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Torben Kling Petersen & Bill Loewe, Cray Inc.

Agenda



- HPC storage performance is under constant strain to deliver required performance
- Modern I/O bound applications tend to be very complex with a mix of large sequential and smaller IOPS bound I/O
- Examine the ability of NXD to improve application performance and manage a mixture of large streaming I/O in conjunction with small, random I/O.

Benefit

Users with complex I/O workloads are always looking for solutions to the problem. NXD is
potentially one of these solutions.

• Experimental setup

- Results
- Summary
- Q&A



NXD Hardware Components



NXD Software Architecture



NXD

- Filter driver implemented as device mapper target driver
- Core library compiled as a Linux kernel module with well defined APIs
- Work at the block layer be transparent to file system and applications
- Core caching function is an OS agnostic portable library with well defined interfaces
- Filter Driver intercept's IO and routes through Cache Management Library for Caching functions

Benchmarking Setup

Single ClusterStor L300N

- 82x Seagate 6TB Enterprise SAS HDDs
- 2x Seagate (ST3200FM0033) SSDs
- High speed interconnect (EDR IB or OmniPath)
- 16 Compute clients on IB or OPA

• Storage Software

- Lustre 2.7.19.12.x8-51 (user app tests run on Lustre 2.5.1)
 - End user app tests used Lustre 2.5.1
- NXD Version 3.1.0.3 (2017.12.20)
 - End user app tests used NXD Version 3.0.10.3

Experimental Setup

1. IOR sequential

IOR -b 32g -t 1m -F -k -m -e -v -v -C -i 5 -o \$OUTFILE

2. Mixed IOR

Sequential IOR: IOR -b 32g -t 1m -F -k -m -e -v -v -C -i 10 -o \$OUTFILE Random IOR: IOR -b 280m -t 4k -F -k -m -e -v -v -C -i 10 -z -o \$OUTFILE

3. Concurrent IOR, increasing load of random

Mixed IOR (above) with 32, 96, 160 and 224 I/O threads respectively

4. SWMR

swmr write ---niter 2500 --testdatafile \$IN_FILE \$OUT_FILE



ANALYZE

NXD / ClusterStor L300N – Write/Rewrite, Aligned I/O

16 nodes, 64 processes total, 1x SSU, (2x SSDs vs. 82x HDDs) running Lustre 2.5.1

Write performance 90 000 80 000 70 000 60 000 S 50 000 O 40 000 30 000 20 000 10 000 0 2 16 32 64 128 1 Transfer Size (KB)

90 000 80 000 70 000 60 000 S 50 000 0 40 000 NytroXD 30 000 GridRAID 20 000 10 000 0 2 32 64 128 8 16 Transfer Size (KB)

Re-write performance

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NXD / ClusterStor L300N – Write/Rewrite, Unaligned I/O

16 nodes, 64 processes total, 1x SSU, (2x SSDs vs. 82x HDDs) running Lustre 2.5.1



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NXD / ClusterStor L300N – Read Performance

16 nodes, 64 processes total, 1x SSU, (2x SSDs vs. 82x HDDs) running Lustre 2.5.1



Unligned I/O

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Experiment 1 – IOR Sequential

Pattern	NXD	Nodes	PPN	Procs	File Size (MB)	Transfer (KB)	Iterations	Write (MB/s)	Read (MB/s)	Write (MiB/s)	Read (MiB/s)
Sequential	Disabled	16	2	32	32,768	1024	5	13,143	11,182	12,535	10,664
Sequential	Enabled	16	2	32	32,768	1024	5	13,014	11,154	12,411	10,637

- No significant decrease in write or read performance
 - Small trend of slightly slower performance with NXD
- Mean impact over 5 iterations < 1%

Experiment 2 – Mixed IOR



Experiment 3 – Concurrent IOR, Increasing Load of Random



Experiment 3 – Concurrent IOR, Increasing Load of Random



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Experiment 4 – SWMR



SWMR - Write

Customer Test Runs – Production Codes and Data



nsive 4M-8 ck and a	3M – 99.9%	< 4K – 0 1%		
nsive			Minimal	No degradation
ck 4M-8	3M – 99.9%	< 4K – 0.1%	Perf gain 2.6% to 6%	Tuning down StripeSize from 4M to 1M and the gains started showing up.
nsive 1M – de	- 90%	< 512K – 10%	Perf gain of 10%	Performance improvement with NXD enabled / Histogram OFF and StripeSize – 1M / RPC 512/7
nsive 10% k	large reads	90% Small Read	Perf gain of 6%	Performance improvement with NXD enabled. No tuning, no pre-load
nsive 10% x large	random block	90% small Read (mostly seq.)	Perf gain of 10% over a run of continuous 100 iterations.	CS performs around 10% Sample size is only 9GiB
al runs 4 MB runs 4 MB	3 for seq	4K for random	Read: 10x increase Write: 6x increase	No degradation
nsive 1 MB ks	3	Mostly 4K	1% with stripe count 1 and stripe size 4m	No degradation over 7 runs
	nsive 1M – hsive 10% hsive 10% k 10% hsive 10% large ll runs 4 ME runs 1 ME	Image: size deImage: point point langeImage: size de10% large readsImage: size de10% random large blockImage: size runs4 MB for seqImage: size hasive ks1 MB	ININ90%< 512K - 10%IN90%Small ReadINSIVE RK10% large reads90% Small ReadINSIVE RK10% random large block90% small Read (mostly seq.)I runs runs4 MB for seq4K for randomI sive runs1 MBMostly 4K	Insive de1M – 90%< 512K – 10%Perf gain of 10%Insive okk10% large reads90% Small ReadPerf gain of 6%Insive okk10% random large block90% small Read (mostly seq.)Perf gain of 10% over a run of continuous 100 iterations.I runs runs4 MB for seq4K for randomRead: 10x increase Write: 6x increaseI sive ks1 MBMostly 4K1% with stripe count 1 and stripe size 4m

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Conclusions



- NXD is transparent to 100% streaming I/O
- NXD can preferentially enhance small block I/O by orders of magnitude with only a modest impact on sequential I/O during mixed I/O workloads
- Performance improvement is very application dependent
 - Careful I/O pattern analysis is required for optimal tuning
 - Example: NXD improved tar -c on a large directory by 43% whereas tar -x was unaffected

• Still limited field experience of NXD

• But the number of customers implementing NXD is growing rapidly

Torben Kling Petersen, PhE tpetersen@cray.com

Q&A

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