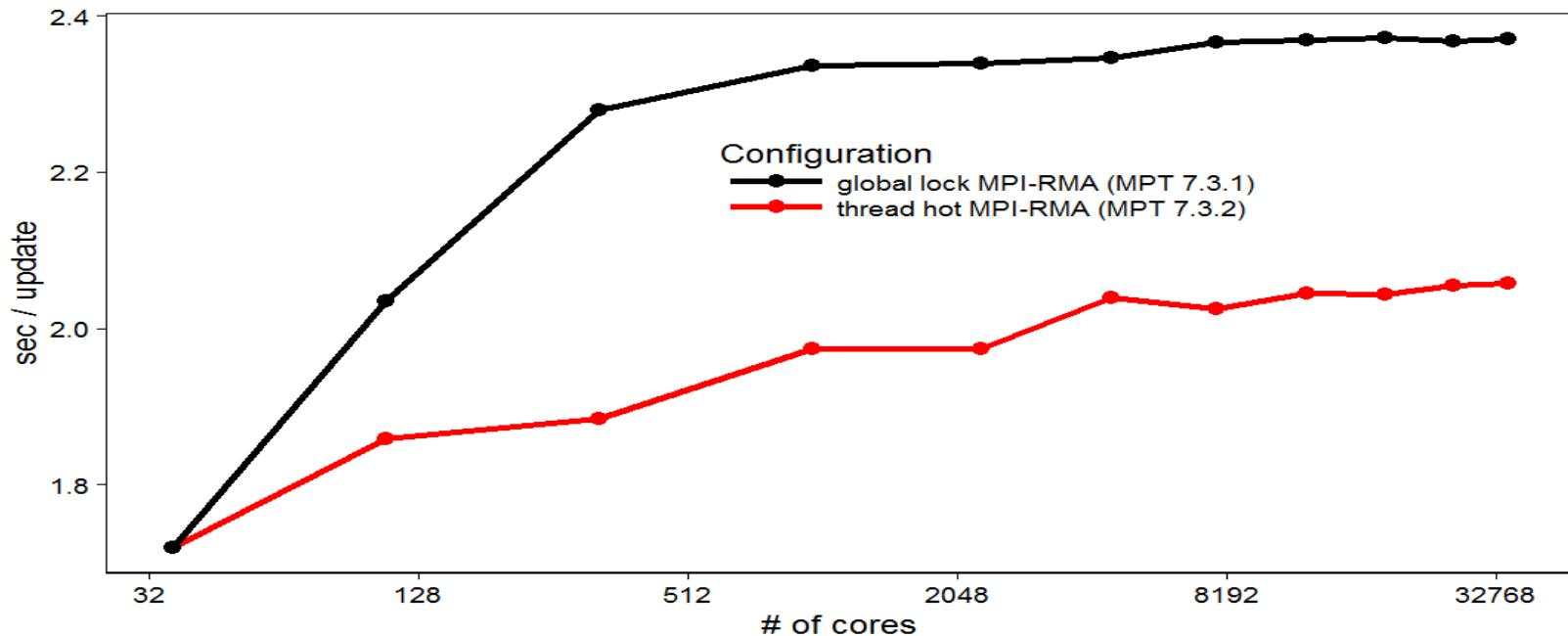
68 Core - 1.4 GHz - 14,688,000 Zones





## XC40 BDW Weak Scaling

1 rank per node - 36 threads per rank - 7,776,000 zones per rank



- Rank reordering

- Cartesian domain optimization for XC topology/placement improves largest run wall time by additional 2.3%

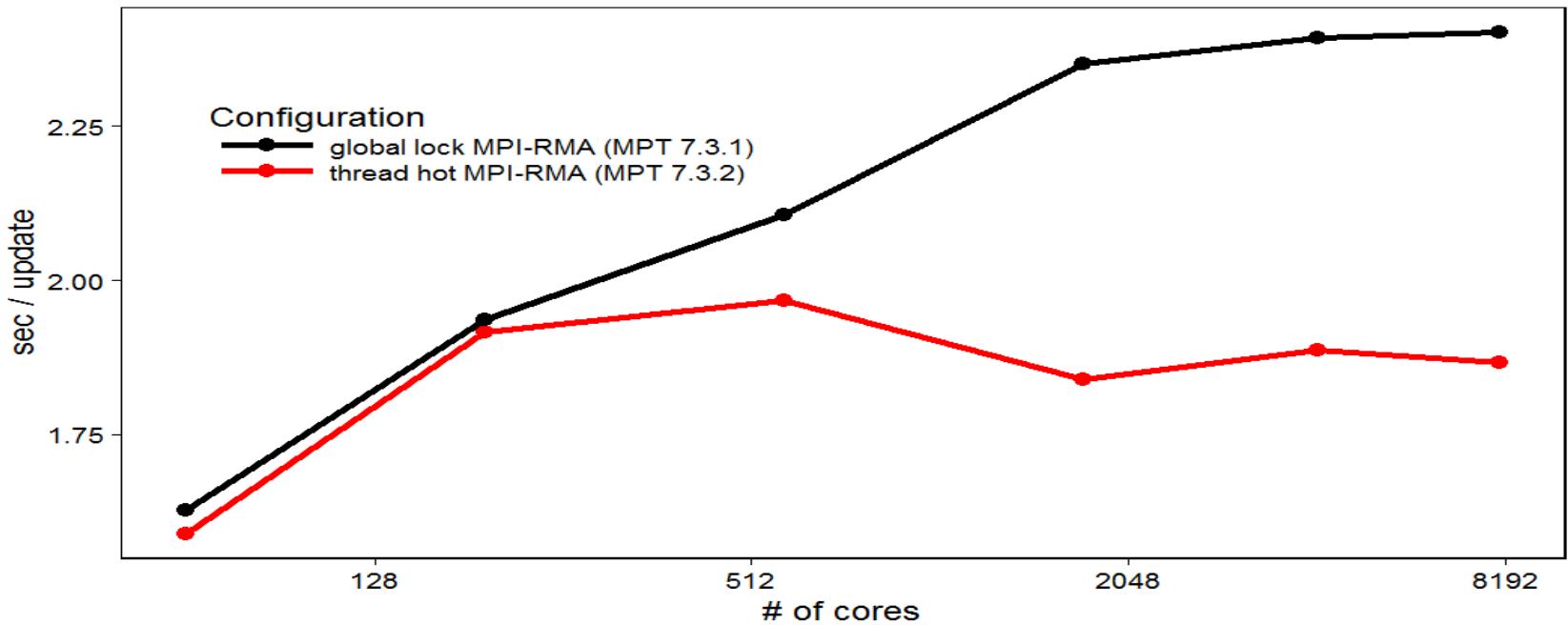
C O M P U T E

| S T O R E

| A N A L Y Z E

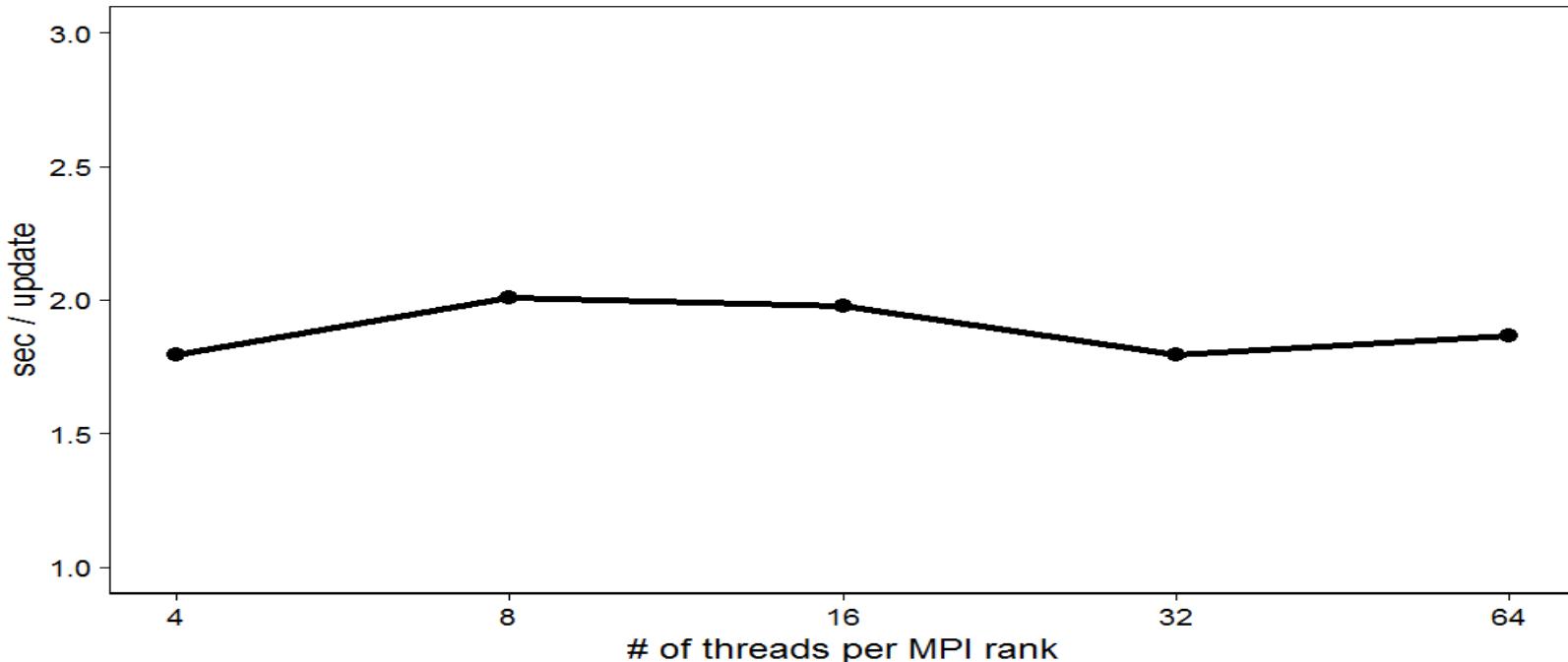
## XC40 KNL Weak Scaling

1 rank per node - 64 threads per rank - 7,776,000 zones per rank



## XC40 KNL Threads/Ranks Comparison

125 Nodes - 8,000 Cores Total - 7,776,000 Zones per Node



- Less than 5% difference between 4 and 64 threads per rank
- Ideal for application like Wombat is 0%

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# Single KNL Node Profile

- 68 threads
- 17x8x4 Patch domain  
(30x30x30 zones per Patches)
- HiMem = 16,076 MB
- 2.97 s/update

Table 1: Profile by Function

Samp%	Samp	Imb.	Imb.	Group
		Samp	Samp%	Function
				Thread=HIDE
100.0%	3,487.0	--	--	Total
75.7%	2,639.0	--	--	USER
22.8%	795.0	61.7	7.3%	compute_fluxes1d\$mod_mhdtdv_
16.4%	572.0	50.2	7.5%	compute_eigenvecs1d\$mod_mhdtdv_
7.0%	244.0	37.7	13.2%	compute_zplane_xyfluxes3d\$mod_mhdtdv_
6.6%	231.0	22.9	9.0%	compute_zfluxes3d\$mod_mhdtdv_
4.0%	140.0	20.4	13.3%	compute_speeds_eigenvals1d\$mod_mhdtdv_
3.0%	103.0	15.4	13.2%	compute_corner_emf3d\$mod_mhdtdv_
2.9%	101.0	23.0	18.2%	compute_zone_averages1d\$mod_mhdtdv_
1.8%	63.0	3.9	5.7%	pack_array3d\$mod_patch_
1.6%	56.0	8.4	12.8%	halfxupdate_states3d\$mod_mhdtdv_
1.5%	53.0	12.2	18.8%	halfxupdate_states3d\$mod_mhdtdv_
1.3%	45.0	8.6	13.5%	halfzupdate_states3d\$mod_mhdtdv_
1.0%	35.0	13.7	25.2%	apply_protections3d\$mod_mhdtdv_
19.3%	672.0	--	--	ETC
7.1%	248.0	1.0	0.8%	_ZL21fullscan_barrier_listii.constprop.2
6.0%	210.0	27.0	11.6%	_fmemalign
1.3%	46.0	33.0	58.9%	_ZL23internal_simple_barrierii.constprop.3
1.2%	42.0	15.9	28.3%	_cray_dset_SN
1.1%	39.0	--	--	_dl_update_slotinfo
5.0%	174.0	--	--	OMP
4.3%	151.0	194.2	49.7%	omp_set_lock
=====				
Total				
=====				
UNHALTED_CORE_CYCLES				
43,959,396,781				
UNHALTED_REFERENCE_CYCLES				
43,669,216,374				
INSTRUCTION_RETIRIED				
30,208,753,315				
LLC_REFERENCES				
2,616,929,294				
LLC_MISSES				
247,223,361				
L3 cache hit,miss ratio				
90.6% hits      9.4% misses				



# Single KNL Node Profile (8X bigger grid)

- 68 threads
- 17x16x16 Patch domain  
(30x30x30 zones per Patches)
- HiMem = 71,550 MB
- 23.93 s/update
  - 8.08X longer time steps
- Spilling outside of MCDRAM cache not a performance issue

Table 1: Profile by Function

Samp%	Samp	Imb.	Imb.	Group	Function	Thread:HIDE
100.0%	24,506.0	--	--	Total		
87.2%	21,381.0	--	--	USER		
25.6%	6,275.0	209.9	3.1%	compute_fluxes1d\$mod_mhdtdv_		
18.0%	4,406.0	175.4	3.5%	compute_eigenvecs1d\$mod_mhdtdv_		
7.3%	1,801.0	133.5	6.4%	compute_zplane_xyfluxes3d\$mod_mhdtdv_		
6.9%	1,683.0	106.8	5.6%	compute_zfluxes3d\$mod_mhdtdv_		
4.5%	1,099.0	69.2	5.9%	compute_speeds_eigenvals1d\$mod_mhdtdv_		
3.5%	868.0	8.7	1.0%	pack_array3d\$mod_patch_		
3.4%	821.0	68.1	7.3%	compute_zone_averages1d\$mod_mhdtdv_		
3.2%	794.0	45.2	5.0%	compute_corner_emf3d\$mod_mhdtdv_		
2.8%	694.0	11.0	1.5%	unpack_array3d\$mod_patch_		
1.7%	411.0	35.3	7.4%	halfupdate_states3d\$mod_mhdtdv_		
1.6%	394.0	29.8	6.6%	halfupdate_states3d\$mod_mhdtdv_		
1.6%	392.0	44.2	9.5%	halfupdate_states3d\$mod_mhdtdv_		
1.4%	333.0	59.6	12.6%	apply_protections3d\$mod_mhdtdv_		
11.0%	2,703.0	--	--	ETC		
7.4%	1,819.0	133.9	6.9%	fmemalign		
1.2%	304.0	52.0	25.5%	ZL21fullscan_barrier_listiii.constprop.2		
1.1%	266.0	26.9	7.9%	cray_dset_SN		
1.7%	421.0	--	--	OMP		
1.6%	397.0	210.7	50.2%	omp_set_lock		
=====						
Total						
=====						
UNHALTED_CORE_CYCLES						
289,624,816,016						
UNHALTED_REFERENCE_CYCLES						
301,384,184,262						
INSTRUCTION_RETIRED						
167,529,858,350						
LLC_REFERENCES						
16,267,607,578						
LLC_MISSES						
1,842,082,443						
L3 cache hit,miss ratio						
88.7% hits    11.3% misses						
=====						



