

Hewlett Packard Enterprise

Optimizing the Cray Graph Engine for Performant Analytics on Cluster, SuperDome Flex, Shasta Systems and Cloud Deployment

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Outline

Background

• Cray Graph Engine (CGE)

• Enabling execution on multiple platforms

• Performant execution across architectures

Recent Optimizations

- Database build
- Query Engine

Deployment Options

• Native and cloud deployments

Benchmarks

- Dictionary build
- Query execution
- Summary

Background : The Cray Graph Engine

- CGE is a scalable parallel graph analytics framework
 - An in-memory semantic graph database
 - Basic graph pattern search
 - Graph-theoretic (whole graph) algorithms
 - A W3C Standards inspired triplestore
 - Uses RDF Data model
 - Uses SPARQL as query language
- Scales-out and scales-up, i.e. built for "vertical scaling" based on parallel and distributed computing principles
 - Most market competitors are scale-out databases.
 - Scale-up => faster with more compute resources
- Fast turn-around on queries, brings interactivity to graph-based discovery
 - Scaling and performance enables interactive analysis of very large datasets (1 – 100s of TBs)

2012: YarcData is formed to build Big Data solutions

2014: Urika-GD – an appliance for graph analytics using the XMT architecture

2016: Urika-GX – graph engine code ported to work on Aries interconnect and x86 architecture.

2017: Urika-XC – graph engine code scales on supercomputers.

2018: Demonstrated scaling and showed CGE as the best performing database on a trillion triples.

2020 : Ongoing work on porting Graph engine ported for HPE hardware architectures such as ProLiant, SuperDomeFlex, InfiniBand, Shasta and HPE Container Platform.

Enabling CGE on Multiple Architectures

• CGE for XC based on Cray Coarray-C++

• Coarray-C++ built on Cray PGAS which leverages XPMEM/DMAPP and Aries interconnect

• Changes required for multi-platform support:

- Developed simplified PGAS library to remove dependence on Cray hardware and software stack
 - Based upon POSIX shared memory and one-sided MPI communication
 - Implements only functionality required by CGE, such as:
 - Creation of symmetric virtual address space across images and symmetric memory allocations/frees
 - Blocking and non-blocking puts/gets
 - Barriers
 - Collectives
 - Direct memory access amongst images on same physical node
- Modified Coarray-C++ to leverage simplified PGAS library
- Advantages of using POSIX shared memory and MPI
 - Portability
 - Performance across architectures

CGE PGAS Library

POSIX shared memory used for symmetric heap

- All images use shared memory to allocate symmetric heap
 - Used for all allocations in CGE, including memory for other allocators CGE maintains
- Enables images on same node to directly read/write to each other's memory without XPMEM

• MPI used for PMI, synchronization and RMA

- MPI window feature used to enable RMA access to symmetric heap
- CGE put/get functions translated into memcpy if on same node or MPI_Put/MPI_Get if to remote node
 - If non-blocking put/get, CGE will track outstanding RMA until next atomic_image_fence()
- atomic_image_fence() implemented using MPI_Win_flush_local for necessary ranks
- sync_all() implemented as an atomic_image_fence() and MPI_Barrier()
- Using MPI enables containerization with Singularity
 - Container, cloud, Infiniband clusters, SuperDome Flex, and Shasta Cray EX supercomputers
 - Without significant loss in performance across architectures

CGE Coarray-C++ Template File

Only implemented features used by CGE

• Removed features such as cofuture and coatomic

• Replaced calls to Cray PGAS library with new versions to CGE PGAS

• New API mimics Cray PGAS which simplified replacement

• New functionality:

- coexchange(): templated utility function provided that performs all-to-all communication of single data elements in a group-aggregated manner
- Separate image fence functions for puts/gets:
 - atomic_image_put_fence()
 - atomic_image_get_fence()
- coexchange and separate put/get fences required corresponding support be added to CGE PGAS library

Recent Optimizations

Dictionary updated to support block-fetching of strings

- Given an array of IDs dictionary uses group-aggregated block fetching of strings
- Used by query results writing, FILTER, Inferencing, etc.

• Inferencer updated to use dictionary block fetching of strings

- Strings used to verify quads being inferred
- Reduced inferencing time of LUBM100K on 512 images from 1796.6 seconds to 158.8 seconds

• Scan/Join operators

- Optimized for shared memory machines to avoid data marshalling
- Modified group-aggregated messaging so images hold data for consecutive images
 - Previously aggregated data from group based on one partner image per group O(num_groups) messages
 - Reduces number of groups an image communicates with -- O(num_groups / group_size) messages
- Reduced message sizes by returning bit-flag for 1/0 statuses rather than full words
- Uses separate fence for remote get to prevent blocking on put of status bits
- LUBM100K Query 9 time on 32 nodes x 16 images per node reduced from 8.6 to 7.0 seconds