# 27 KNL Nodes Profile

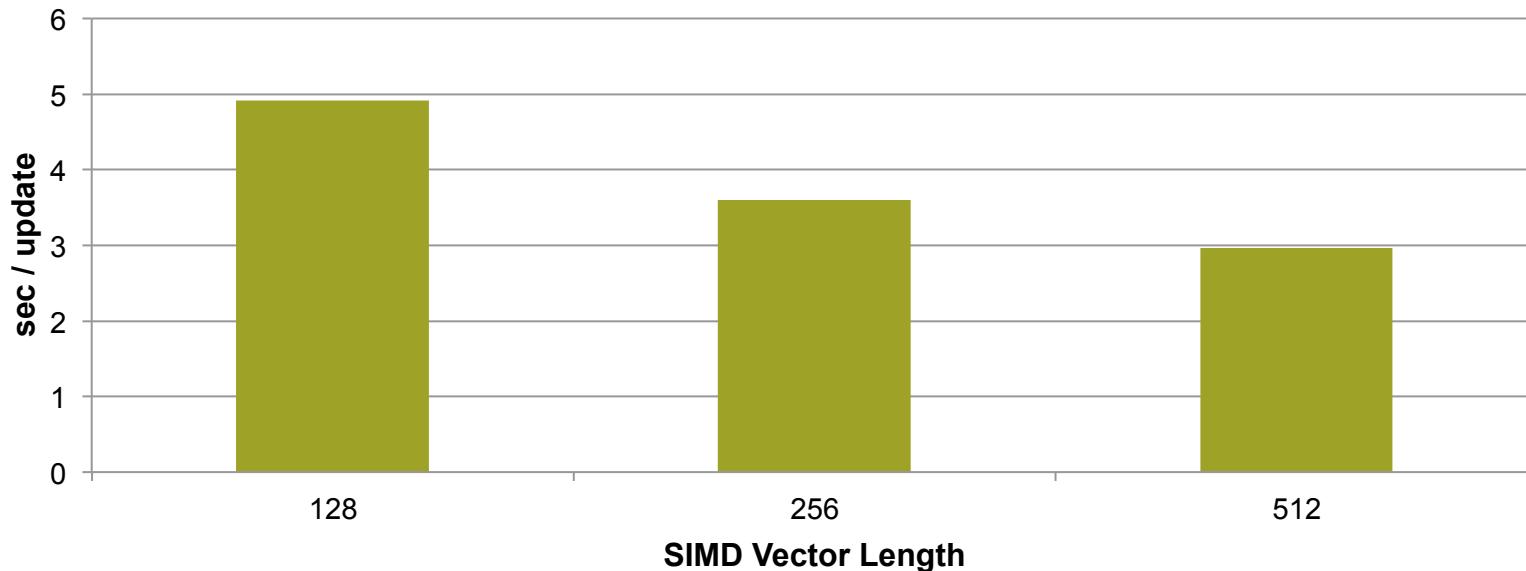
- 68 threads per rank
- 27 ranks
- 17x16x16 Patch domain (30x30x30 zones per Patches)
- HiMem = 32,678 MB
- 3.43 s/update

Table 1: Profile by Function

Samp%	Samp	Imb.	Imb.	Group
	Samp	Samp%	Function	
			PE=HIDE	
100.0%	4,330.7	--	--	Total
60.6%	2,622.6	--	--	USER
16.4%	710.0	49.0	6.7%	compute_fluxes1d\$mod_mhdtdv_
12.0%	519.8	42.2	7.8%	compute_eigenvecs1d\$mod_mhdtdv_
5.2%	225.8	7.2	3.2%	initrma\$mod_decomposition_
5.1%	220.3	32.7	13.4%	compute_zplane_xyfluxes3d\$mod_mhdtdv_
4.7%	205.3	21.7	9.9%	compute_zfluxes3d\$mod_mhdtdv_
2.7%	117.0	19.0	14.5%	compute_speeds_eigenvals1d\$mod_mhdtdv_
2.1%	91.6	26.4	23.3%	compute_zone_averages1d\$mod_mhdtdv_
2.0%	88.5	11.5	11.9%	compute_corner_emf3d\$mod_mhdtdv_
1.2%	53.9	8.1	13.6%	pack_array3d\$mod_patch_
1.1%	49.7	7.3	13.2%	halfyupdate_states3d\$mod_mhdtdv_
1.1%	45.9	7.1	13.9%	halfzupdate_states3d\$mod_mhdtdv_
1.0%	44.9	5.1	10.7%	halfxupdate_states3d\$mod_mhdtdv_
25.5%	1,104.1	--	--	ETC
12.8%	553.3	154.7	22.7%	ZL21fullscan_barrier_listii.constprop.2
4.6%	199.7	43.3	18.5%	fmemalign
2.8%	120.8	32.2	21.8%	dl_update_slotinfo
1.2%	53.4	32.6	39.4%	update_get_addr
9.3%	402.6	--	--	OMP
4.5%	194.9	224.1	55.5%	omp_set_lock
4.4%	189.6	63.4	26.0%	cray\$mt_barrier_part_prime_wait_others
4.6%	201.3	--	--	MPI
3.9%	168.0	34.0	17.5%	mpi_put

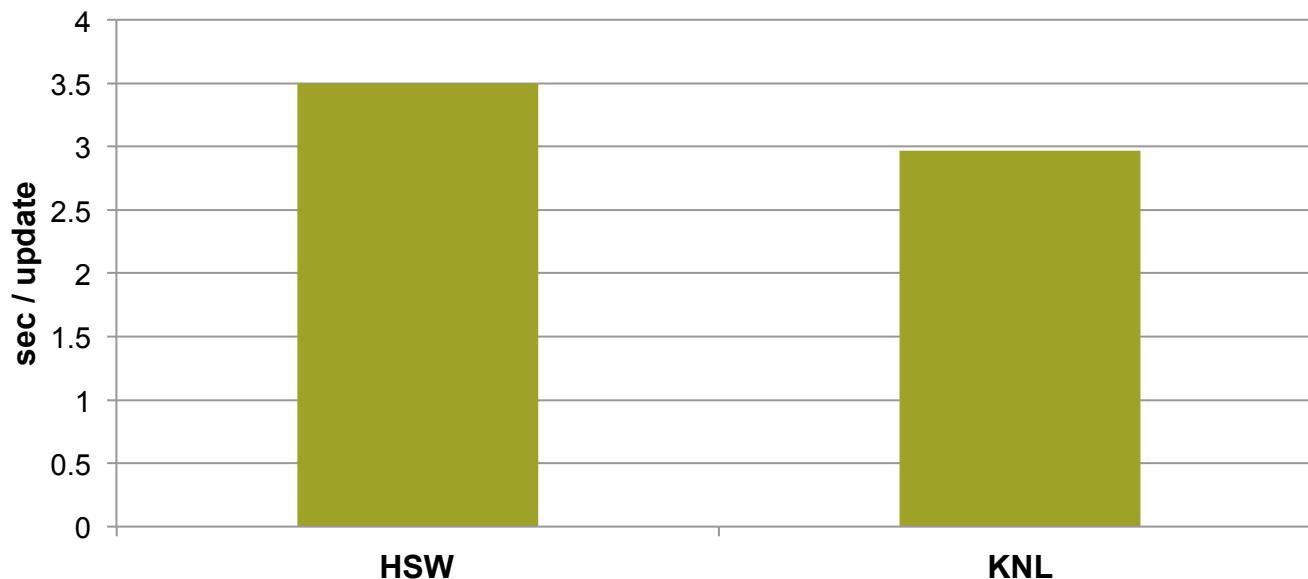
## KNL Performance - SIMD Vector Length

14,688,000 zones - 68 threads



- Vector length is important
- Significant effort went into making solver loops (compute and conv) vectorize

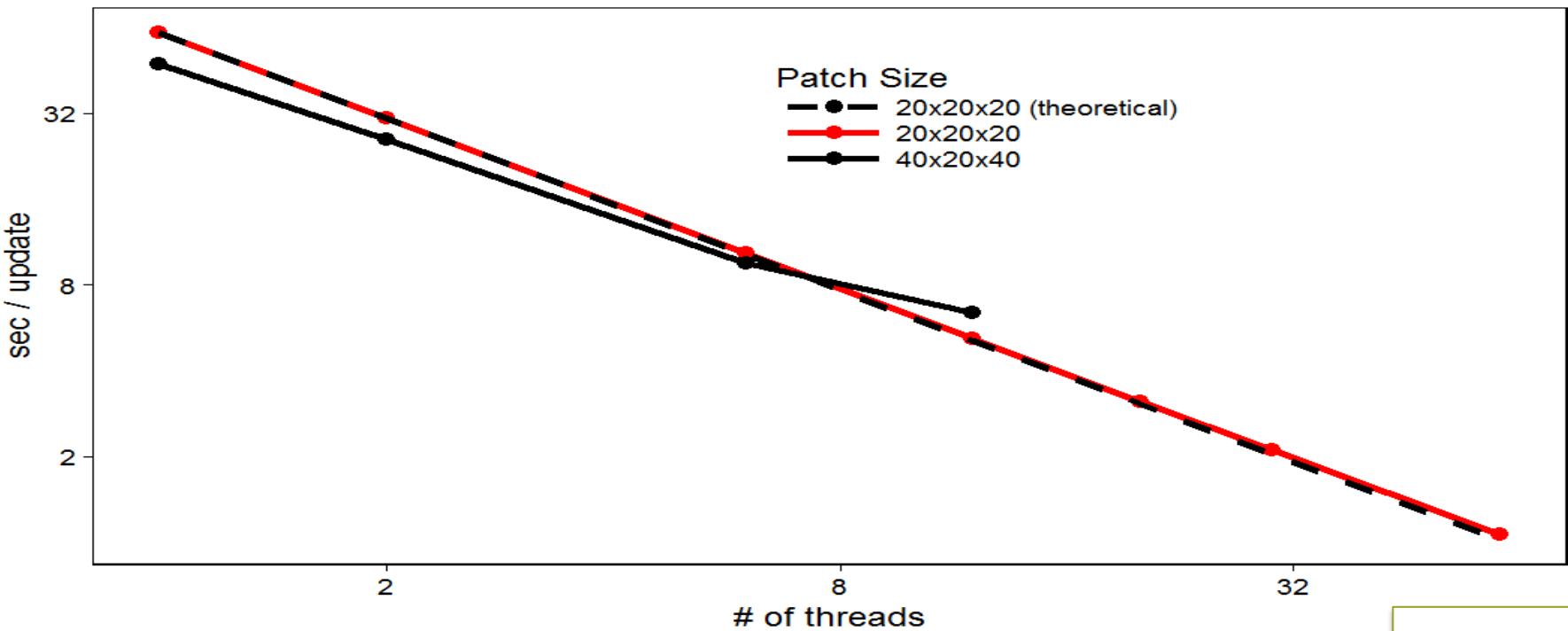
## HSW/KNL Node Performance Comparison 14,688,000 zones



- For this problem size a single KNL node is ~15% faster than a HSW node
- Have measured as high as 20% for other problem sizes
- Multi-node comparisons trend similar (KNL is 5-20% faster depending on exact problem)

## KNC Thread Strong Scaling

60 Core - 1 GHz - 480,000 Zones



# SPMD OpenMP Pros



- Much less overhead than low level !\$OMP PARALLEL DO and !\$OMP DO
- Scoping is handled with language standard
- In very complicated applications, the number of directives and code modifications is much less

# SPMD OpenMP Cons



- Users must really understand threading and the implications of how the language handles global and local storage which would become shared by all threads and private to a thread.



# Analysis of BGW Kernels

A study of higher level OpenMP

# Introduction



- BGW test has three kernels
- Reveal does very well on one
- Optimization improved other two significantly

# Profiling the BGW test

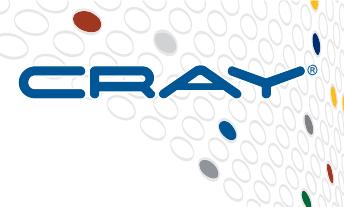


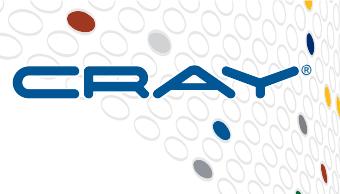
Table 2: Inclusive and Exclusive Time in Loops (from -hprofile\_generate)

Loop Incl	Time	Loop Hit	Loop Trips	Loop Trips	Loop Trips	Function=/.LOOP[.]
Incl Time	(Loop Adj.)		Trips	Trips	Trips	
Time%			Avg	Min	Max	
-----						
94.5%	55.048452	0.003011	1	96.0	96	hackakernel_.LOOP.04.li.121
91.5%	53.298164	0.015635	96	240.0	240	hackakernel_.LOOP.08.li.166
91.5%	53.282529	0.055862	23,040	100.0	100	hackakernel_.LOOP.09.li.170
91.4%	53.226667	53.226667	2,304,000	8,000.0	8,000	hackakernel_.LOOP.10.li.177
2.8%	1.658162	0.713467	96	240.0	240	hackakernel_.LOOP.11.li.192
1.6%	0.944695	0.944695	22,944	6,000.0	6,000	hackakernel_.LOOP.12.li.215
0.1%	0.081882	0.000093	96	100.0	100	hackakernel_.LOOP.06.li.138
0.1%	0.081789	0.081789	9,600	8,000.0	8,000	hackakernel_.LOOP.07.li.145
0.0%	0.005393	0.005393	96	6,000.0	6,000	hackakernel_.LOOP.13.li.243
0.0%	0.000204	0.000204	96	6,000.0	6,000	hackakernel_.LOOP.05.li.129
0.0%	0.000000	0.000000	1	240.0	240	hackakernel_.LOOP.02.li.79
0.0%	0.000000	0.000000	1	240.0	240	hackakernel_.LOOP.03.li.85
0.0%	0.000000	0.000000	1	96.0	96	hackakernel_.LOOP.01.li.60
=====						

Loop B

Loop C

Loop A



# Kernel A loop

```
138. + 1 2-----<          do my_igp = 1, ngpown      (100)
139.   1 2               if (my_igp .gt. ncouls .or. my_igp .le. 0) cycle
140.   1 2
141.   1 2               igmax=ncouls
142.   1 2
143.   1 2               mygpvar1 = CONJG(leftvector(my_igp,n1))
144.   1 2
145.   1 2 Vr2----<          do ig = 1, igmax      (800)
146.   1 2 Vr2           matngmatmgpD(ig,my_igp) = rightvector(ig,n1) * mygpvar1
147.   1 2 Vr2---->
148.   1 2----->          enddo
                           enddo
```

Which of these two loop should we parallelize?  
Are there some inefficiencies?

# Kernel B loop

```

165.      1 A-----<>          schdt_array = 0D0
166. + 1 b-----<          do ifreq=1,nFreq           (240)
167.      1 b
168.      1 b
169.      1 b
170. + 1 b b-----<          schDt = (0D0,0D0)
171.      1 b b
172.      1 b b
173.      1 b b
174.      1 b b
175.      1 b b
176.      1 b b
177.      1 b b Vr3--<          do my_igp = 1, ngpown   (100)
178.      1 b b Vr3
179.      1 b b Vr3
180.      1 b b Vr3
181.      1 b b Vr3
182.      1 b b Vr3-->
183.      1 b b
184.      1 b b----->          igmax=ncouls
185.      1 b
186.      1 b----->          schDtt = (0D0,0D0)
187.      1
188. + 1          do ig = 1, igmax           (8000)
189.      1          I_epsRggp_int = I_epsR_array(ig,my_igp,ifreq)
190.      1          I_epsAggp_int = I_epsA_array(ig,my_igp,ifreq)

```

---

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Once again which loop is best to target for parallelization

# Kernel C loop

```

192. + 1 2-----<          do ifreq=1,nFreq           (240)
193.   1 2
194.   1 2          schDt = schDt_array(ifreq)
195.   1 2
196.   1 2          cedifft_zb = dFreqGrid(ifreq)
197.   1 2          cedifft_coh = CMPLX(cedifft_zb,0D0) - dFreqBrd(ifreq)
198.   1 2
199.   1 2          if (ifreq .ne. 1) then
200.     1 2          cedifft_zb_right = cedifft_zb
201.     1 2          cedifft_zb_left = dFreqGrid(ifreq-1)
202.     1 2          schDt_right = schDt
203.     1 2          schDt_left = schDt_array(ifreq-1)
204.     1 2          schDt_avg = 0.5D0 * ( schDt_right + schDt_left )
205.     1 2          schDt_lin = schDt_right - schDt_left
206.     1 2          schDt_lin2 = schDt_lin/(cedifft_zb_right-cedifft_zb_left)
207.     1 2          endif
208.     1 2
209.   1 2          ! The below two lines are for sigma1 and sigma3
210.     1 2          if (ifreq .ne. nFreq) then
211.       1 2          schDi(:) = schDi(:) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(:)-cedifft_coh)
212.       1 2          schDi_corb(:) = schDi_corb(:) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(:)-cedifft_coh)
213.       1 2          endif
214.       1 2          if (ifreq .ne. 1) then
215.         1 2          do iw = 1, nfreqeval           (6000)

```

Once again which loop is best to target for parallelization

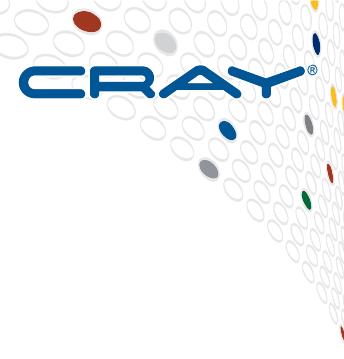
Array Assignment

# Kernel C loop -continued



```
216. 1 2 V      !These lines are for sigma2
217. 1 2 V          intfact=abs((wx(iw)-cedifft_zb_right)/(wx(iw)-cedifft_zb_left))
218. 1 2 V          if (intfact .lt. 1d-4) intfact = 1d-4
219. 1 2 V          if (intfact .gt. 1d4) intfact = 1d4
220. 1 2 V          intfact = -log(intfact)
221. 1 2 V          sch2Di(iw) = sch2Di(iw) - CMPLX(0.d0,prefactor) * schDt_avg * intfact
222. 1 2 V      !These lines are for sigma4
223. 1 2 V          if (flag_occ) then
224. 1 2 V              intfact=abs((wx(iw)+cedifft_zb_right)/(wx(iw)+cedifft_zb_left))
225. 1 2 V              if (intfact .lt. 1d-4) intfact = 1d-4
226. 1 2 V              if (intfact .gt. 1d4) intfact = 1d4
227. 1 2 V              intfact = log(intfact)
228. 1 2 V              schDt_lin3 = (schDt_left + schDt_lin2*(-wx(iw)-cedifft_zb_left))*intfact
229. 1 2 V          else
230. 1 2 V              schDt_lin3 = (schDt_left + schDt_lin2*(wx(iw)-cedifft_zb_left))*intfact
231. 1 2 V          endif
232. 1 2 V          schDt_lin3 = schDt_lin3 + schDt_lin
233. 1 2 V          schDi_cor(iw) = schDi_cor(iw) - CMPLX(0.d0,prefactor) * schDt_lin3
234. 1 2 V----->      enddo
235. 1 2          endif
236. 1 2----->      enddo
```

# Kernel A



```
do my_igp = 1, ngpown
    if (my_igp .gt. ncouls .or. my_igp .le. 0) cycle

    igmax=ncouls

    mygpvar1 = CONJG(leftvector(my_igp,n1))

    do ig = 1, igmax
        matngmatmgpD(ig,my_igp) = rightvector(ig,n1) * mygpvar1
    enddo
enddo
```

## Automatically inserted by Reveal

```
138.      1           ! Directive inserted by Cray Reveal. May be incomplete.  
139.      1 M-----< !$OMP parallel do default(none) &  
140.      1 M           !$OMP& private (ig,igmax,mygpvarl,my_igp) &  
141.      1 M           !$OMP& shared (leftvector,matngmatmgpd,n1,ncouls,ngpown,rightvector)  
142. + 1 M m-----< do my_igp = 1, ngpown  
143.     1 M m           if (my_igp .gt. ncouls .or. my_igp .le. 0) cycle  
144.     1 M m  
145.     1 M m           igmax=ncouls  
146.     1 M m  
147.     1 M m           mygpvarl = CONJG(leftvector(my_igp,n1))  
148.     1 M m  
149.     1 M m Vr2----< do ig = 1, igmax  
150.     1 M m Vr2       matngmatmgpD(ig,my_igp) = rightvector(ig,n1) * mygpvarl  
151.     1 M m Vr2---->  
152.     1 M m----->> enddo  
                           enddo
```

# Kernel A



Automatically changed by Levesque

```
137.      1 M-----< !$OMP PARALLEL DO
138.      1 M imV----<           do ig = 1, ncouls
139. + 1 M imV ir4--<           do my_igp = 1, min(ncouls,ngpown)
140.      1 M imV ir4           matngmatmgpD(ig,my_igp) = rightvector(ig,n1) * CONJG(leftvector(my_igp,n1))
141.      1 M imV ir4-->       enddo
142.      1 M imV---->>       enddo
```



```
137. 1 M-----< !$OMP PARALLEL DO
138. 1 M imV-----<      do ig = 1, ncouls
139. + 1 M imV ir4--<      do my_igp = 1, min(ncouls,ngpown)
140. 1 M imV ir4          matngmatmgpD(ig,my_igp) = rightvector(ig,n1) * CONJG(leftvector(my_igp,n1))
141. 1 M imV ir4-->      enddo
142. 1 M imV----->      enddo
```

1:1 Default text Editing disabled  
/Users/levesque/Documents/A-Reveal.rtf

Today, 8:09:55 AM 1,165 bytes RTF ▼ Converted Current Locale (UTF-8) ▼ PC

```
138. 1           ! Directive inserted by Cray Reveal. May be incomplete.
139. 1 M-----< !$OMP parallel do default(None)
140. 1 M           !$OMP& private (ig,igmax,mygpvar1,my_igp)
141. 1 M           !$OMP& shared (leftvector,matngmatmgpD,n1,ncouls,ngpown,rightvector)
142. + 1 M m-----<      do my_igp = 1, ngpown
143. 1 M m           if (my_igp .gt. ncouls .or. my_igp .le. 0) cycle
144. 1 M m
145. 1 M m           igmax=ncouls
146. 1 M m
147. 1 M m           mygpvar1 = CONJG(leftvector(my_igp,n1))
148. 1 M m
149. 1 M m Vr2----<      do ig = 1, igmax
150. 1 M m Vr2          matngmatmgpD(ig,my_igp) = rightvector(ig,n1) * mygpvar1
151. 1 M m Vr2---->      enddo
152. 1 M m----->      enddo
```

# Kernel B



```
schdt_array = 0D0
do ifreq=1,nFreq

    schDt = (0D0,0D0)

    do my_igp = 1, ngpown

        if (my_igp .gt. ncouls .or. my_igp .le. 0) cycle

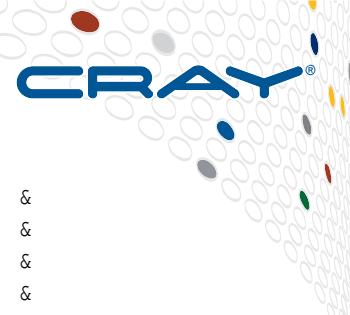
        igmax=ncouls

        schDtt = (0D0,0D0)
        do ig = 1, igmax
            I_epsRggp_int = I_epsR_array(ig,my_igp,ifreq)
            I_epsAggp_int = I_epsA_array(ig,my_igp,ifreq)
            schD=I_epsRggp_int-I_epsAggp_int
            schDtt = schDtt + matngmatmgpD(ig,my_igp)*schD
        enddo
        schdt_array(ifreq) = schdt_array(ifreq) + schDtt
    enddo

enddo
```

# Kernel B

Automatically inserted by Reveal



```
170.    1           ! Directive inserted by Cray Reveal. May be incomplete.
171.    1 M-----< !$OMP parallel do default(none) &
172.    1 M           !$OMP& private (ifreq,ig,igmax,i_epsaggp_int,i_epsrggp_int,my_igp,
173.    1 M           !$OMP&             schd,schdt,schdtt) &
174.    1 M           !$OMP& shared  (i_epsa_array,i_epsr_array,matngmatmgpd,ncouls,nfreq,
175.    1 M           !$OMP&             ngpown,schdt_array) &
176. + 1 M m-----<           do ifreq=1,nFreq
177.    1 M m
178.    1 M m           schDt = (0D0,0D0)
179.    1 M m
180. + 1 M m 4-----<           do my_igp = 1, ngpown
181.    1 M m 4
182.    1 M m 4           if (my_igp .gt. ncouls .or. my_igp .le. 0) cycle
183.    1 M m 4
184.    1 M m 4           igmax=ncouls
185.    1 M m 4
186.    1 M m 4           schDtt = (0D0,0D0)
187.    1 M m 4 Vr3--<           do ig = 1, igmax
188.    1 M m 4 Vr3           I_epsRggp_int = I_epsR_array(ig,my_igp,ifreq)
189.    1 M m 4 Vr3           I_epsAggp_int = I_epsA_array(ig,my_igp,ifreq)
190.    1 M m 4 Vr3           schD=I_epsRggp_int-I_epsAggp_int
191.    1 M m 4 Vr3           schDtt = schDtt + matngmatmgpD(ig,my_igp)*schD
192.    1 M m 4 Vr3-->           enddo
193.    1 M m 4           schdt_array(ifreq) = schdt_array(ifreq) + schDtt
194.    1 M m 4----->           enddo
195.    1 M m
196.    1 M m----->>           enddo
```

COMPUTE

STORE

ANALYZE

# Kernel C



```
do ifreq=1,nFreq
    schDt = schDt_array(ifreq)
    cedifft_zb = dFreqGrid(ifreq)
    cedifft_coh = CMPLX(cedifft_zb,0D0) - dFreqBrd(ifreq)

    if (ifreq .ne. 1) then
        cedifft_zb_right = cedifft_zb
        cedifft_zb_left = dFreqGrid(ifreq-1)
        schDt_right = schDt
        schDt_left = schDt_array(ifreq-1)
        schDt_avg = 0.5D0 * ( schDt_right + schDt_left )
        schDt_lin = schDt_right - schDt_left
        schDt_lin2 = schDt_lin/(cedifft_zb_right-cedifft_zb_left)
    endif

! The below two lines are for sigma1 and sigma3
    if (ifreq .ne. nFreq) then
        schDi(:) = schDi(:) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(:)-cedifft_coh)
        schDi_corb(:) = schDi_corb(:) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(:)-cedifft_cor)
    endif
    if (ifreq .ne. 1) then
        do iw = 1, nfreqeval
!These lines are for sigma2
        intfact=abs((wxi(iw)-cedifft_zb_right)/(wxi(iw)-cedifft_zb_left))
        if (intfact .lt. 1d-4) intfact = 1d-4
        if (intfact .gt. 1d4) intfact = 1d4
        intfact = -log(intfact)
        sch2Di(iw) = sch2Di(iw) - CMPLX(0.d0,prefactor) * schDt_avg * intfact
```

# Kernel C



```
!These lines are for sigma4
    if (flag_occ) then
        intfact=abs((wx(iw)+cedifft_zb_right)/(wx(iw)+cedifft_zb_left))
        if (intfact .lt. 1d-4) intfact = 1d-4
        if (intfact .gt. 1d4) intfact = 1d4
        intfact = log(intfact)
        schDt_lin3 = (schDt_left + schDt_lin2*(-wx(iw)-cedifft_zb_left))*intfact
    else
        schDt_lin3 = (schDt_left + schDt_lin2*(wx(iw)-cedifft_zb_left))*intfact
    endif
    schDt_lin3 = schDt_lin3 + schDt_lin
    schDi_cor(iw) = schDi_cor(iw) - CMPLX(0.d0,prefactor) * schDt_lin3
enddo
endif
enddo
```

# Kernel C



```
202. + 1 2-----<      do ifreq=1,nFreq
203.   1 2
204.   1 2          schDt = schDt_array(ifreq)
205.   1 2
206.   1 2          cedifft_zb = dFreqGrid(ifreq)
207.   1 2          cedifft_coh = CMPLX(cedifft_zb,0D0)- dFreqBrd(ifreq)
208.   1 2
209.   1 2          if (ifreq .ne. 1) then
210.   1 2              cedifft_zb_right = cedifft_zb
211.   1 2              cedifft_zb_left = dFreqGrid(ifreq-1)
212.   1 2              schDt_right = schDt
213.   1 2              schDt_left = schDt_array(ifreq-1)
214.   1 2              schDt_avg = 0.5D0 * ( schDt_right + schDt_left )
215.   1 2              schDt_lin = schDt_right - schDt_left
216.   1 2              schDt_lin2 = schDt_lin/(cedifft_zb_right-cedifft_zb_left)
217.   1 2          endif
218.   1 2
219.   1 2          ! The below two lines are for sigma1 and sigma3
220.   1 2          if (ifreq .ne. nFreq) then
221.   1 2              schDi(:) = schDi(:) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(:)-cedifft_coh)
222.   1 2              schDi_corb(:) = schDi_corb(:) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(:)-cedifft_cor)
223.   1 2          endif
224.   1 2          if (ifreq .ne. 1) then
```

# Kernel C

Automatically inserted by Reveal



```
225. 1 2           ! Directive inserted by Cray Reveal. May be incomplete.
226. 1 2 M-----< !$OMP parallel do default(none)
227. 1 2 M           !$OMP&    private (intfact,iw,schdt_lin3)           &
228. 1 2 M           !$OMP&    shared   (cedifft_zb_left,cedifft_zb_right,flag_occ,nfrequeval,  &
229. 1 2 M           !$OMP&          prefactor,sch2di,schdi_cor,schdt_avg,schdt_left,      &
230. 1 2 M           !$OMP&          schdt_lin,schdt_lin2,wxi)                    &
231. 1 2 M mV----<           do iw = 1, nfrequeval
232. 1 2 M mV       !These lines are for sigma2
233. 1 2 M mV           intfact=abs((wxi(iw)-cedifft_zb_right)/(wxi(iw)-cedifft_zb_left))
234. 1 2 M mV           if (intfact .lt. 1d-4) intfact = 1d-4
235. 1 2 M mV           if (intfact .gt. 1d4) intfact = 1d4
236. 1 2 M mV           intfact = -log(intfact)
237. 1 2 M mV           sch2Di(iw) = sch2Di(iw) - CMPLX(0.d0,prefactor) * schDt_avg * intfact
238. 1 2 M mV       !These lines are for sigma4
239. 1 2 M mV           if (flag_occ) then
240. 1 2 M mV               intfact=abs((wxi(iw)+cedifft_zb_right)/(wxi(iw)+cedifft_zb_left))
241. 1 2 M mV               if (intfact .lt. 1d-4) intfact = 1d-4
242. 1 2 M mV               if (intfact .gt. 1d4) intfact = 1d4
243. 1 2 M mV               intfact = log(intfact)
244. 1 2 M mV               schDt_lin3 = (schDt_left + schDt_lin2*(-wxi(iw)-cedifft_zb_left))*intfact
245. 1 2 M mV           else
246. 1 2 M mV               schDt_lin3 = (schDt_left + schDt_lin2*(wxi(iw)-cedifft_zb_left))*intfact
247. 1 2 M mV           endif
248. 1 2 M mV               schDt_lin3 = schDt_lin3 + schDt_lin
249. 1 2 M mV               schDi_cor(iw) = schDi_cor(iw) - CMPLX(0.d0,prefactor) * schDt_lin3
250. 1 2 M mV---->           enddo
251. 1 2           endif
252. 1 2----->           enddo
```

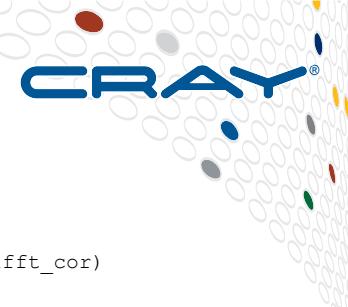
# Kernel C

Automatically modified by Levesque



```
192. 1           ! Directive inserted by Cray Reveal. May be incomplete.
193. 1 M-----< !$OMP parallel do default(none)
194. 1 M           !$OMP&    private (cedifft_coh,cedifft_zb,cedifft_zb_left,
195. 1 M           !$OMP&      cedifft_zb_right,ifreq,intfact,iw,schdt,schdt_avg,
196. 1 M           !$OMP&      schdt_left,schdt_lin,schdt_lin2,schdt_lin3,
197. 1 M           !$OMP&      schdt_right)
198. 1 M           !$OMP&    shared  (cedifft_cor,dfreqbrd,dfreqgrid,flag_occ,nfreq,
199. 1 M           !$OMP&      nfreqeval,pref,prefactor,sch2di,schdi,schdi_cor,
200. 1 M           !$OMP&      schdi_corb,schdt_array,wxi)
201. 1 M mV-----<
202. 1 M mV 4-----<
203. 1 M mV 4
204. 1 M mV 4           schDt = schDt_array(ifreq)
205. 1 M mV 4
206. 1 M mV 4           cedifft_zb = dFreqGrid(ifreq)
207. 1 M mV 4           cedifft_coh = CMPLX(cedifft_zb,0D0)- dfreqBrd(ifreq)
208. 1 M mV 4
209. 1 M mV 4           if (ifreq .ne. 1) then
210. 1 M mV 4               cedifft_zb_right = cedifft_zb
211. 1 M mV 4               cedifft_zb_left = dFreqGrid(ifreq-1)
212. 1 M mV 4               schDt_right = schDt
213. 1 M mV 4               schDt_left = schDt_array(ifreq-1)
214. 1 M mV 4               schDt_avg = 0.5D0 * ( schDt_right + schDt_left )
215. 1 M mV 4               schDt_lin = schDt_right - schDt_left
216. 1 M mV 4               schDt_lin2 = schDt_lin/(cedifft_zb_right-cedifft_zb_left)
217. 1 M mV 4           endif
218. 1 M mV 4
219. 1 M mV 4           ! The below two lines are for sigma1 and sigma3
220. 1 M mV 4           if (ifreq .ne. nFreq) then
221. 1 M mV 4               schDi(iw) = schDi(iw) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(iw)-cedifft_coh)
222. 1 M mV 4               schDi_corb(iw) = schDi_corb(iw) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(iw)-cedifft_cor)
223. 1 M mV 4           endif
224. 1 M mV 4           if(ifreq.ne.1)then
```