Deployment Options

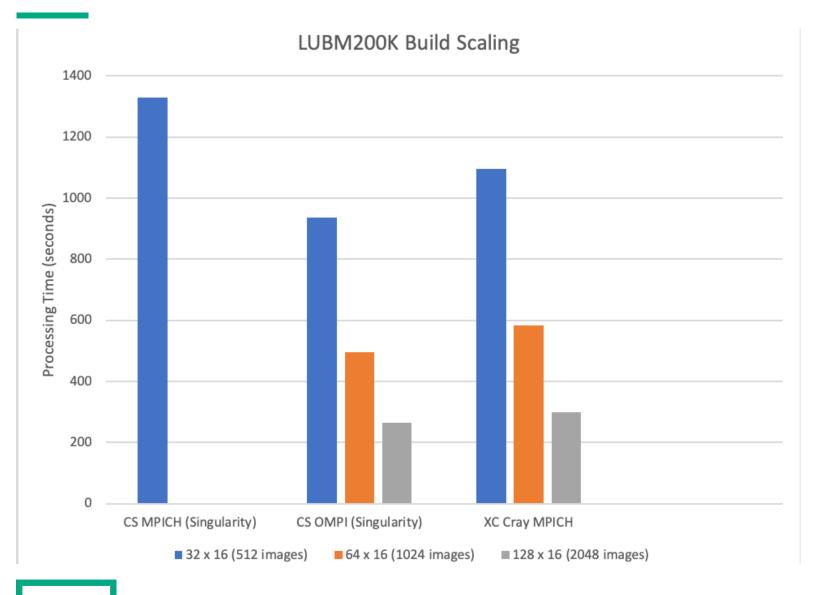
Build and execute natively using Cray MPICH

- Method used for executing on XC and EX
- Enables optimal performance for given hardware using Cray MPI libraries

Containerization using Singularity

- Two models for MPI applications:
 - Bind host MPI bind-mounted into container
 - Hybrid MPI built into container and host MPI interacts with container MPI
- CGE uses hybrid model for portability but requires container MPI to be configured for performance
 - Container includes: Mellanox OFED, UCX/OFI and MPI
- Containers created for both OpenMPI and MPICH
 - OpenMPI 4.1 using UCX
 - Requires host OpenMPI install
 - MPICH 3.4 using OFI
 - Easy execution using slurm
 - Sometimes fails with OFI timeouts

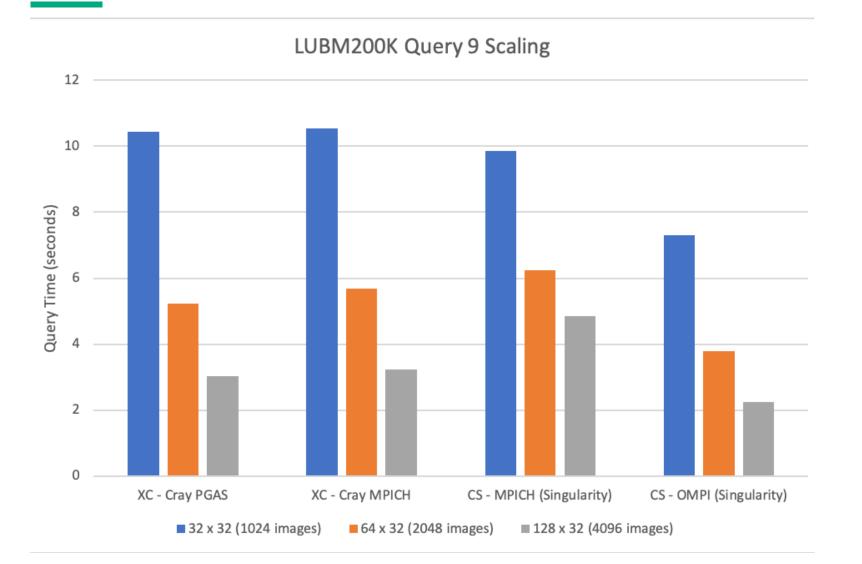
Dictionary Build Step Scaling



• Database build time excluding file I/O

- Parse, Sync, Sort, Update and Inference
- Scales well by node count
- CGE using OMPI in Singularity is fastest
- CGE using MPICH fails to build using > 32 nodes
 - OFI timeout errors

LUBM200K Query 9 Scaling



- Strong scaling for all deployments except CS-MPICH
 - 64 to 128 nodes had reduced scaling
- Performance with Cray MPICH matches Cray PGAS
- CGE on CS using OMPI is fastest
- CGE on CS using MPICH required new coexchange
 - OFI timeouts for small all-to-all puts/gets

Query Performance on EX, AWS and SDF

• Benchmarked on AWS and Cray EX using LUBM100K (~18.2 billion quads)

- AWS parallel cluster with c5n.18xlarge nodes
- Cray EX dual socket nodes with AMD 64-core EPYC processors
- 32 nodes for both with either 16 or 32 images per node

Computing Platform	MPI	Query 9 time in seconds	
		512 images	1024 images
AWS efa cluster (singularity)	OpenMPI	11.60	8.40
Cray EX with Slingshot	Cray MPICH	6.17	4.24

• SuperDome Flex (SDF) with LUBM25K (~4.6 billion quads)

Computing Platform	MPI	Query 9 time in seconds	
		96 images	192 images
8 socket SDF (singularity)	MPICH	16.45	12.81

Summary

- Described effort for enabling CGE to execute on multiple platforms
 - Simplified PGAS and Coarray-C++ built on top of POSIX shared memory and one-sided MPI
 - Portability and performance using POSIX and MPI
- Recent optimizations to CGE
 - Dictionary build, inferencing and variable bindings
- Demonstrated dictionary build and query performance and scaling
 - CGE MPI based PGAS performance matches Cray PGAS
 - Containerized CGE performance on CS can match or exceed native XC performance
- CGE now deployable either natively or in a Singularity container
 - Native version for XC/EX enables leveraging Cray MPICH
 - Container enables easy deployment across platforms as well as in the cloud



• Questions?