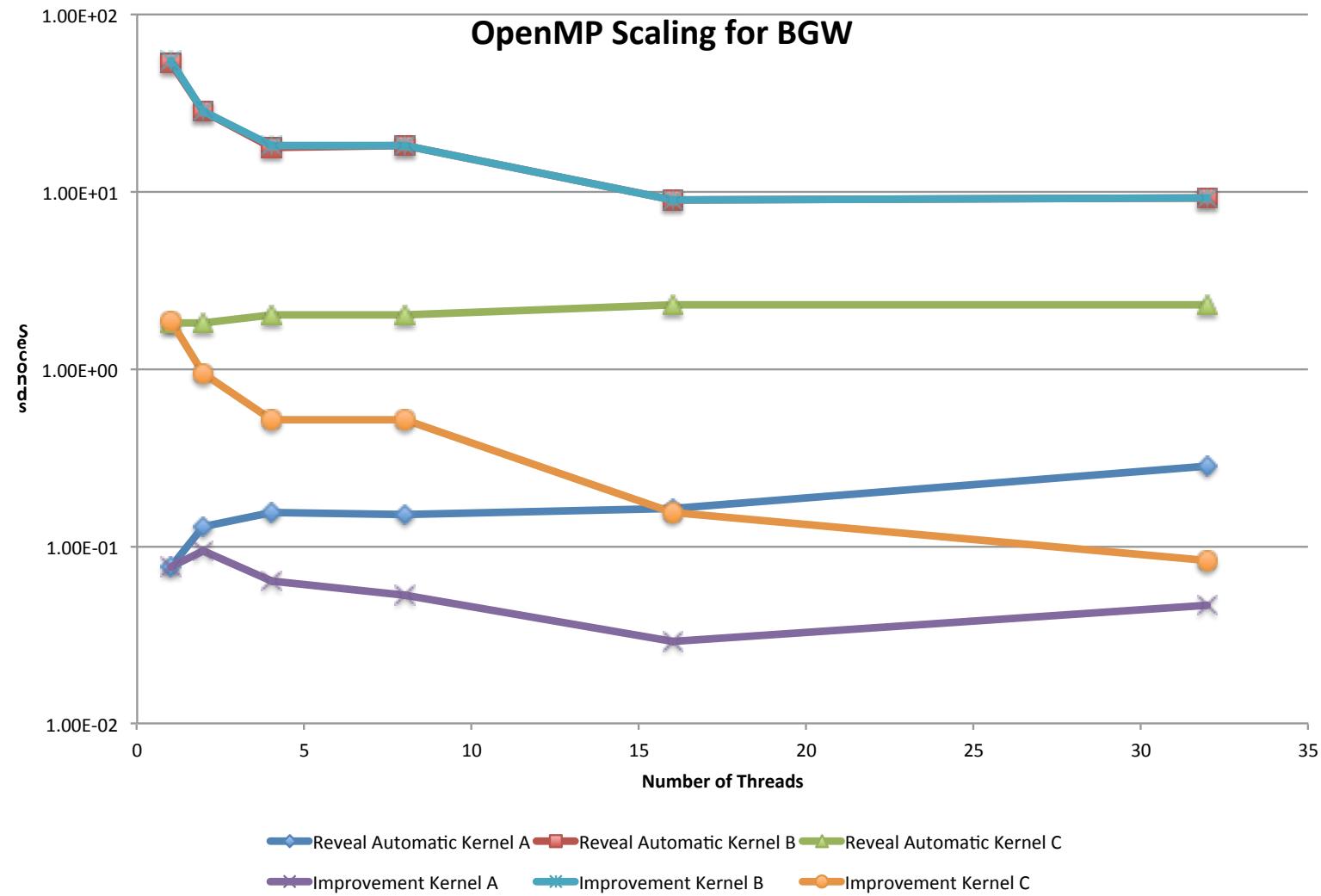
# Kernel C



```
219.    1 M mV 4      ! The below two lines are for sigma1 and sigma3
220.    1 M mV 4      if (ifreq .ne. nFreq) then
221.    1 M mV 4          schDi(iw) = schDi(iw) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(iw)-cediffit_coh)
222.    1 M mV 4          schDi_corb(iw) = schDi_corb(iw) - CMPLX(0.d0,pref(ifreq)) * schDt / ( wxi(iw)-cediffit_cor)
223.    1 M mV 4      endif
224.    1 M mV 4      if(ifreq.ne.1)then
225.    1 M mV 4          !These lines are for sigma2
226.    1 M mV 4              intfact=abs((wxi(iw)-cediffit_zb_right)/(wxi(iw)-cediffit_zb_left))
227.    1 M mV 4              if (intfact .lt. 1d-4) intfact = 1d-4
228.    1 M mV 4              if (intfact .gt. 1d4) intfact = 1d4
229.    1 M mV 4              intfact = -log(intfact)
230.    1 M mV 4          sch2Di(iw) = sch2Di(iw) - CMPLX(0.d0,prefactor) * schDt_avg * intfact
231.    1 M mV 4      !These lines are for sigma4
232.    1 M mV 4      if (flag_occ) then
233.    1 M mV 4          intfact=abs((wxi(iw)+cediffit_zb_right)/(wxi(iw)+cediffit_zb_left))
234.    1 M mV 4          if (intfact .lt. 1d-4) intfact = 1d-4
235.    1 M mV 4          if (intfact .gt. 1d4) intfact = 1d4
236.    1 M mV 4          intfact = log(intfact)
237.    1 M mV 4          schDt_lin3 = (schDt_left + schDt_lin2*(-wxi(iw)-cediffit_zb_left))*intfact
238.    1 M mV 4      else
239.    1 M mV 4          schDt_lin3 = (schDt_left + schDt_lin2*(wxi(iw)-cediffit_zb_left))*intfact
240.    1 M mV 4      endif
241.    1 M mV 4          schDt_lin3 = schDt_lin3 + schDt_lin
242.    1 M mV 4          schDi_coriw) = schDi_coriw) - CMPLX(0.d0,prefactor) * schDt_lin3
243.    1 M mV 4      endif
244.    1 M mV 4----->      enddo
245.    1 M mV----->>      enddo
```

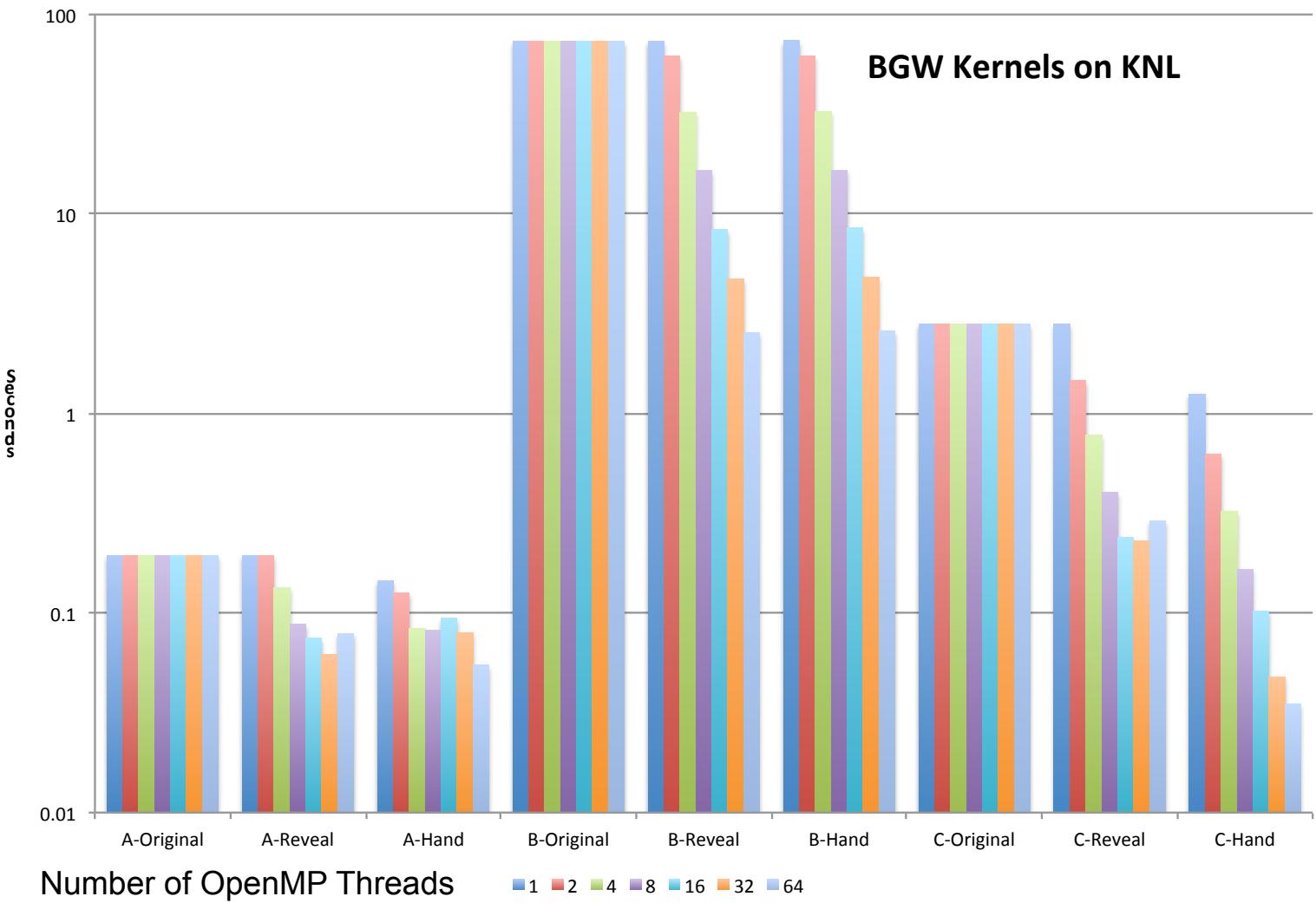


## OpenMP Scaling for BGW



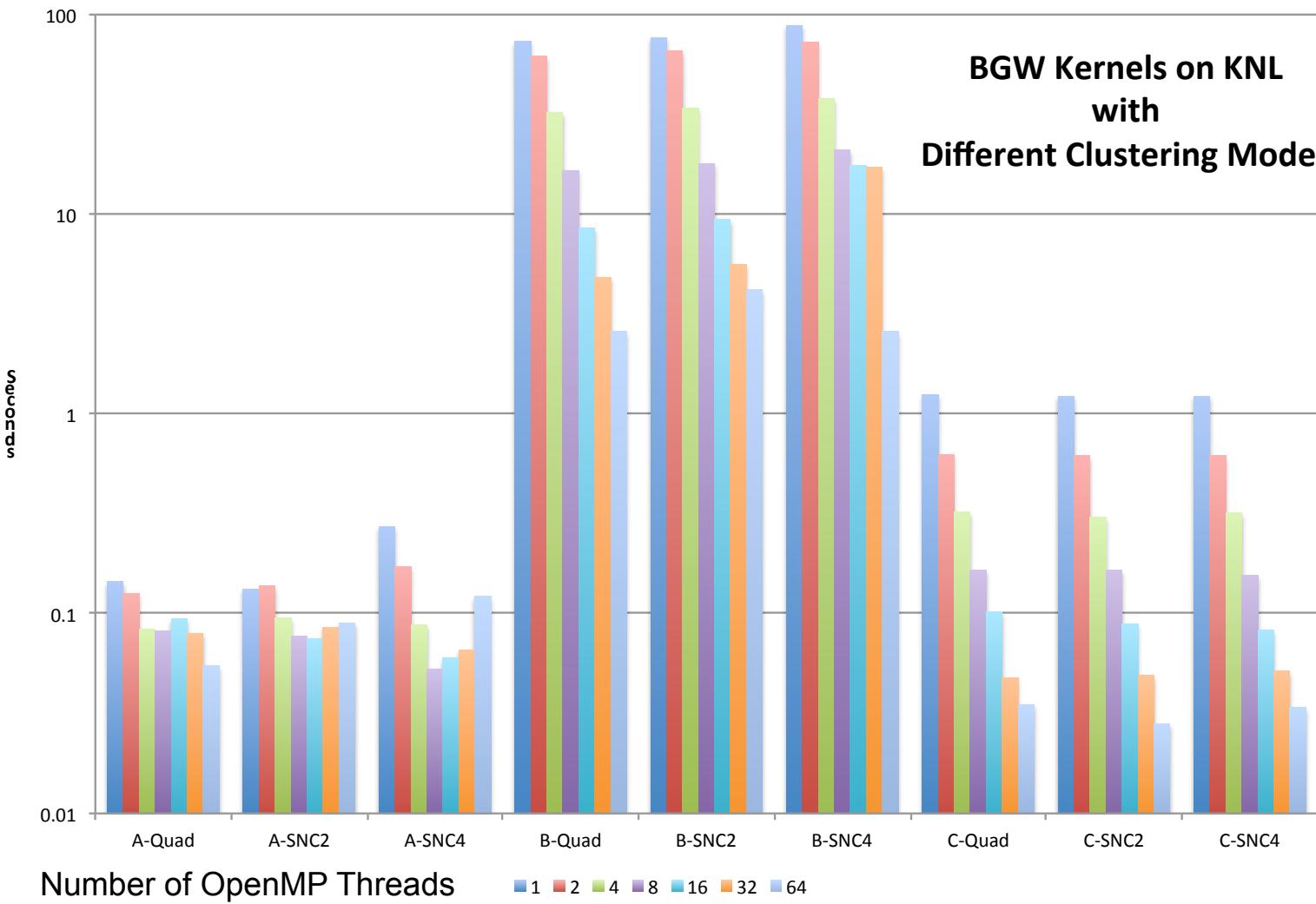


## BGW Kernels on KNL





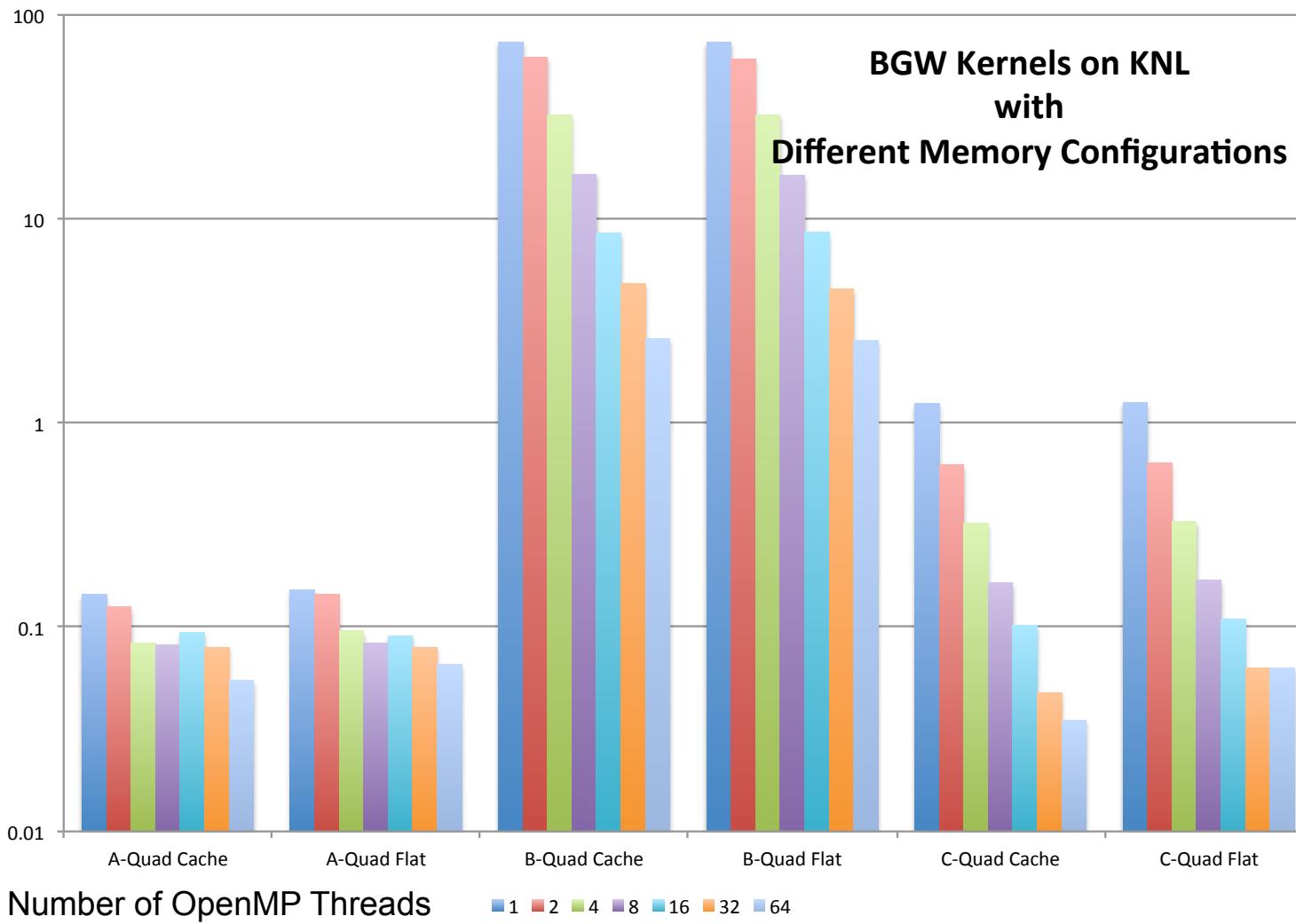
## BGW Kernels on KNL with Different Clustering Modes





## BGW Kernels on KNL with Different Memory Configurations

Seconds





# Nekbone

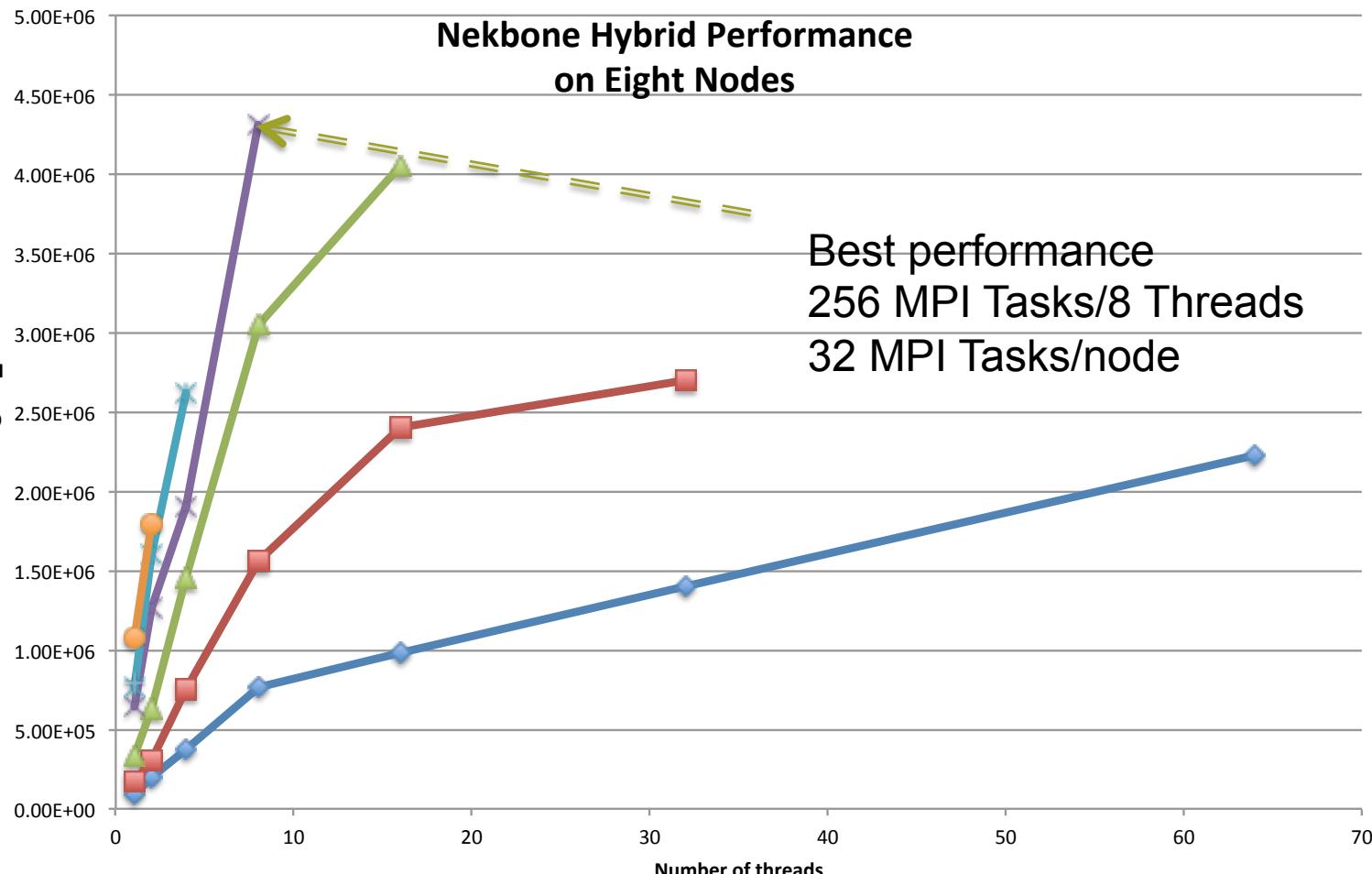
High order, incompressible  
Navier-Stokes solver based on the spectral  
element method



## Nekbone Hybrid Performance on Eight Nodes

Best performance  
256 MPI Tasks/8 Threads  
32 MPI Tasks/node

MFLOPS



Number of MPI Tasks across 8 nodes

32 64 128 256 512 1024 2048



# Comment on Application

- **Excellent OpenMP**

- SPMD style with consistent memory access pattern
  - Come to CUG and see my tutorial
- Only example that shows Hybrid a clear winner over all MPI



# S3D

Flow solver for performing direct numerical simulation (DNS) of turbulent combustion.

# Details of Test

- **OpenACC/OpenMP code written for Titan**
  - Zero modifications from Titan code to KNL
  - OpenACC – OpenMP controlled by #IFDEF
- **Test Problems ( Small and Medium fit in MCDRAM)**
  - Small Grid (6.5 GB/node)
    - Size of Titan Acceptance test
    - 192\*192\*192 on 16 nodes
  - Medium Grid (10 GB/node)
    - 384\*384\*192 on 16 nodes
  - Large Grid (24 GB/node)
    - 768\*384\*384 on 16 nodes
- **Ran on 8,16,32,64 ,128 and 256 nodes**
  - Weak Scaling employed for number of MPI tasks

# Comment on Application



## Reasonably Good OpenMP

- It is the traditional approach with high level Parallel DO
- One Parallel loop may not have same memory access pattern as the next
  - Could cause caching issues
- Could be improved with SPMD style OpenMP to maintain constant memory access across parallel loops

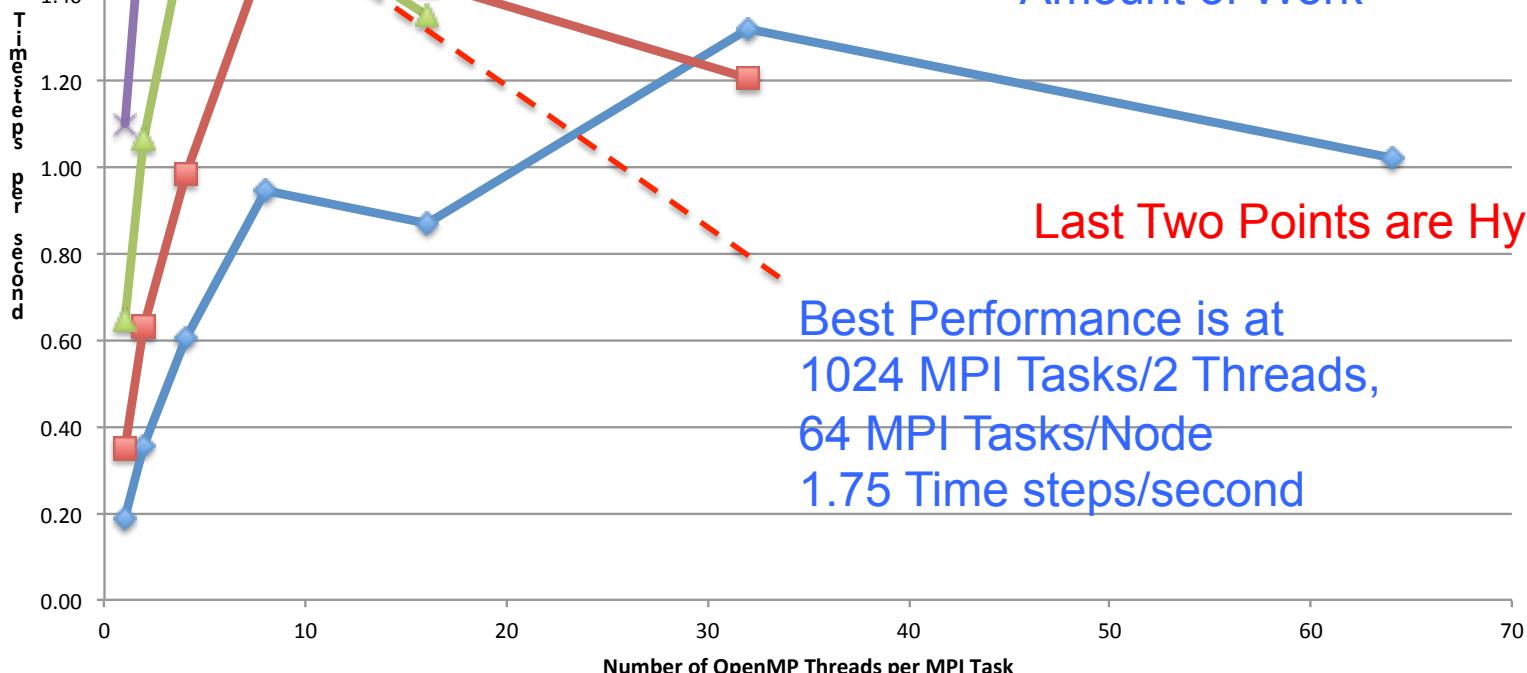


## S3D - Small Grid - Hybrid across 16 nodes for Several MPI counts

Weak Scaling – each  
MPI Task does equal  
Amount of Work

Last Two Points are Hyper-threads

Best Performance is at  
1024 MPI Tasks/2 Threads,  
64 MPI Tasks/Node  
1.75 Time steps/second



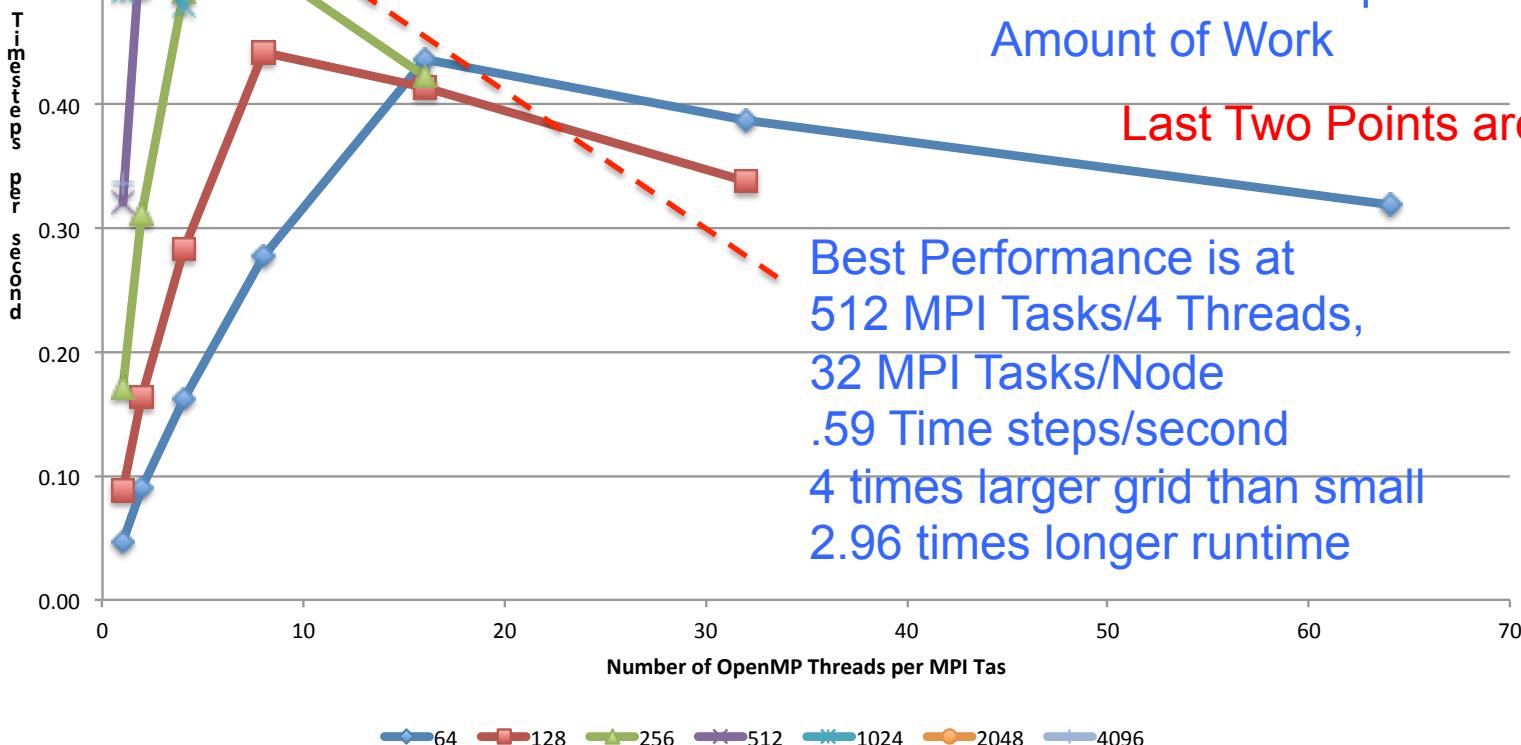


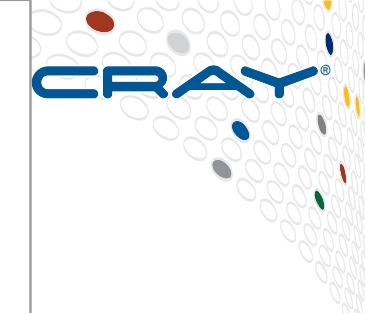
# S3D - Medium Grid - Hybrid across 16 nodes for Several MPI counts

Weak Scaling – each  
MPI Task does equal  
Amount of Work

Last Two Points are Hyper-threads

Best Performance is at  
512 MPI Tasks/4 Threads,  
32 MPI Tasks/Node  
.59 Time steps/second  
4 times larger grid than small  
2.96 times longer runtime



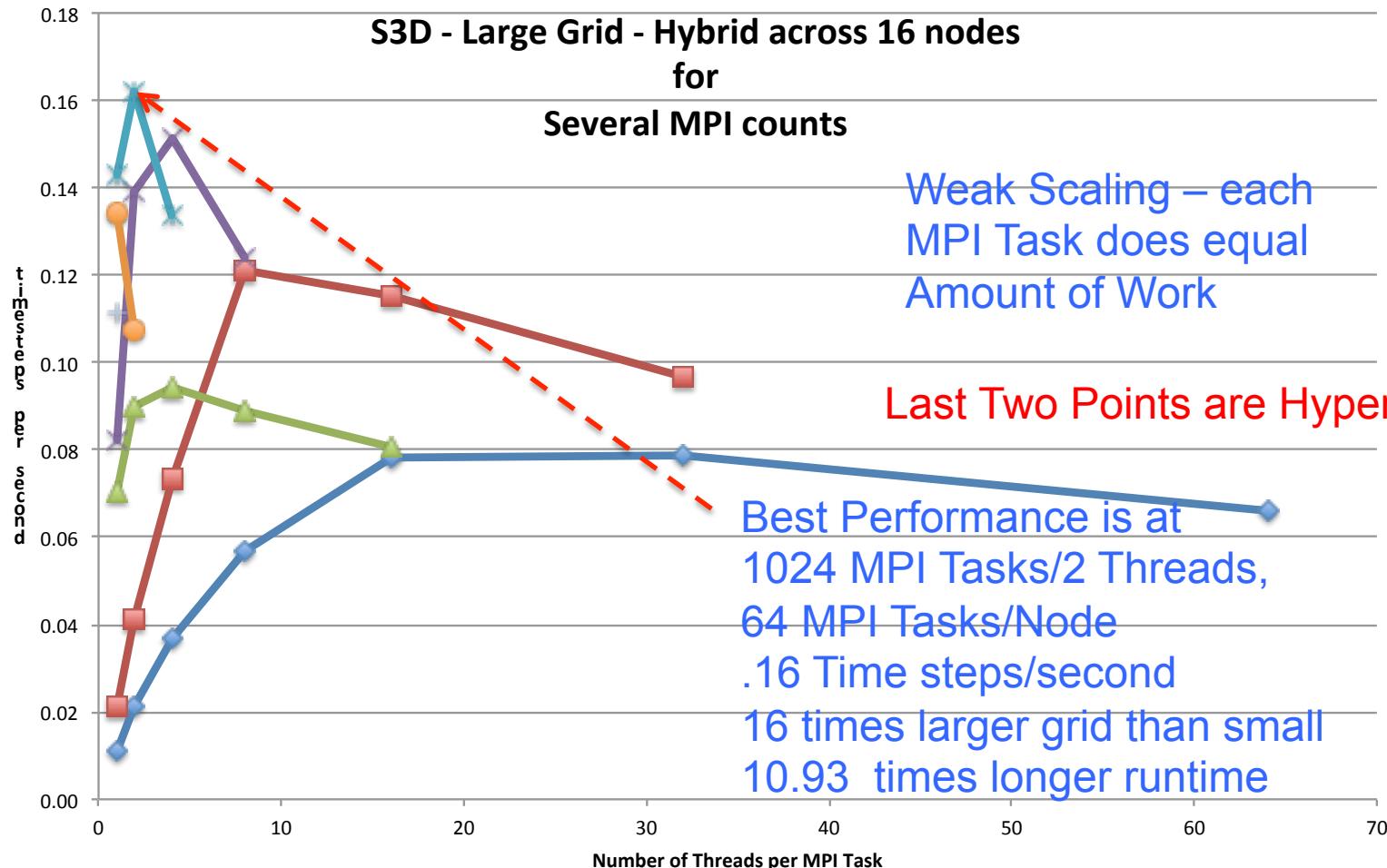


## S3D - Large Grid - Hybrid across 16 nodes for Several MPI counts

Weak Scaling – each  
MPI Task does equal  
Amount of Work

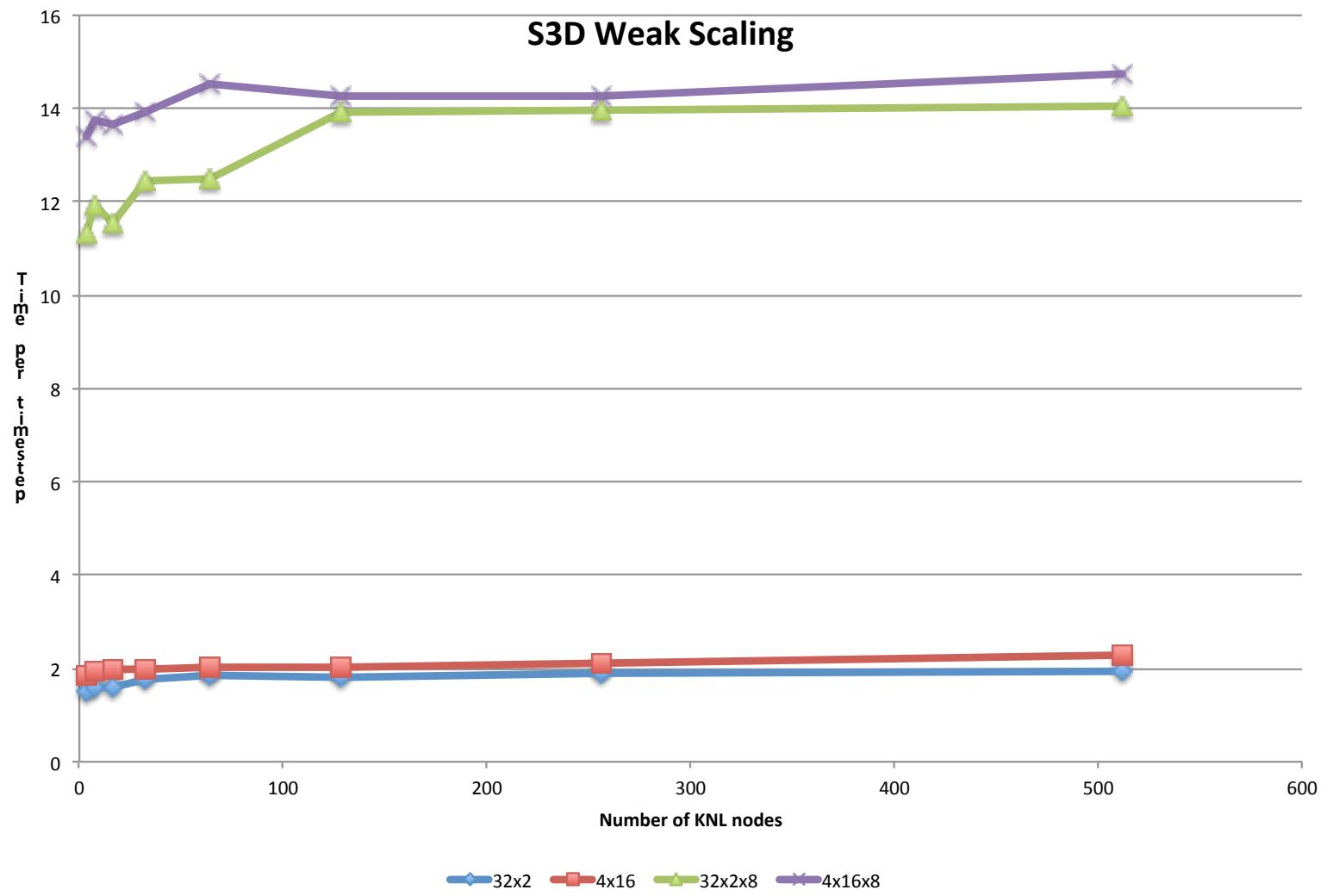
Last Two Points are Hyper-threads

Best Performance is at  
1024 MPI Tasks/2 Threads,  
64 MPI Tasks/Node  
.16 Time steps/second  
16 times larger grid than small  
10.93 times longer runtime





## S3D Weak Scaling



# Conclusions



- **MPI works very well on KNL**
  - MPI sweet spot is always greater than 8, most times 32-64
- **Threading is extremely sensitive to NUMA features of the 2-D tile communication grid and Cache Home Agent (CHA)**
- **Advantages of using MCDRAM for cache over flat is highly dependent upon cache friendliness of application.**
  - Himeno not cache friendly - factor of 4-5
  - S3D is cache friendly – factor of 1.5-2.0

# Should you use Hyper-Threads

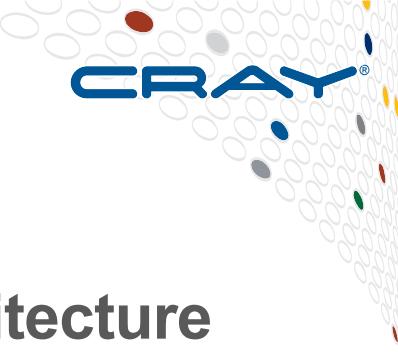


- Not a definite yes or no
  - Extremely code dependent
  - HT = 2 is most likely to be best
    - Caveat – Cray doesn't recommend running MPI across Hyper-Threads
      - Although several applications show best performance on small node counts
  - While HT helps in hiding latency, it can put increase cache pressure on all levels of cache
  - Needs to be tested

# How many MPI ranks/node

- Unless SPMD or extremely good high level OpenMP/threading is used, at least 8 MPI ranks per node seems to be the rule with most applications getting best performance with 64 MPI ranks per node
- Too few MPI ranks will not be able to fully utilize the network bandwidth of number of outstanding references, leaving MPI performance on the table
- Be very careful with single node performance studies
  - Must consider off-node component – especially if the ultimate goal is 1000s of nodes
  - OpenMP and MPI trade-off must be done on as many nodes as possible

# OpenMP on KNL



- OpenMP is best when used within a tile
- KNL certainly has a Non-Uniform Memory Architecture
- OpenMP would be best way to employ Hyper-threads



# Conclusions

- **OpenMP is a tool, not a requirement**
  - First find a sweet spot with MPI than add OpenMP
  - Do not think you need to have > 8 OpenMP threads
- **Vectorization is critical to getting good performance on KNL**
- **One must look hard at how and when to use MCDRAM**
- **No good guidance on when Hyperthreading will help**
- **KNL has a significant number of performance “knobs” that can be tuned for an individual code**
  - While lessons learned have narrowed the search space some, a significant numbers of combinations remain



# Vectorization on KNL

- KNL vector register length is twice that of Haswell
- KNL scalar performance is about 1/3 that of Haswell
- KNL performance is highly correlated with good vectorization
- Without good-excellent vectorization, the slow scalar performance will significantly degrade overall performance
- Vectorizing loops with lots of indirection and/or strides will probably be slower than Xeon scalar performance

# MCDRAM is critical to performance

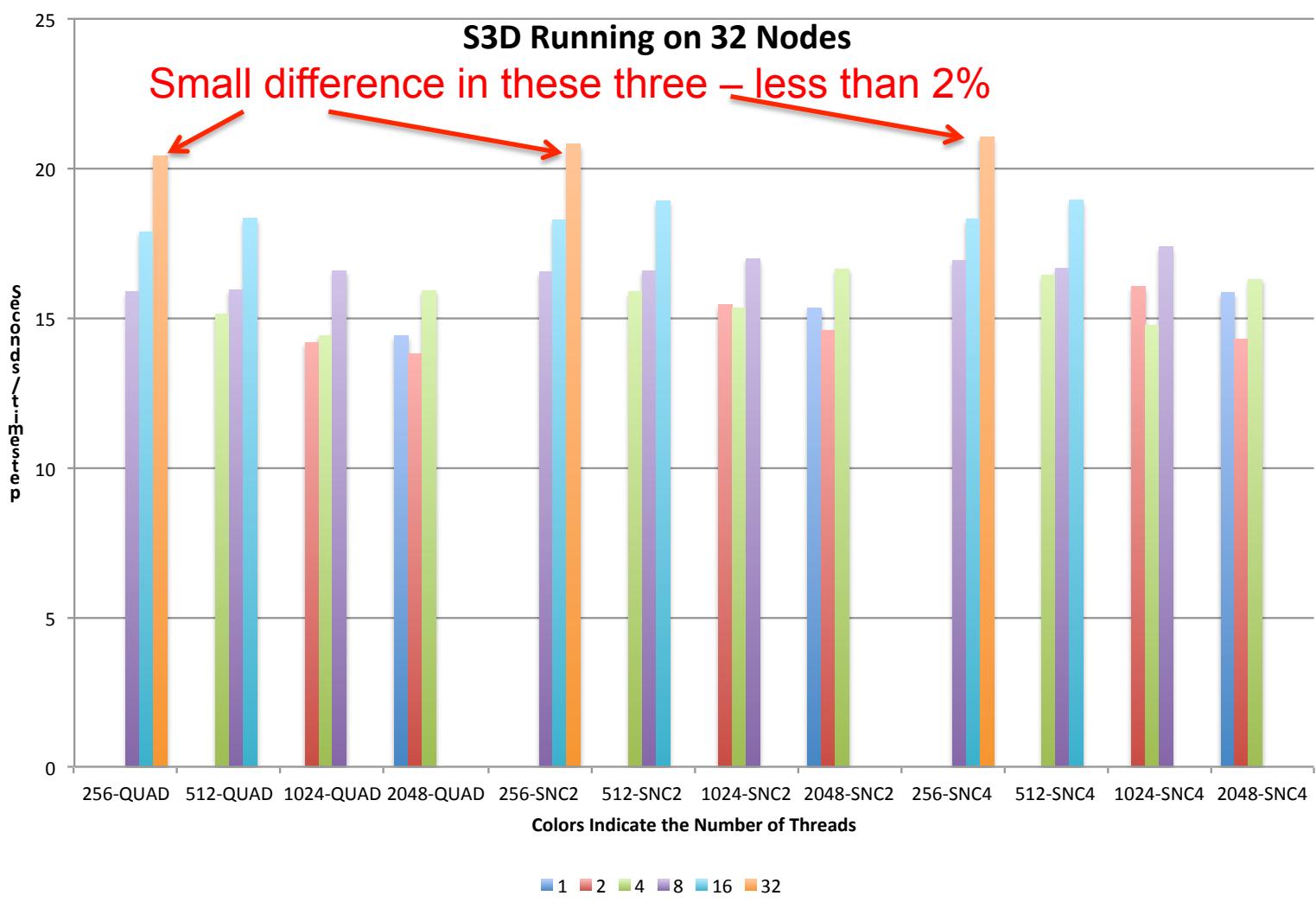


- If an application can fit within MCDRAM, then using flat memory with numactl – membind is the best way to go.
- If an application cannot fit with MCDRAM, then one should consider cache, split, equal and/or flat
  - This is extremely hard, since reserving some of MCDRAM for memory space will reduce last level cache for other variables
  - When using MCDRAM as cache, one must consider/examine issues with the direct mapped cache at scale



## S3D Running on 32 Nodes

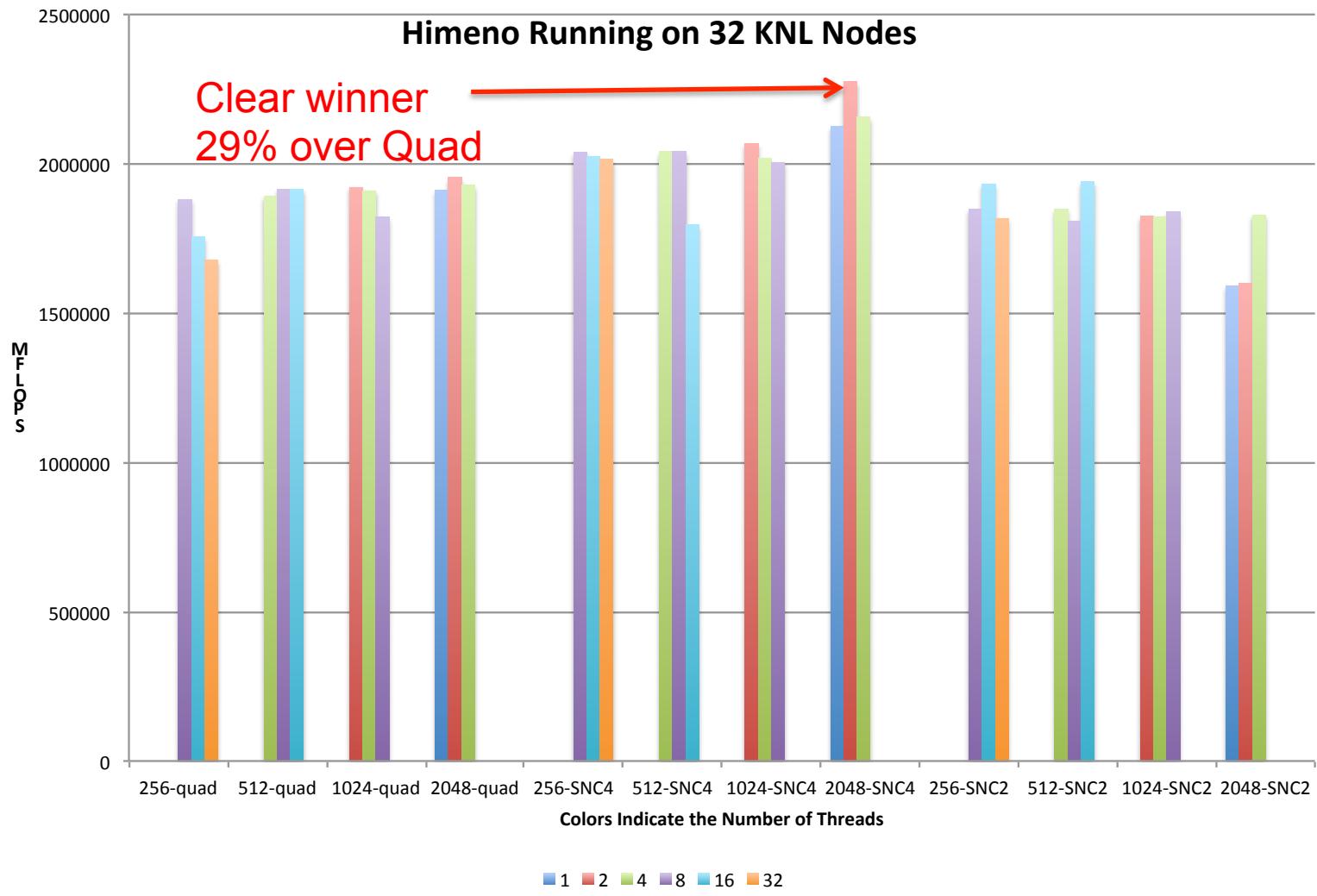
Small difference in these three – less than 2%





## Himeno Running on 32 KNL Nodes

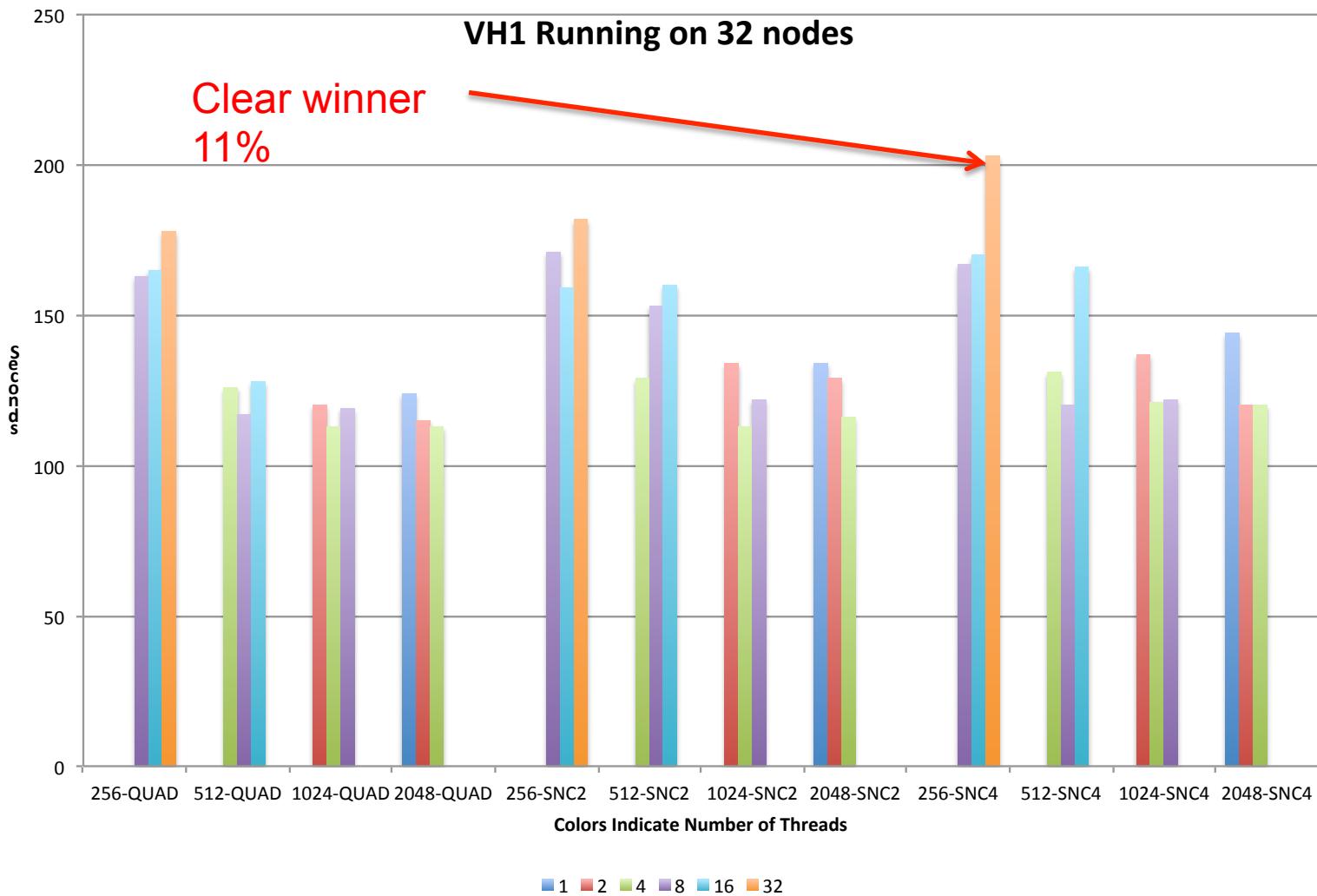
Clear winner  
29% over Quad



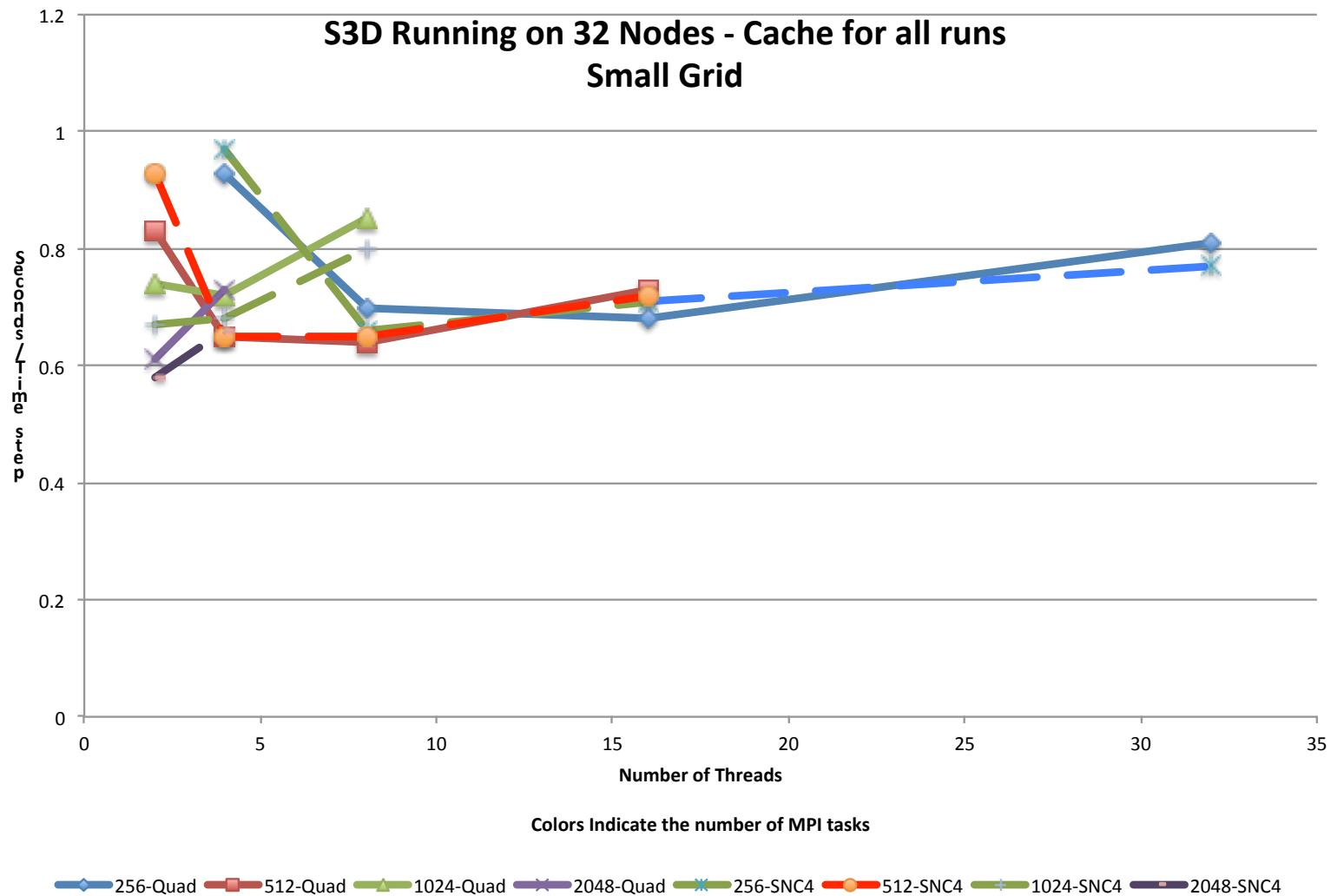


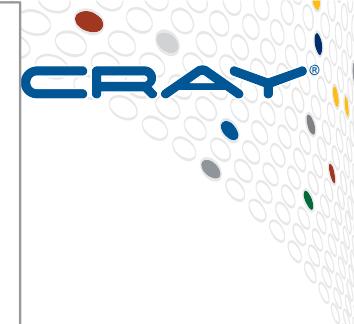
## VH1 Running on 32 nodes

Clear winner  
11%



# S3D Running on 32 Nodes - Cache for all runs Small Grid

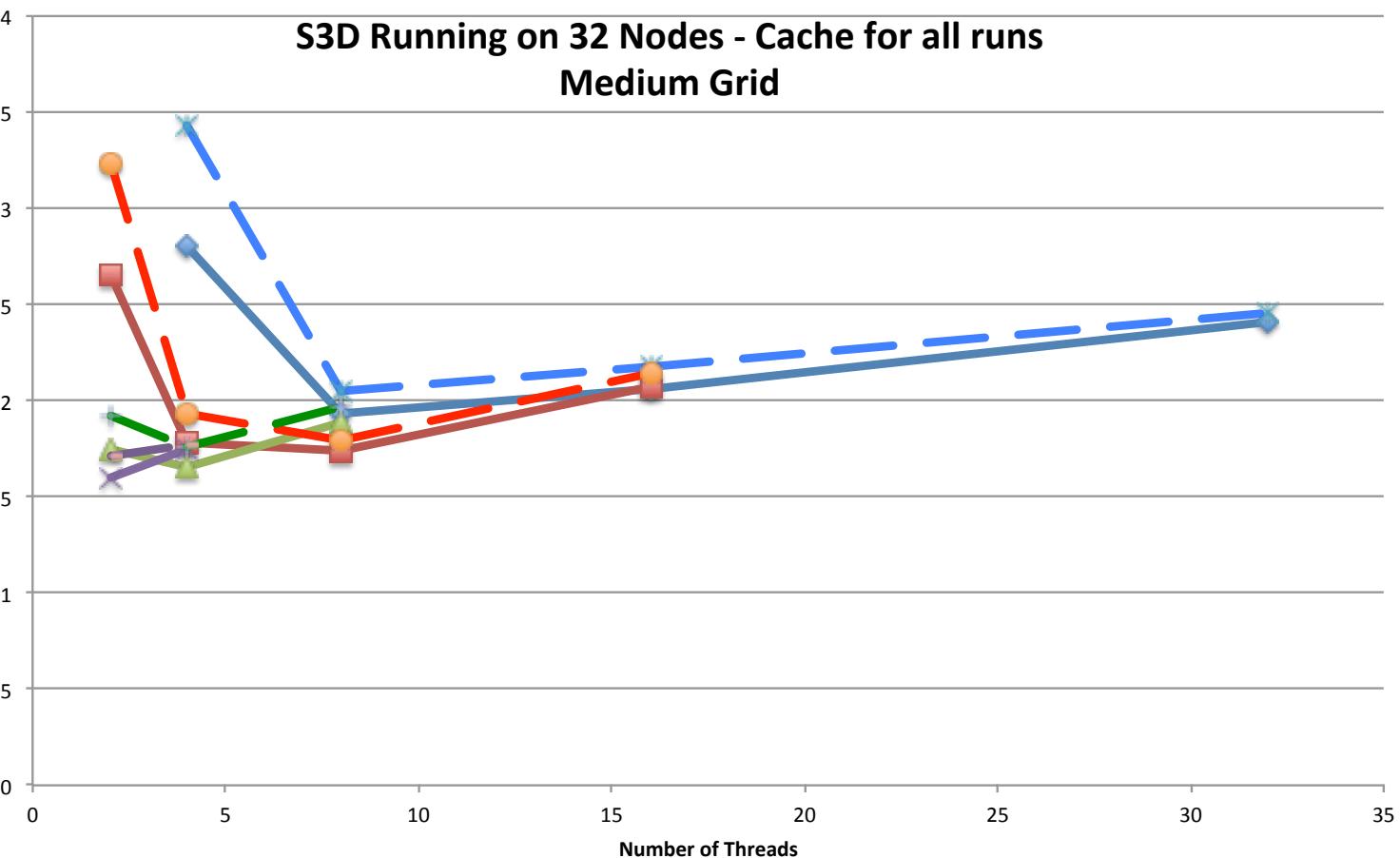




# S3D Running on 32 Nodes - Cache for all runs

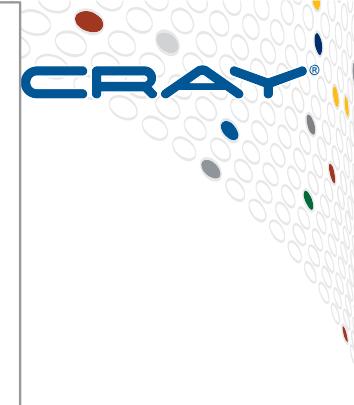
## Medium Grid

Seconds / Time step



Colors Indicate number of MPI Tasks

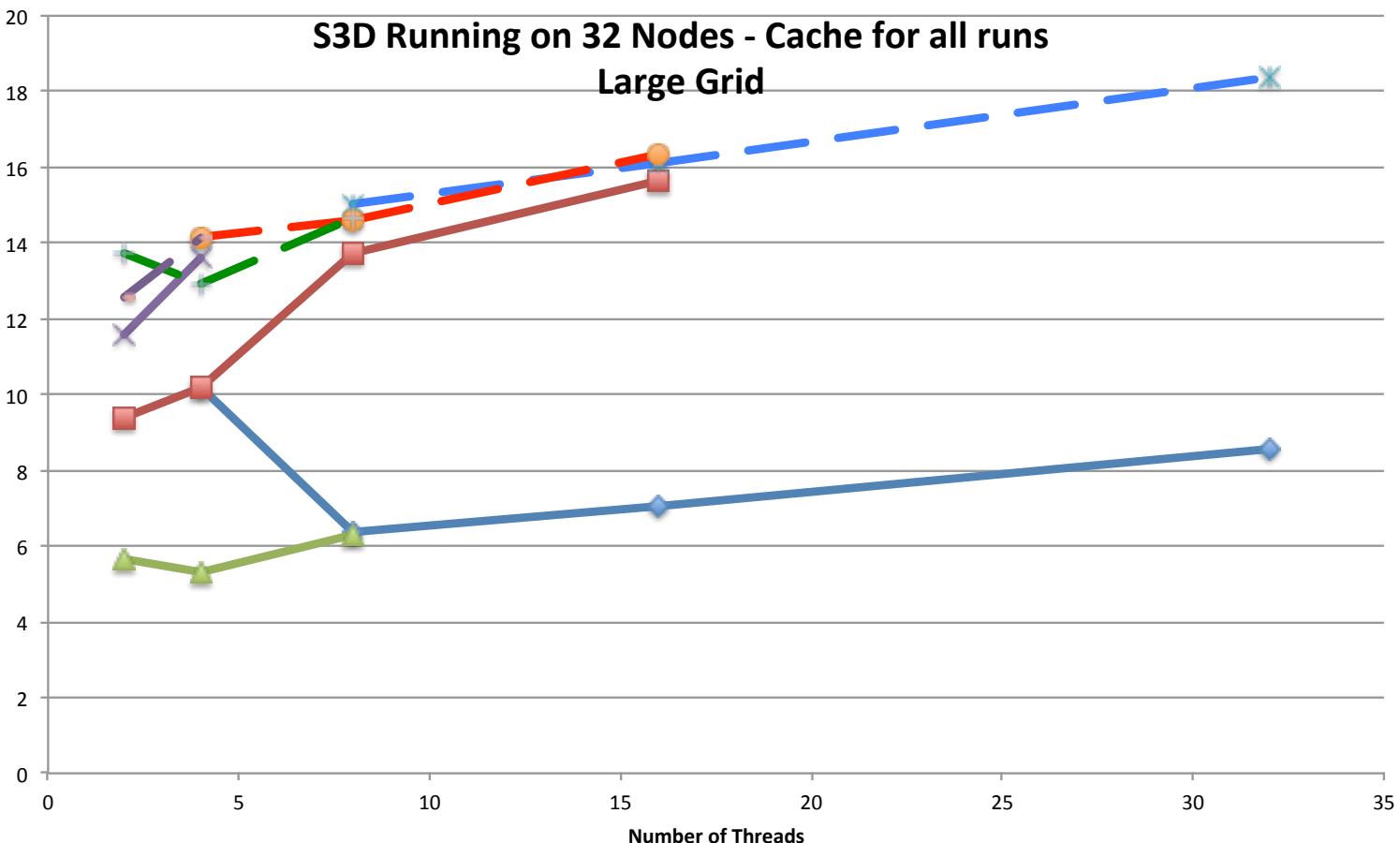
— 256-Quad — 512-Quad — 1024-Quad — 2048-Quad — 256-SNC4 — 512-SNC4 — 1024-SNC4 — 2048-SNC4



# S3D Running on 32 Nodes - Cache for all runs

## Large Grid

Seconds / Timestep



Colors Indicate the Number of MPI tasks

—♦— 256-Quad —■— 512-Quad —▲— 1024-Quad —×— 2048-Quad —●— 256-SNC4 —○— 512-SNC4 —+— 1024-SNC4 —— 2048-SNC4

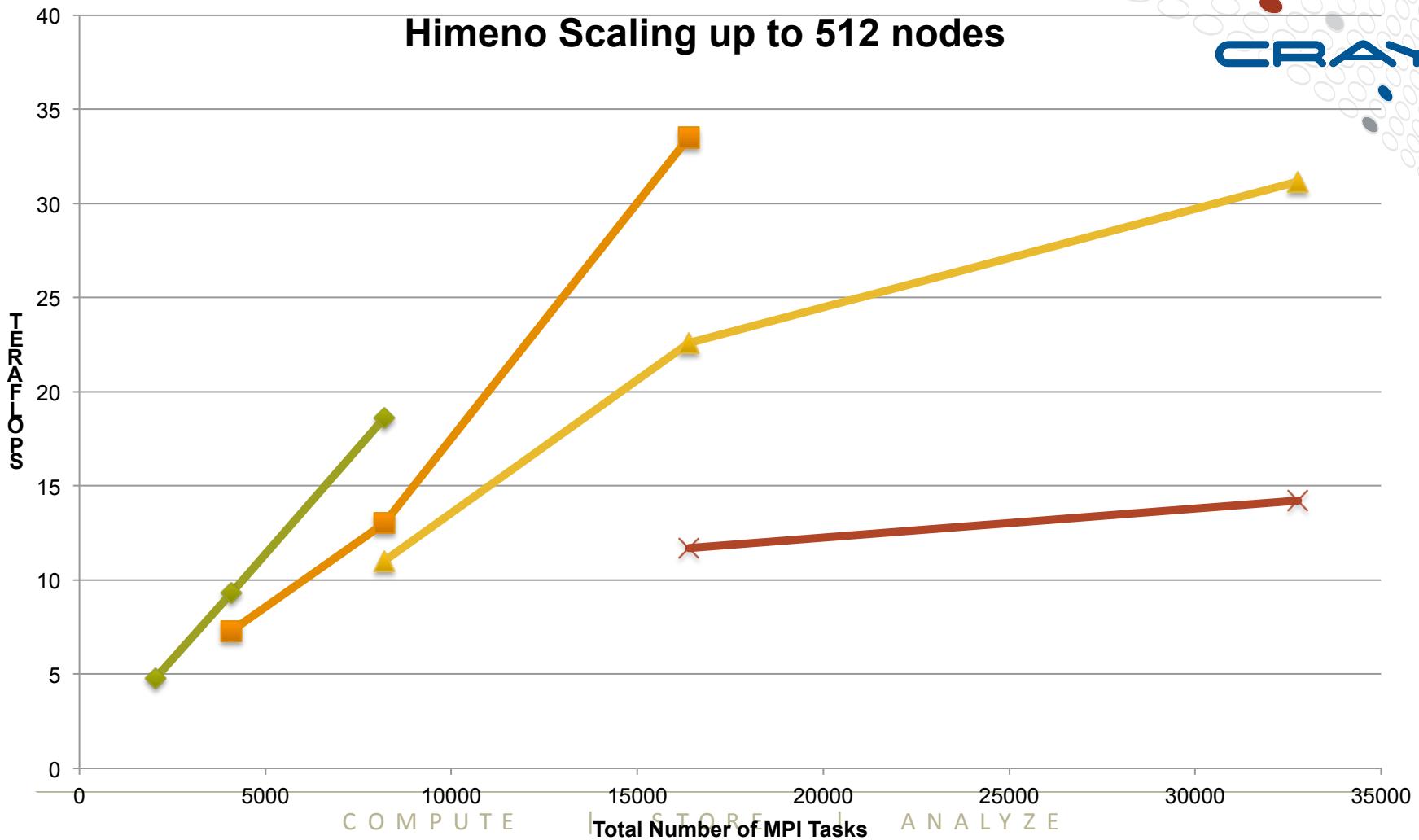
# Examine MPI scaling of existing application across all nodes, all cores



- Regardless of your current threading, examine the scaling of MPI only across all the cores of all the nodes that you would like to use.
- There will be a MPI sweet spot, the point at which MPI stops scaling.
- To date a large majority of the applications have shown that at least one MPI task/tile is the best hybrid configuration
  - 32 (34) MPI tasks – 2,4, or 8 threads
  - 64 (68) MPI tasks – 2,4 threads
- While 128 (136) MPI tasks on a nodes occasionally is faster, Cray does not recommend doing MPI on Hyper-threads

# Himeno Scaling up to 512 nodes

CRAY®

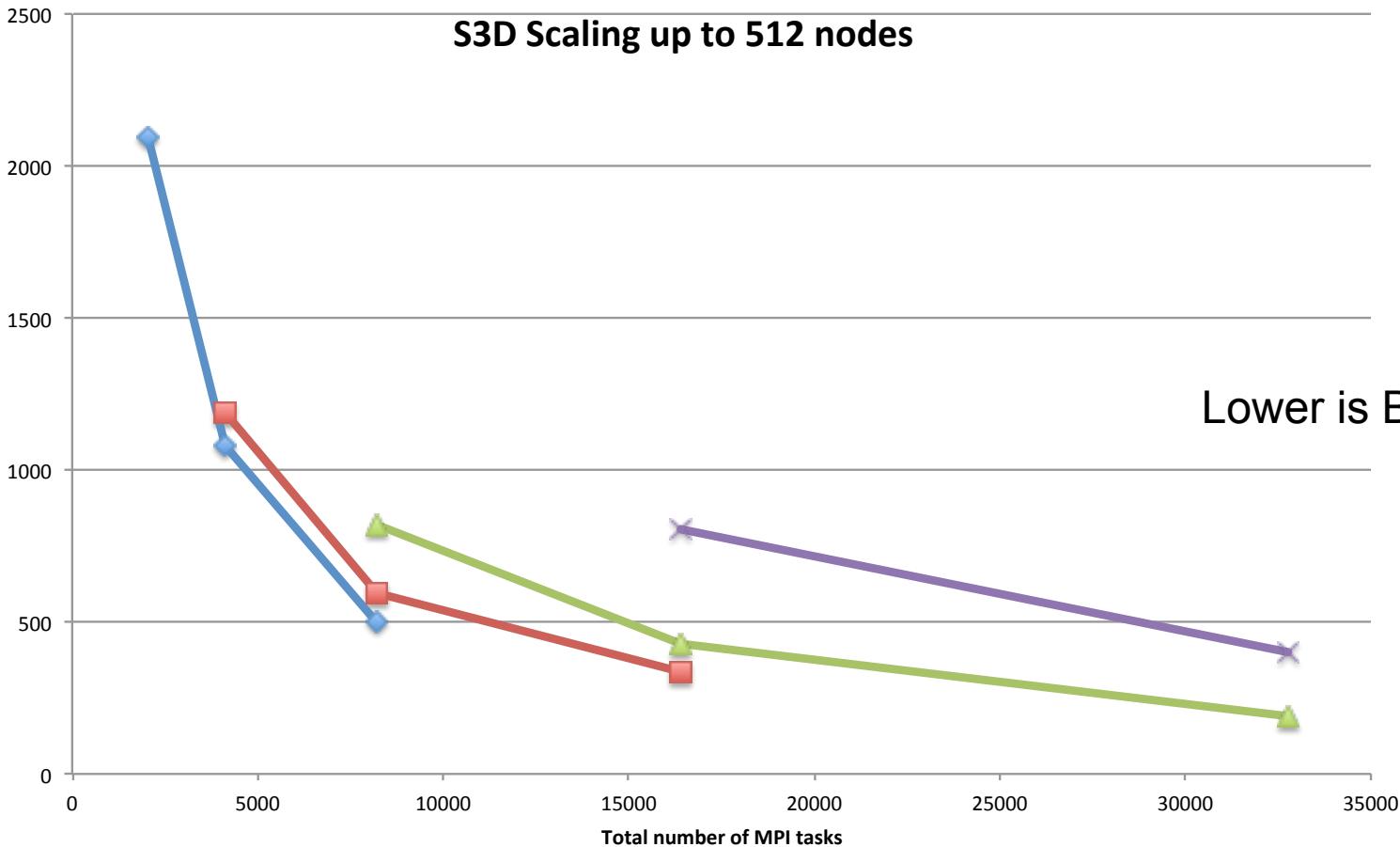




## S3D Scaling up to 512 nodes

Lower is Better

Seconds



Colors indicate the number of MPI tasks/node

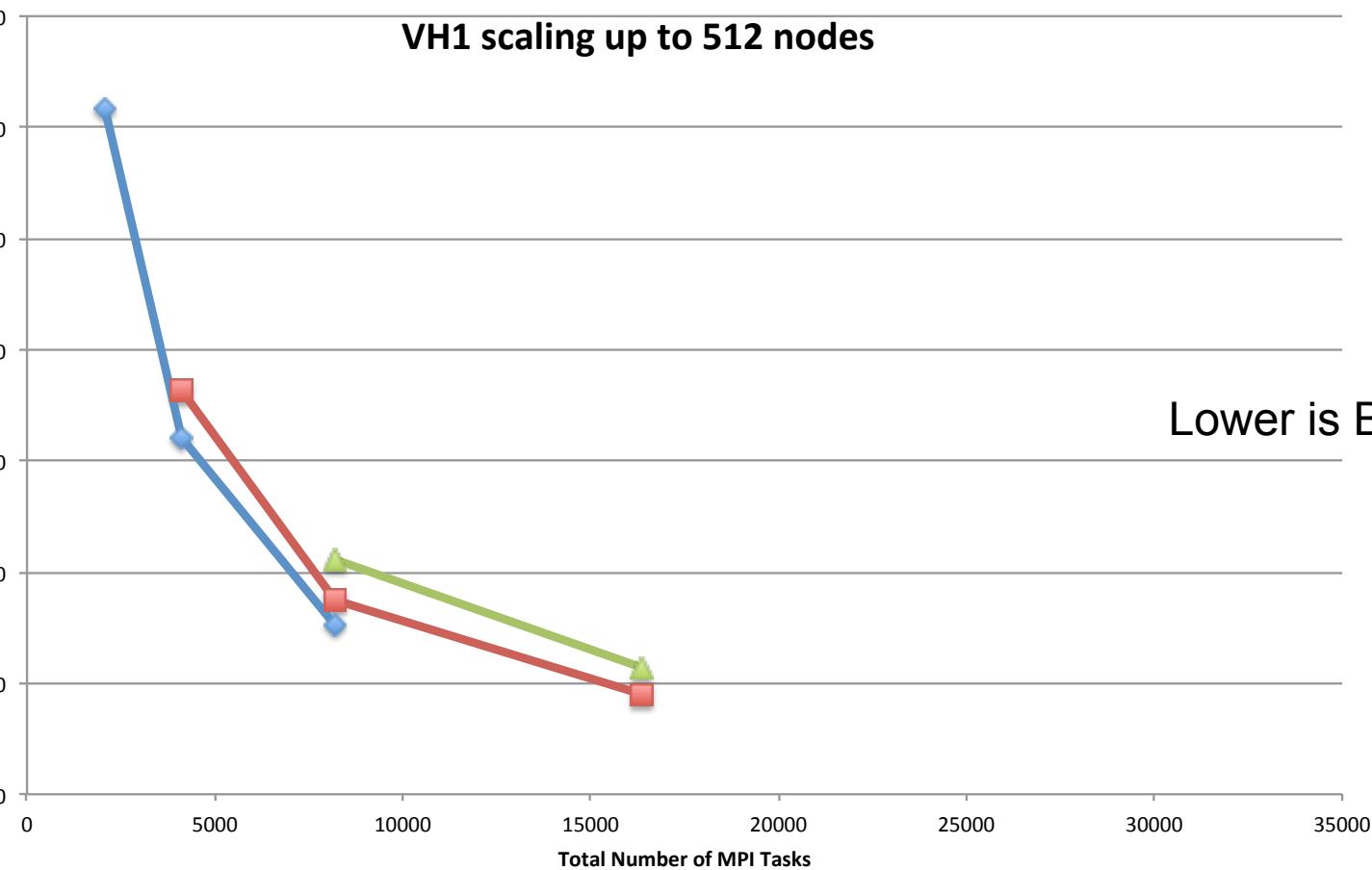
—◆— 16   —■— 32   —▲— 64   —×— 128



## VH1 scaling up to 512 nodes

Seconds

Lower is Better



Colors Indicate number of MPI Tasks/Node

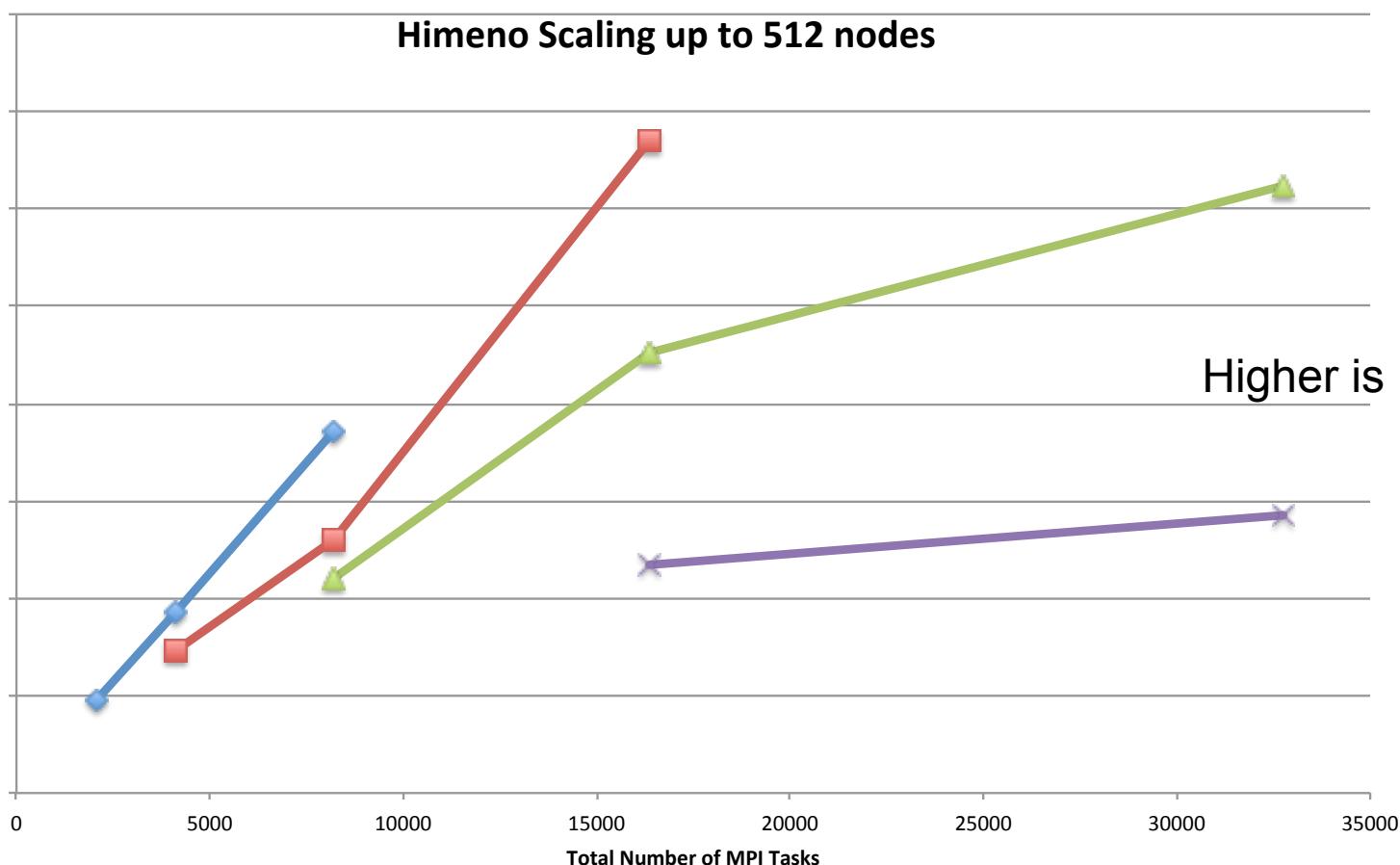
—♦— 16   —■— 32   —▲— 64   —×— 128



## Himeno Scaling up to 512 nodes

TERAFLOPS

Higher is Better



Colors indicate the number of MPI Tasks/node

—♦— 16   —■— 32   —▲— 64   —×— 128

# Investigate load imbalance



C O M P U T E

|

S T O R E

|

A N A L Y Z E



# Scalar Load Imbalance is exaggerated on KNL

Table 1: Profile by Function

Samp%	Samp	Imb.	Imb.	Group
		Samp	Samp%	Function
				PE=HIDE
100.0%	190,766.8	--	--	Total
87.8%	167,553.3	--	--	MPI
52.0%	99,181.7	20,759.3	17.3%	MPI_Waitall
12.6%	24,068.4	23,281.6	49.2%	MPI_Allgather
10.7%	20,407.1	10,906.9	34.9%	MPI_Allreduce
3.3%	6,245.7	1,399.3	18.3%	mpi_waitall
2.2%	4,110.2	48.8	1.2%	mpi_bcast
2.0%	3,795.6	11,579.4	75.4%	MPI_Isend
1.7%	3,298.1	55.9	1.7%	MPI_BARRIER
1.2%	2,199.5	639.5	22.5%	MPI_ALLREDUCE
9.4%	17,949.0	--	--	USER
2.0%	3,887.4	1,543.6	28.4%	ipcress_getmu_all\$ipcress_module_
1.5%	2,782.6	20,929.4	88.4%	hypre_BinarySearch
0.8%	1,594.7	4,774.3	75.0%	hypre_BoomerAMGBuildCoarseOperator
0.5%	939.8	1,084.2	53.6%	hypre_BoomerAMGRelax
0.4%	776.4	4,092.6	84.1%	hypre_BoomerAMGBuildInterp
0.3%	626.9	101.1	13.9%	inside_com3\$derivatives_common\$derivative_module_
0.2%	428.9	302.1	41.4%	hypre_CSRMatrixMatvec
0.2%	391.5	741.5	65.5%	hypre_BoomerAMGCoarsen
0.2%	291.3	59.7	17.0%	derivatives_common\$derivative_module_
0.1%	281.3	80.7	22.3%	teos_iter_3\$teos_module_
0.0%	0.0	1.0	100.0%	hypre_BoomerAMGSetThreshold
2.5%	4,807.5	--	--	ETC
0.7%	1,245.7	493.3	28.4%	EXP_Z
0.4%	760.1	118.9	13.5%	_RTON



# Scalar Load Imbalance is exaggerated on KNL

